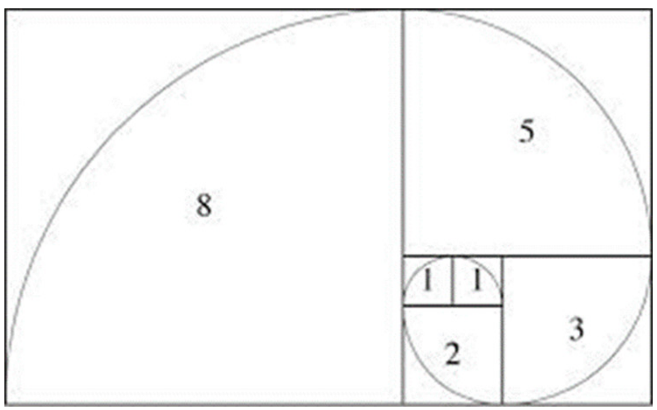
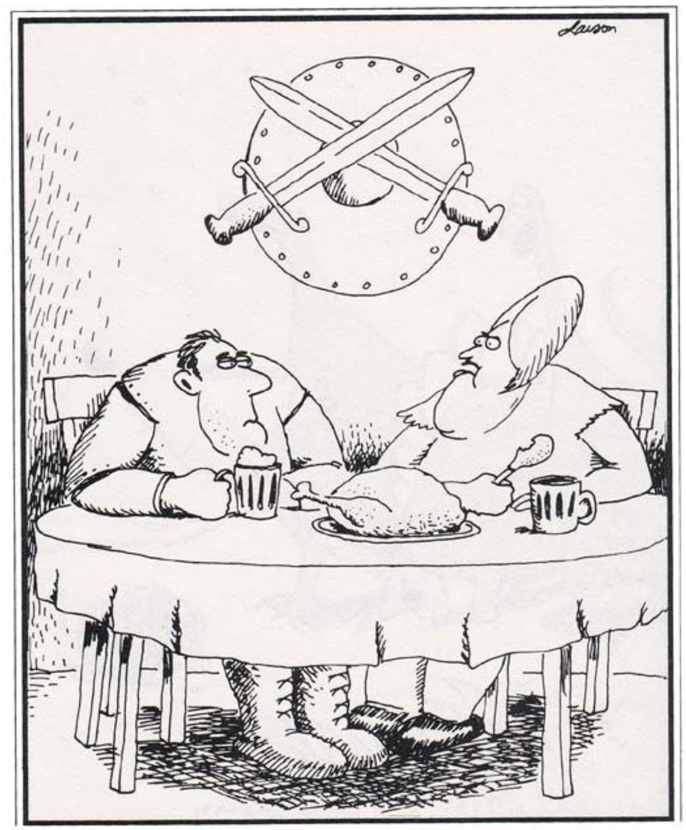
CSE 311: Foundations of Computing

Lecture 15: Recursion & Strong Induction
Applications: Fibonacci & Euclid





"And another thing . . . I want you to be more assertive! I'm tired of everyone calling you Alexander the Pretty-Good!"

See Edstem post about 1-1 meetings with TAs not about current HW

Last class: Inductive Proofs In 5 Easy Steps

1. "Let $P(n)$ be... . We will show that $P(n)$ is true for all integers $n \geq b$ by induction."
2. "Base Case:" Prove $P(b)$
3. "Inductive Hypothesis:

   Assume that for some arbitrary integer $k \geq b$,

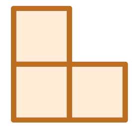
   $P(k)$ is true"
4. "Inductive Step:" Prove that $P(k+1)$ is true:

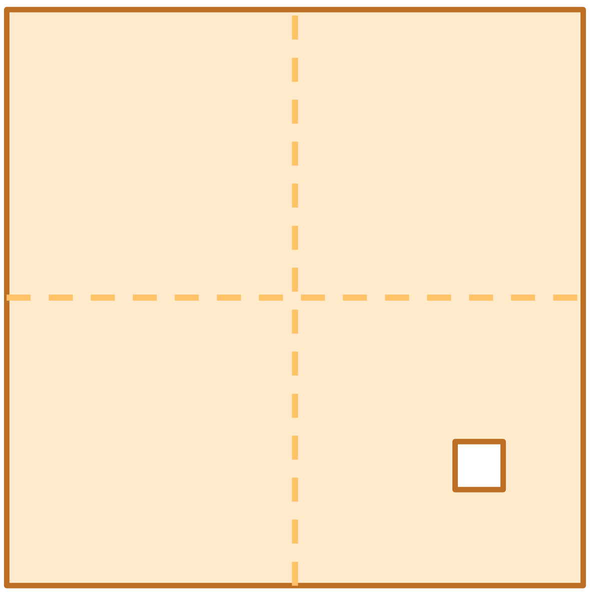
   Use the goal to figure out what you need.

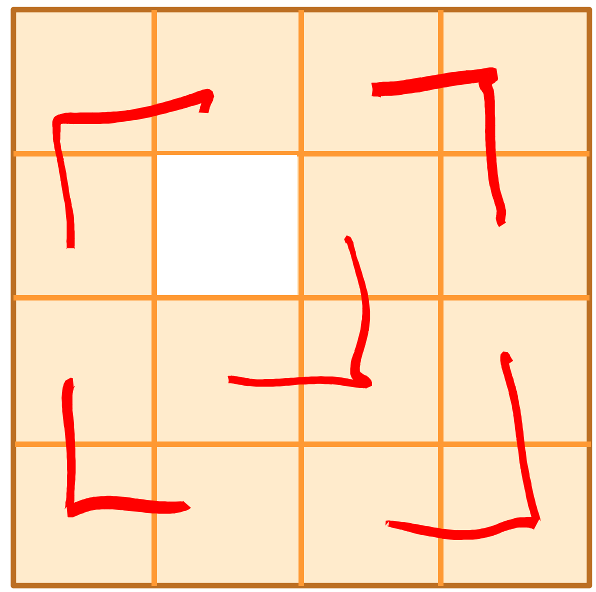
   Make sure you are using I.H. and point out where you are using it. (Don't assume $P(k+1)$!!)

5. "Conclusion: $P(n)$ is true for all integers $n \geq b$"

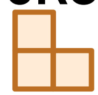
Checkerboard Tiling

- Prove that a $2^n \times 2^n$ checkerboard with one square removed can be tiled with: 



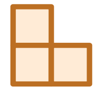


Checkerboard Tiling

1. **Let** P(n) **be "Any** $2^n \times 2^n$ **checkerboard with one square removed can be tiled with**  **".**
   We prove P(n) **for all** $n \geq 1$ **by induction on** n.

Base Case $(n=1)$: $2^1 \times 2^1$

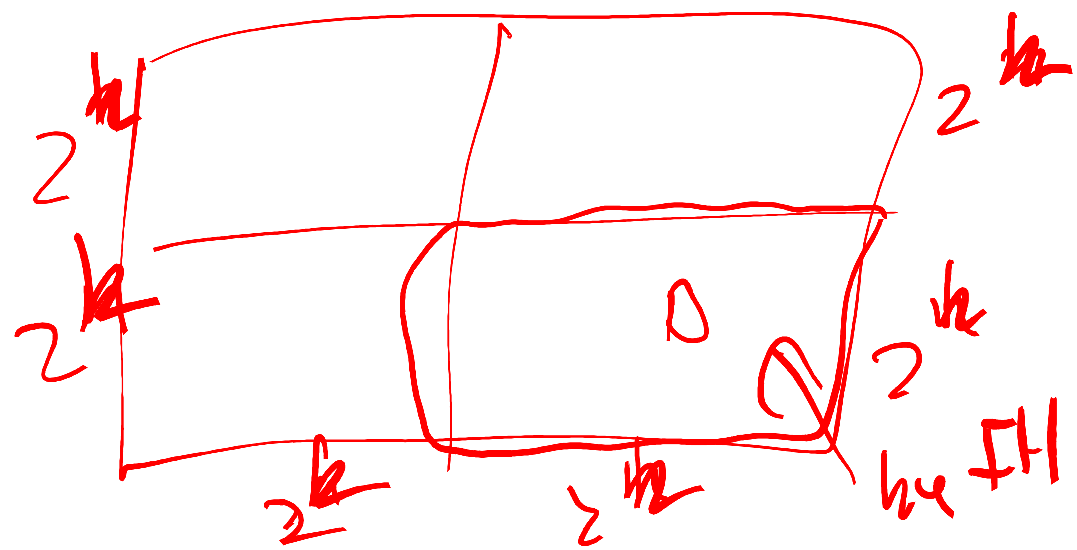
Checkerboard Tiling

1. **Let** P(n) **be "Any** $2^n \times 2^n$ **checkerboard with one square removed can be tiled with**  **" .**
   We prove P(n) **for all** $n \geq 1$ **by induction on** n.
2. **Base Case:** n=1 

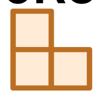
3. Induction Hypothesis: Suppose that P(k) is true for some arbitrary $k \geq 1$

P(k+1):

$2^{k+1} \times 2^{k+1}$



Checkerboard Tiling

1. **Let** $P(n)$ **be "Any** $2^n \times 2^n$ **checkerboard with one square removed can be tiled with**  **" .**
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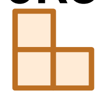
2. **Base Case:** $n=1$ 
3. **Inductive Hypothesis: Assume** $P(k)$ **for some arbitrary integer** $k \geq 1$

Checkerboard Tiling

1. **Let** $P(n)$ **be "Any** $2^n \times 2^n$ **checkerboard with one square removed can be tiled with**  **".**
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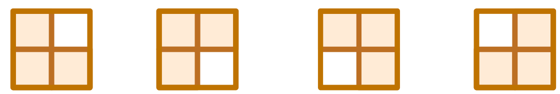
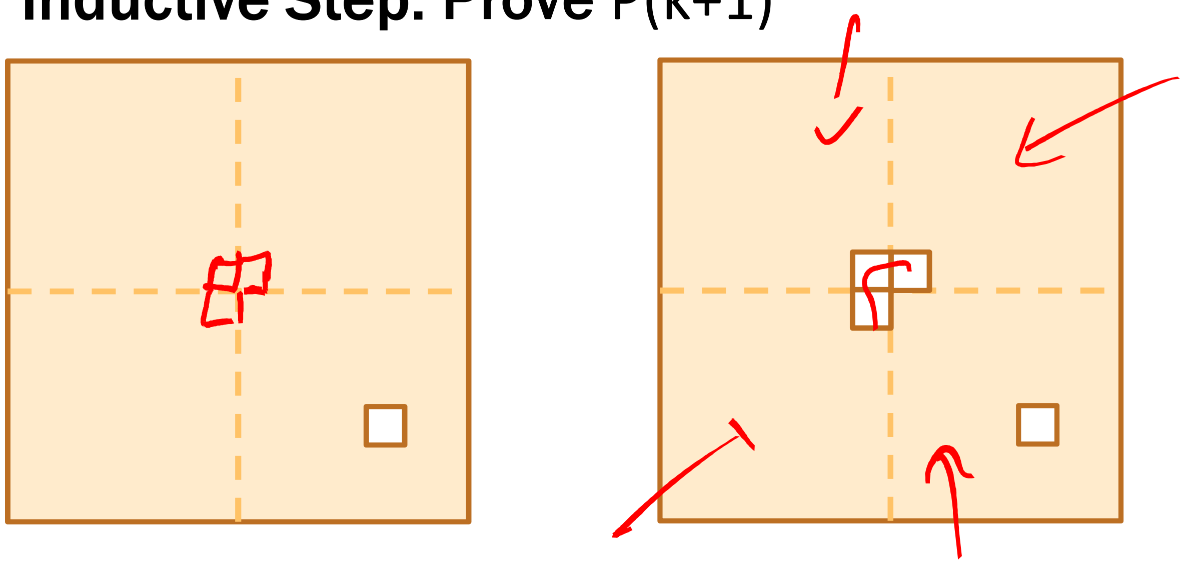
2. **Base Case:** $n=1$ 
3. **Inductive Hypothesis: Assume** $P(k)$ **for some arbitrary integer** $k \geq 1$
4. **Inductive Step: Prove** $P(k+1)$

$2^{2n}-1 \equiv 0 \pmod 3$



Apply IH to each quadrant then fill with extra tile.

Recall: Induction Rule of Inference

Domain: Natural Numbers

$$\begin{array}{c} P(0) \\ \forall k\ (P(k) \longrightarrow P(k+1)) \\ \hline \therefore \forall n\ \ P(n) \end{array}$$

How do the givens prove P(5)?

$$P(0) \xrightarrow{P(0)\to P(1)} P(1) \xrightarrow{P(1)\to P(2)} P(2) \xrightarrow{P(2)\to P(3)} P(3) \xrightarrow{P(3)\to P(4)} P(4) \xrightarrow{P(4)\to P(5)} P(5)$$

Recall: Induction Rule of Inference

Domain: Natural Numbers

$$\begin{array}{c} P(0) \\ \forall k\ (P(k) \longrightarrow P(k+1)) \\ \hline \therefore \forall n\ \ P(n) \end{array}$$

How do the givens prove P(5)?

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We made it harder than we needed to …

When we proved $P(2)$ **we knew BOTH** $P(0)$ **and** $P(1)$

When we proved $P(3)$ **we knew** $P(0)$ **and** $P(1)$ **and** $P(2)$

When we proved $P(4)$ **we knew** $P(0)$, $P(1)$, $P(2)$, $P(3)$

etc.

That's the essence of the idea of Strong Induction.

Strong Induction

$$\frac{P(0) \qquad \forall k \left(\overbrace{\forall j \left(0 \le j \le k \to P(j)\right)}^{\text{all previous values}} \to \overset{\text{next value}}{P(k+1)} \right)}{\therefore \forall n\, P(n)}$$

Strong Induction

$$\frac{P(0) \qquad \forall k \left(\overbrace{\forall j\,(0 \le j \le k \rightarrow P(j))}^{Q(n)} \rightarrow P(k+1) \right)}{\therefore\ \forall n\, P(n)}$$

Strong induction for P follows from ordinary induction for Q where

$$Q(k) := \ \forall j\,(0 \le j \le k \rightarrow P(j))$$

Note that $Q(0) = P(0)$ and $Q(k+1) \equiv Q(k)\ \wedge P(k+1)$ and $\forall n\, Q(n) \equiv \forall n\, P(n)$

Last class: Inductive Proofs In 5 Easy Steps

1. "Let $P(n)$ be… . We will show that $P(n)$ is true for all integers $n \geq b$ by induction."
2. "Base Case:" Prove $P(b)$
3. "Inductive Hypothesis:

   Assume that for some arbitrary integer $k \geq b$,

   $P(k)$ is true"
4. "Inductive Step:" Prove that $P(k+1)$ is true:

   Use the goal to figure out what you need.

   Make sure you are using I.H. and point out where you are using it. (Don't assume $P(k+1)$!!)

5. "Conclusion: $P(n)$ is true for all integers $n \geq b$"

Strong Inductive Proofs In 5 Easy Steps

1. "Let $P(n)$ be... . We will show that $P(n)$ is true for all integers $n \geq b$ by *strong* induction."
2. "Base Case:" Prove $P(b)$
3. "Inductive Hypothesis:

   Assume that for some arbitrary integer $k \geq b$,

   $P(j)$ *is true for every integer* j *from* b to k"
4. "Inductive Step:" Prove that $P(k+1)$ is true:

   Use the goal to figure out what you need.

   Make sure you are using I.H. (that $P(b), \dots, P(k)$ are true) and point out where you are using it.
   (Don't assume $P(k+1)$!!)
5. "Conclusion: $P(n)$ is true for all integers $n \geq b$"

Recall: Fundamental Theorem of Arithmetic

Every integer > 1 has a unique prime factorization

48 = 2 • 2 • 2 • 2 • 3
591 = 3 • 197
45,523 = 45,523
321,950 = 2 • 5 • 5 • 47 • 137
1,234,567,890 = 2 • 3 • 3 • 5 • 3,607 • 3,803

We use strong induction to prove that a factorization into primes exists, but not that it is unique.

Every integer ≥ 2 is a product of (one or more) primes.

Every integer ≥ 2 is a product of (one or more) primes.

1. **Let** P(n) **be** "n is a product of some list of primes". **We will show that** P(n) **is true for all integers** n $\geq$ 2 **by strong induction.**

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1. **Let** P(n) **be** "n is a product of some list of primes". **We will show that** P(n) **is true for all integers** n $\geq$ 2 **by strong induction.**
2. **Base Case** (n=2): 2 **is prime, so it is a product of (one) prime.**
   Therefore P(2) **is true.**

Inductive Hypothesis: Assume that for some arbitrary integer $k \geq 2$

That $P(j)$ is true for all integers j with $2 \leq j \leq k$

i.e. that every integer from 2 to k is a product of primes.

Every integer ≥ 2 is a product of (one or more) primes.

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3. **Inductive Hyp: Suppose that for some arbitrary integer** k ≥ 2, P(j) **is true for every integer** j **between** 2 **and** k

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   Goal: Show P(k+1); **i.e.** k+1 **is a product of primes**

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   <u>**Case:** k+1 **is prime**</u>**: Then by definition** k+1 **is a product of primes**

   <u>**Case:** k+1 **is composite:**</u> **Then** k+1=ab **for some integers** a **and** b **where** 2 ≤ a, b ≤ k.

by IH. P(a) and P(b) both true

Every integer ≥ 2 is a product of (one or more) primes.

1. **Let** P(n) **be** "n is a product of some list of primes". **We will show that** P(n) **is true for all integers** $n \geq 2$ **by strong induction.**
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   Case: k+1 is composite: Then k+1=ab **for some integers** a **and** b **where** $2 \leq a, b \leq k$. **By our IH,** P(a) **and** P(b) **are true so we have**

   $$a = p_1 p_2 \cdots p_r \text{ and } b = q_1 q_2 \cdots q_s$$

   for some primes $p_1, p_2, \ldots, p_r, q_1, q_2, \ldots, q_s$.

   Thus, $k+1 = ab = p_1 p_2 \cdots p_r q_1 q_2 \cdots q_s$ **which is a product of primes.**

   Since $k \geq 2$, **one of these cases must happen and so** P(k+1) **is true.**

Every integer ≥ 2 is a product of (one or more) primes.

1. **Let** $P(n)$ **be** "n is a product of some list of primes". **We will show that** $P(n)$ **is true for all integers** $n \geq 2$ **by strong induction.**
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   $\boxed{\textbf{Goal: Show } P(k+1)\textbf{; i.e. } k+1 \textbf{ is a product of primes}}$

   <u>Case: k+1 is prime</u>: Then by definition $k+1$ **is a product of primes**

   <u>Case: k+1 is composite:</u> Then $k+1=ab$ **for some integers** a **and** b **where** $2 \leq a, b \leq k$. **By our IH,** $P(a)$ **and** $P(b)$ **are true so we have**

   $a = p_1p_2 \cdots p_r$ **and** $b = q_1q_2 \cdots q_s$

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   Thus, $k+1 = ab = p_1p_2 \cdots p_rq_1q_2 \cdots q_s$ **which is a product of primes.**

   Since $k \geq 2$, **one of these cases must happen and so** $P(k+1)$ **is true.**
5. **Thus** $P(n)$ **is true for all integers** $n \geq 2$, **by strong induction.**

Strong Induction is particularly useful when…

…we need to analyze methods that on input k make a recursive call for an input different from $k - 1$.

e.g.: Recursive Modular Exponentiation:

- For exponent $k > 0$ it made a recursive call with exponent $\mathrm{j} = k/2$ when k was even or $\mathrm{j} = k - 1$ when k was odd.

Fast Exponentiation

```
public static int FastModExp(int a, int k, int modulus) {

        if (k == 0) {
            return 1;

         } else if ((k % 2) == 0) {
            long temp = FastModExp(a,k/2,modulus);
            return (temp * temp) % modulus;

        } else {
            long temp = FastModExp(a,k-1,modulus);
            return (a * temp) % modulus;
        }

}
```

$$a^{2j} \bmod m = \left(a^{j} \bmod m\right)^{2} \bmod m$$

$$a^{2j+1} \bmod m = \big((a \bmod m) \cdot (a^{2j} \bmod m)\big) \bmod m$$

Strong Induction is particularly useful when…

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e.g.: Recursive Modular Exponentiation:

- For exponent $k > 0$ it made a recursive call with exponent $\mathrm{j} = k/2$ when k was even or $\mathrm{j} = k-1$ when k was odd.

We won't analyze this particular method by strong induction, but we could.

However, we will use strong induction to analyze other functions with recursive definitions.

Recursive definitions of functions

- $\boxed{0! = 1;}$ $(n+1)! = (n+1)\cdot n!$ **for all** $n \geq 0$.

- $F(0) = 0$; $F(n+1) = F(n) + 1$ **for all** $n \geq 0$.

  $\therefore F(n) = n$

- $G(0) = 1$; $G(n+1) = 2 \cdot G(n)$ **for all** $n \geq 0$.

  $\therefore G(n) = 2^n$

- $H(0) = 1$; $H(n+1) = 2^{H(n)}$ **for all** $n \geq 0$.

  $H(n) = 2^{2^{\cdot^{\cdot^{\cdot^{2}}}}}$ height n

Prove $n! \leq n^n$ for all $n \geq 1$

$(n+1)! = (n+1) \cdot n!$
$(0+1)! = 1 \cdot 0!$

1. **Let** P(n) **be** "n! ≤ n^n". **We will show that** P(n) **is true for all integers** n ≥ 1 **by induction.**
2. **Base Case** (n=1): $1!=1\cdot 0!=1\cdot 1=1=1^1$ **so** P(1) **is true.**
3. **Inductive Hypothesis: Suppose that** P(k) **is true for some arbitrary integer** k ≥ 1**. I.e., suppose** $k! \leq k^k$.

Prove $n! \leq n^n$ for all $n \geq 1$

1. **Let** P(n) **be** "n! ≤ n^n". **We will show that** P(n) **is true for all integers** n ≥ 1 **by induction.**
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3. **Inductive Hypothesis: Suppose that** P(k) **is true for some arbitrary integer** k ≥ 1. **I.e., suppose** k! ≤ k^k.
4. **Inductive Step:**

   $\boxed{\textbf{Goal: Show } P(k+1)\textbf{, i.e. show } (k+1)! \leq (k+1)^{k+1}}$

$(k+1)!$	$= (k+1)\cdot k!$	**by definition of !**
	$\leq (k+1)\cdot k^k$	**by the IH**
	$\leq (k+1)\cdot (k+1)^k$	**since** k ≥ 0
	$= (k+1)^{k+1}$	

   Therefore P(k+1) **is true.**
5. **Thus** P(n) **is true for all** n ≥ 1, **by induction.**

More Recursive Definitions

Suppose that $h\colon \mathbb{N} \to \mathbb{R}$.

Then we have familiar summation notation:

$$\sum_{i=0}^{0} h(i) = h(0)$$

$$\sum_{i=0}^{n+1} h(i) = h(n+1) + \sum_{i=0}^{n} h(i) \quad \textbf{for } n \geq 0$$

There is also product notation:

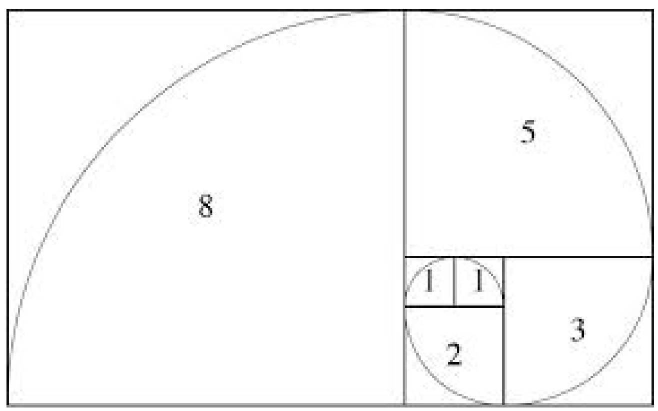
$$\prod_{i=0}^{0} h(i) = h(0)$$

$$\prod_{i=0}^{n+1} h(i) = h(n+1) \cdot \prod_{i=0}^{n} h(i) \quad \textbf{for } n \geq 0$$

$$n! = \prod_{i=1}^{n} i \qquad n \geq 1$$

Fibonacci Numbers

$$f_0 = 0$$
$$f_1 = 1$$
$$f_n = f_{n-1} + f_{n-2} \text{ for all } n \geq 2$$



$f_2 = 1$

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ... $f_3 = 2$ $f_4 = 3$ 5 8

Fibonacci Numbers

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golden mean

1.618 - -

Tamás Görbe
@TamasGorbe

A Mathematician's Way* of Converting Miles to Kilometers

1.609

$$3 \text{ mi} \approx 5 \text{ km}$$
$$5 \text{ mi} \approx 8 \text{ km}$$
$$8 \text{ mi} \approx 13 \text{ km}$$

f_n mi $\approx$ f_{n+1} km

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

1. **Let** P(n) **be** "$f_n < 2^n$". **We prove that** P(n) **is true for all integers** $n \geq 0$ **by strong induction.**

$$\boldsymbol{f_0 = 0 \quad f_1 = 1}$$
$$\boldsymbol{f_n = f_{n-1} + f_{n-2}} \text{ for all } \boldsymbol{n \geq 2}$$

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

1. Let P(n) be "$f_n < 2^n$". We prove that P(n) is true for all integers $n \geq 0$ by strong induction.
2. Base Case: $f_0=0 < 1= 2^0$ so P(0) is true.

$\boldsymbol{f_0 = 0} \quad \boldsymbol{f_1 = 1}$
$\boldsymbol{f_n = f_{n-1} + f_{n-2}}$ for all $\boldsymbol{n \geq 2}$

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

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2. **Base Case:** $f_0=0 < 1= 2^0$ **so** P(0) **is true.**
3. **Inductive Hypothesis: Assume that for some arbitrary integer** $k \geq 0$**, we have** $f_j < 2^j$ **for every integer** j **from** 0 **to** k**.**

P(k+1)

$f_0 = 0 \quad f_1 = 1$
$f_n = f_{n-1} + f_{n-2}$ for all $n \geq 2$

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

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4. **Inductive Step:** **Goal: Show** P(k+1); **that is,** $f_{k+1} < 2^{k+1}$

$$\boldsymbol{f_0 = 0 \qquad f_1 = 1}$$
$$\boldsymbol{f_n = f_{n-1} + f_{n-2}} \text{ for all } \boldsymbol{n \geq 2}$$

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

1. **Let** P(n) **be "**$f_n < 2^n$**". We prove that** P(n) **is true for all integers** $n \geq 0$ **by strong induction.**
2. **Base Case:** $f_0=0 < 1= 2^0$ **so** P(0) **is true.**
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4. **Inductive Step:** $\boxed{\textbf{Goal: Show } P(k+1)\textbf{; that is, } f_{k+1} < 2^{k+1}}$

   Case k+1 = 1: $f_1 = 1 < 2 = 2^1 = 2^{k+1}$

   Case k+1 ≥ 2:

$f_0 = 0 \quad f_1 = 1$

$f_n = f_{n-1} + f_{n-2}$ for all $n \geq 2$

Bounding Fibonacci I: $f_n < 2^n$ for all $n \geq 0$

1. **Let** P(n) **be "**$f_n < 2^n$**". We prove that** P(n) **is true for all integers** $n \geq 0$ **by strong induction.**
2. **Base Case:** $f_0=0 < 1= 2^0$ **so** P(0) **is true.**
3. **Inductive Hypothesis: Assume that for some arbitrary integer** $k \geq 0$, **we have** $f_j < 2^j$ **for every integer** j **from** 0 **to** k.
4. **Inductive Step:** $\boxed{\textbf{Goal: Show } P(k+1)\textbf{; that is, } f_{k+1} < 2^{k+1}}$

   Case k+1 = 1: **Then** $f_1 = 1 < 2 = 2^1$ **so** P(k+1) **is true here.**

   Case k+1 ≥ 2:

$$f_{k+1} = \underbrace{f_k}_{< 2^k} + \underbrace{f_{k-1}}_{< 2^{k-1}}$$

by def'n

$n \geq 0$

$$< 2^k + 2^{k-1} < 2^k + 2^k = 2^{k+1}$$

$f_0 = 0 \quad f_1 = 1$

$f_n = f_{n-1} + f_{n-2}$ for all $n \geq 2$

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Didn't use IH

<u>Case k+1 = 1</u>: **Then** $f_1 = 1 < 2 = 2^1$ **so** P(k+1) **is true here.**

<u>Case k+1 ≥ 2</u>: **Then** $f_{k+1} = f_k + f_{k-1}$ **by definition**

$< 2^k + 2^{k-1}$ **by the IH since** $k-1 \geq 0$

$< 2^k + 2^k = 2 \cdot 2^k$

$= 2^{k+1}$

so P(k+1) **is true in this case.**

These are the only cases so P(k+1) **follows.**

$$\boldsymbol{f_0 = 0 \quad f_1 = 1}$$
$$\boldsymbol{f_n = f_{n-1} + f_{n-2}} \text{ for all } \boldsymbol{n \geq 2}$$

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   <u>Case</u> <u>$k+1 = 1$</u>**: Then** $f_1 = 1 < 2 = 2^1$ **so** P(k+1) **is true here.**

   <u>Case</u> <u>$k+1 \geq 2$</u>**: Then** $f_{k+1} = f_k + f_{k-1}$ **by definition**

   $< 2^k + 2^{k-1}$ **by the IH since** $k-1 \geq 0$

   $< 2^k + 2^k = 2 \cdot 2^k = 2^{k+1}$

   so P(k+1) **is true in this case.**

   These are the only cases so P(k+1) **follows.**
5. **Therefore by strong induction,**

   $f_n < 2^n$ **for all integers** $n \geq 0$.

$$f_0 = 0 \quad f_1 = 1$$
$$f_n = f_{n-1} + f_{n-2} \text{ for all } n \geq 2$$