CSE 311: Foundations of Computing

Lecture 23: FSMs with Output, Minimization





Vending Machine



Enter 15 cents in dimes or nickels Press S or B for a candy bar



Vending Machine, v0.1



B, S

Basic transitions on N (nickel), D (dime), B (butterfinger), S (snickers)

Vending Machine, v0.2



Adding output to states: N – Nickel, S – Snickers, B – Butterfinger

Vending Machine, v1.0



Adding additional "unexpected" transitions to cover all symbols for each state

- Many different FSMs (DFAs) for the same problem
- Take a given FSM and try to reduce its state set by combining states
 - Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this

State Minimization Algorithm

- Put states into groups
- Try to find groups that can be collapsed into one state
 - states can keep track of information that isn't necessary to determine whether to accept or reject
- Group states together until we can prove that collapsing them can change the accept/reject result
 - find a specific string x such that:

starting from state A, following edges according to x ends in accept starting from state B, following edges according to x ends in reject

– (algorithm below could be modified to show these strings)

State Minimization Algorithm

1. Put states into groups based on their outputs (whether they accept or reject)

State Minimization Algorithm

- 1. Put states into groups based on their outputs (whether they accept or reject)
- 2. Repeat the following until no change happens
 - a. If there is a symbol a so that not all states in a group
 G agree on which group a leads to, split G into smaller
 groups based on which group the states go to on a





3. Finally, convert groups to states



present		nex	t stat	output	
state	0	1	2	3	
SO	SO	S1	S2	S3	1 accept
S1	SO	S3	S1	S5	0 reject
S2	S1	S3	S2	S4	1 1
S3	S1	S0	S4	S5	0
S4	SO	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)



present	1	next	output		
state	0	1	2	3	
- S0	SO	S1	S2	S3	1
S1	SO	S3	S1	S5	0
- S2	S1	S3	S2	S4	1
S3	S1	SO	S4	S5	0
— S4	SO	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)



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SO	SO	S1	S2	S3	1
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S2	S1	S3	S2	S4	1
S3	S1	SO	S4	S5	0
S4	SO	S1	S2	S5	1
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S3	S1	SO	S4	S5	0
S4	SO	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

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S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)





Finally convert groups to states:

Can combine states S0-S4 and S3-S5.

In table replace all S4 with S0 and all S5 with S3

Minimized Machine



present	1	next	output					
state	0	1	2	3				
SO	SO	S1	S2	S3	1			
S1	SO	S3	S1	S3	0			
S2	S1	S3	S2	SO	1			
S3	S1	SO	SO	S3	0			
state								
transition table								

A Simpler Minimization Example



The set of all binary strings with # of 1's \equiv # of 0's (mod 2).

A Simpler Minimization Example



Split states into accept/reject groups

Every symbol causes the DFA to go from one group to the other so neither group needs to be split

Minimized DFA



The set of all binary strings with # of 1's \equiv # of 0's (mod 2).

= The set of all binary strings with even length.

Nondeterministic Finite Automata (NFA)

- Graph with start state, final states, edges labeled by symbols (like DFA) but
 - Not required to have exactly 1 edge out of each state
 labeled by each symbol— can have 0 or >1
 - Also can have edges labeled by empty string $\boldsymbol{\epsilon}$
- Definition: x is in the language recognized by an NFA if and only if <u>some</u> valid execution of the machine gets to an accept state



Consider This NFA



Consider This NFA



What language does this NFA accept?

10(10)* U 111 (0 U 1)*

NFA ϵ -moves

01 ...



NFA ε-moves

Strings over {0,1,2} w/even # of 2's OR sum to 0 mod 3



NFA for set of binary strings with a 1 in the 3rd position from the end



Compare with the smallest DFA



- Generalization of DFAs
 - drop two restrictions of DFAs
 - every DFA is an NFA
- Seem to be more powerful
 - designing is easier than with DFAs

• Seem related to regular expressions