## CSE 311: Foundations of Computing

Lecture 22: Finite State Machines


## Last class: Strings this machine says are OK?



The set of all binary strings that end in 0

## Finite State Machines

- States
- Transitions on input symbols
- Start state and final states
- The "language recognized" by the machine is the set of strings that reach a final state from the start

| Old State | 0 | 1 |
| :---: | :---: | :---: |
| $\mathrm{~s}_{0}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{1}$ |
| $\mathrm{~s}_{1}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{2}$ |
| $\mathrm{~s}_{2}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{3}$ |
| $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ |



Finite State Machines

- Each machine designed for strings over some fixed alphabet $\Sigma$.
- Must have a transition defined from each state for every symbol in $\Sigma$.

| Old State | 0 | 1 |
| :---: | :---: | :---: |
| $\mathrm{~s}_{0}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{1}$ |
| $\mathrm{~s}_{1}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{2}$ |
| $\mathrm{~s}_{2}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{3}$ |
| $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ |



## What language does this machine recognize?

| Old State | 0 | 1 |
| :---: | :---: | :---: |
| $\mathrm{~s}_{0}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{1}$ |
| $\mathrm{~s}_{1}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{2}$ |
| $\mathrm{~s}_{2}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{3}$ |
| $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ |



## What language does this machine recognize?

The set of all binary strings that contain 111
or don't end in 1

| Old State | 0 | 1 |
| :---: | :---: | :---: |
| $\mathrm{~s}_{0}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{1}$ |
| $\mathrm{~s}_{1}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{2}$ |
| $\mathrm{~s}_{2}$ | $\mathrm{~s}_{0}$ | $\mathrm{~s}_{3}$ |
| $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ | $\mathrm{~s}_{3}$ |



Strings over $\{0,1,2\}$
$M_{1}$ : Strings with an even number of 2's


$$
\begin{aligned}
& \text { So: strings ileum } 4 \text { of } 25 \\
& S_{1}:- \text { odd } 4 \text { of } 25
\end{aligned}
$$

Strings over $\{0,1,2\}$
$M_{1}$ : Strings with an even number of 2's


## State Machine Design Recipe

Given a language, how do you design a state machine for it?

Create states to remember enough
(about the portion of the input string that it has already seen) to correctly answer "accept/reject" on the whole string after seeing the rest.

Add labeled edges to show how the memory (state) should be updated for each new symbol.

Strings over $\{0,1,2\}$
$M_{2}$ : Strings where the sum of digits $\bmod 3$ is 0


So: Strings whose Sum of digits mols iso
$s_{1}$ : Strings
$S_{2}$ : Strings $\qquad$

Strings over $\{0,1,2\}$
$M_{2}$ : Strings where the sum of digits $\bmod 3$ is 0

Strings over $\{0,1,2\}$
$M_{2}$ : Strings where the sum of digits $\bmod 3$ is 0


What language does this machine recognize?


- even \# of digits (length)
- even \#of Os and evenatas or add \# OS and ad el $-10^{*} \cup 01^{*}$
001000 events ofd\#1s
- 2 ne bullet pons and Cunt have more than 2 consecutive Os or 1 s


## What language does this machine recognize?



The set of all binary strings with \# of 1's 三 \# of 0's (mod 2)
(both are even or both are odd).

Can you think of a simpler description?

Strings over $\{0,1,2\}$
$M_{1}$ : Strings with an even number of 2's

$M_{2}$ : Strings where the sum of digits $\bmod 3$ is 0


Strings over $\{0,1,2\} \mathrm{W}$ / even number of 2 's and mod 3 sum 0


Strings over $\{0,1,2\} \mathrm{w}$ / even number of 2s and mod 3 sum 0


Strings over $\{0,1,2\} \mathbf{w} /$ even number of $2, s \operatorname{OR} \bmod 3$ sum 0


The set of binary strings with a 1 in the $3^{\text {rd }}$ position from the start

The set of binary strings with a 1 in the $3^{\text {rd }}$ position from the start


The set of binary strings with a 1 in the $3^{\text {rd }}$ position from the end

3 bit shift register "Remember the last three bits"


The set of binary strings with a 1 in the $3^{\text {rd }}$ position from the end


The set of binary strings with a 1 in the $3^{\text {rd }}$ position from the end


The beginning versus the end


