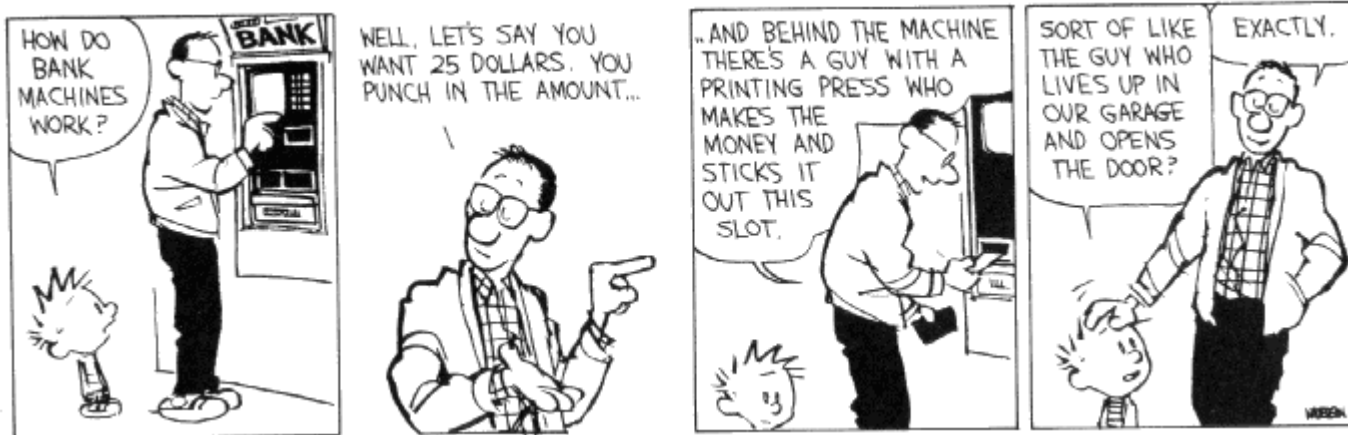


# CSE 311: Foundations of Computing

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## Lecture 21: DFAs and Finite State Machines with Output

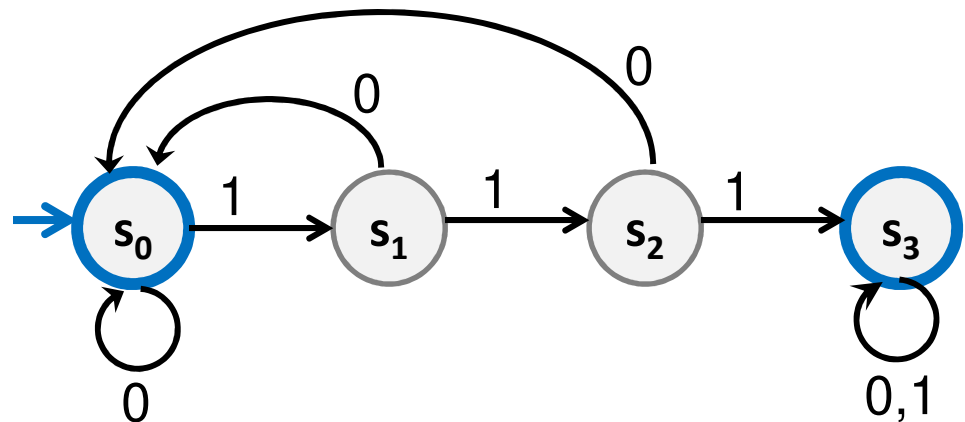


# Finite State Machines

(DFA)

- States
- Transitions on input symbols
- Start state and final states
- The “language recognized” by the machine is the set of strings that reach a final state from the start

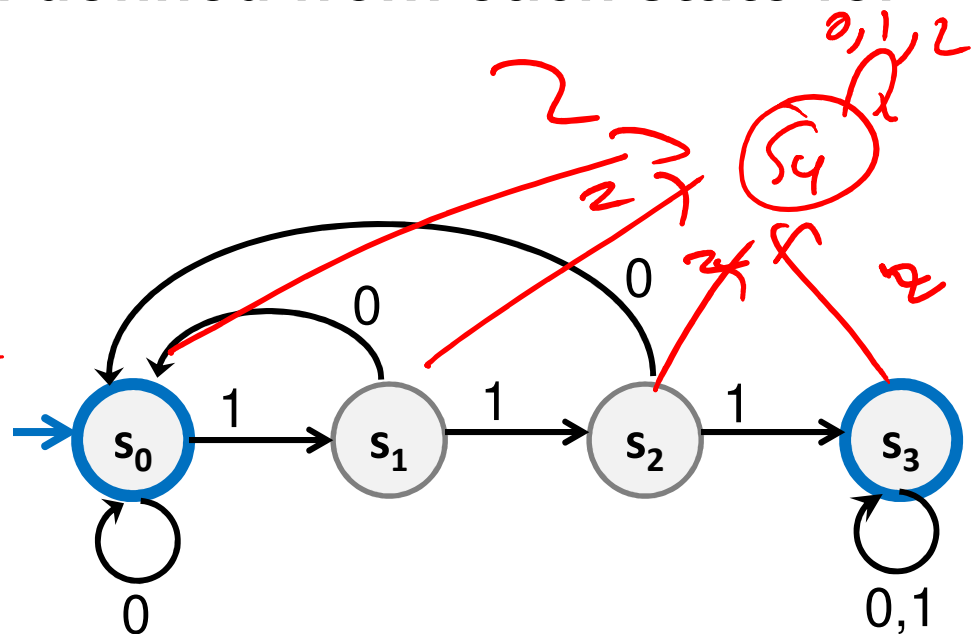
Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$



# Finite State Machines

- Each machine designed for strings over some fixed alphabet  $\Sigma$ .
- Must have a transition defined from each state for **every** symbol in  $\Sigma$ .

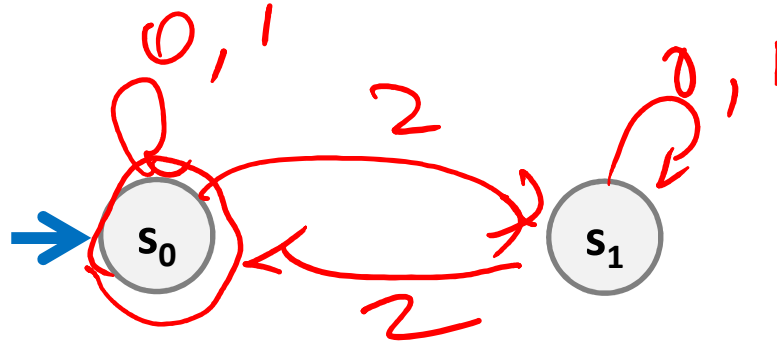
Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$



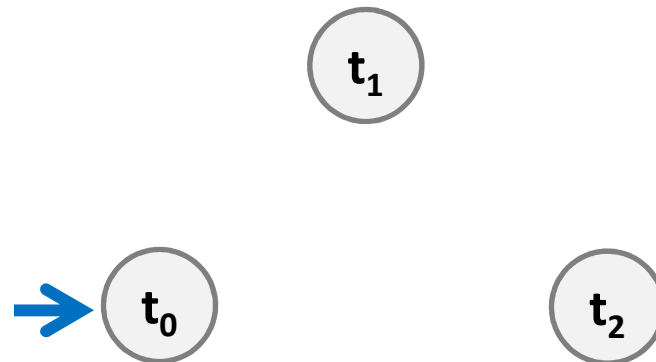
## Strings over $\{0, 1, 2\}$

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$M_1$ : Strings with an even number of 2's



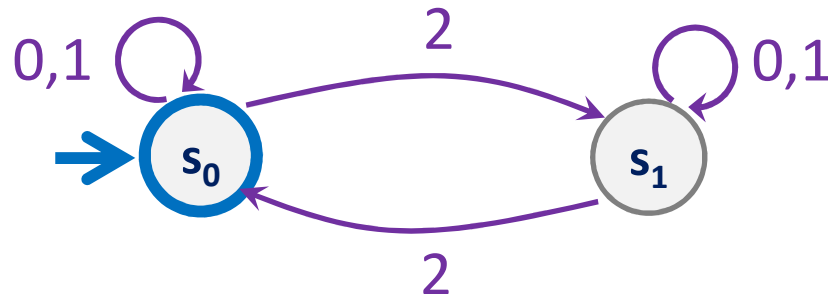
$M_2$ : Strings where the sum of digits mod 3 is 0



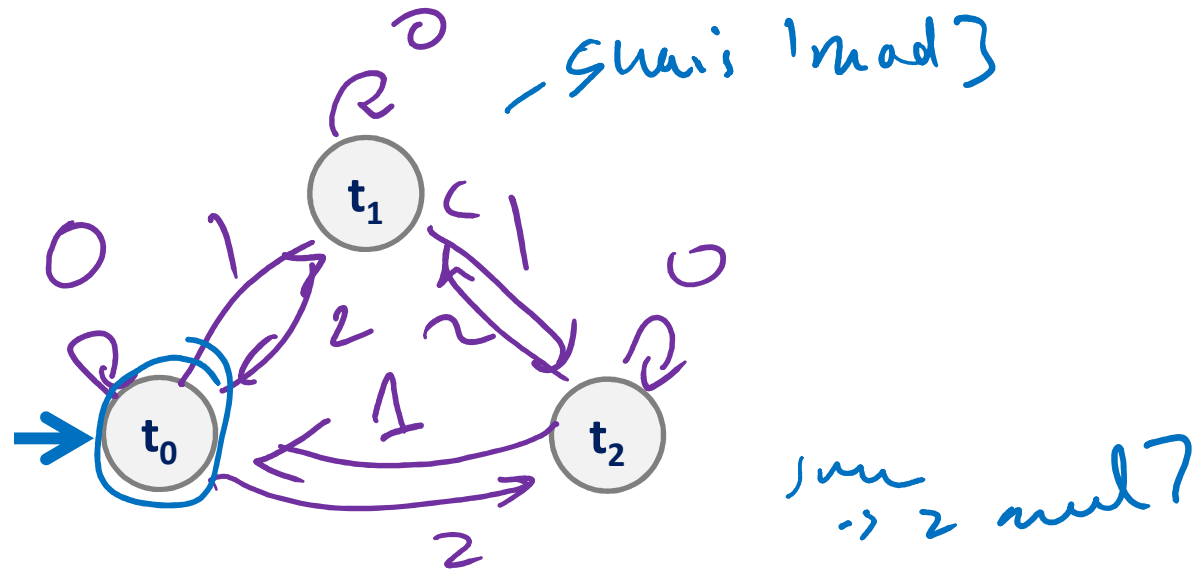
## Strings over $\{0, 1, 2\}$

---

$M_1$ : Strings with an even number of 2's



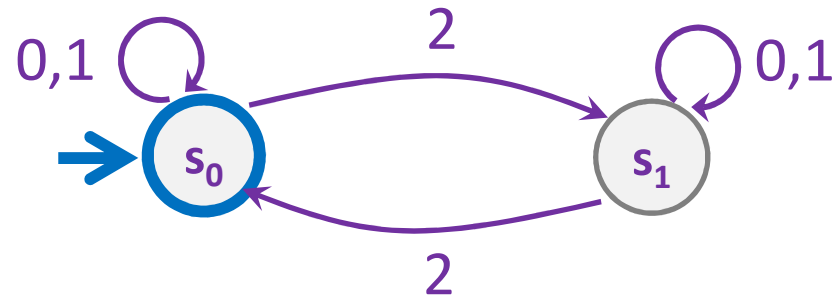
$M_2$ : Strings where the sum of digits mod 3 is 0



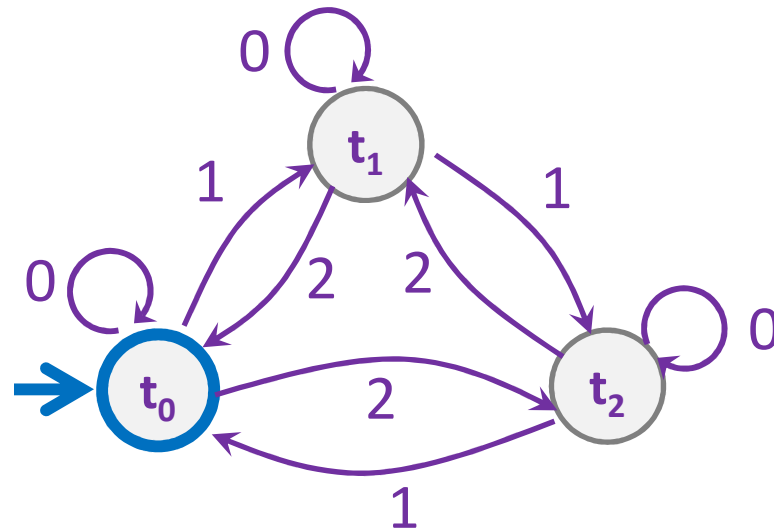
## Strings over $\{0, 1, 2\}$

---

$M_1$ : Strings with an even number of 2's

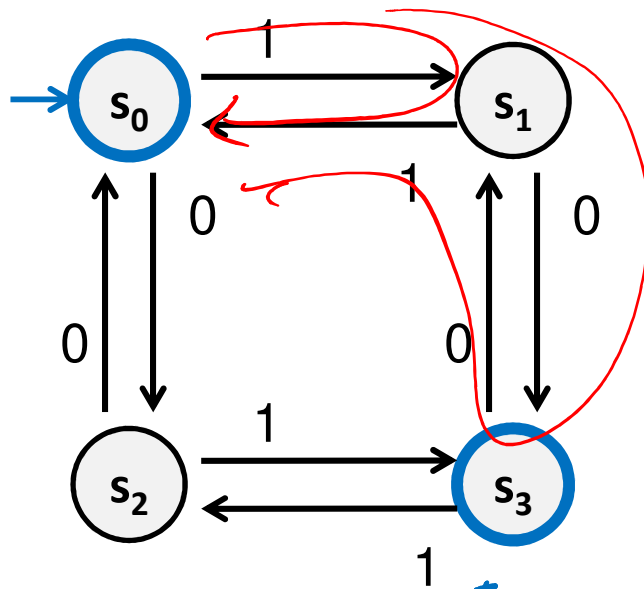


$M_2$ : Strings where the sum of digits mod 3 is 0



# What language does this machine recognize?

$s_0 = \text{even \# of 1's and even \# of 0's}$



$s_3 = \text{odd \# of 1's and odd \# of 0's}$

parity of  
#1's = parity of # of 0's

~~Alternates 1's and 0's  
all 1's for all 0's~~

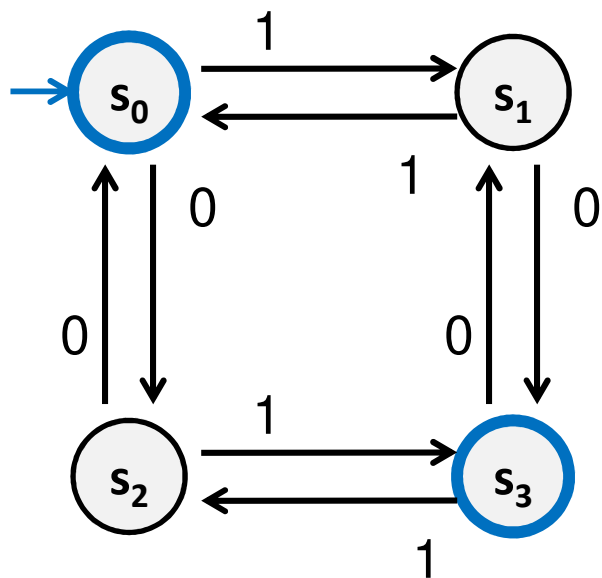
~~String 1<sup>st</sup> char is different  
from last~~

strings with even # of  
chars ✓

~~equal #1's and 0's~~

# What language does this machine recognize?

---



The set of all binary strings with  $\# \text{ of } 1\text{'s} \equiv \# \text{ of } 0\text{'s} \pmod{2}$   
(both are even or both are odd).

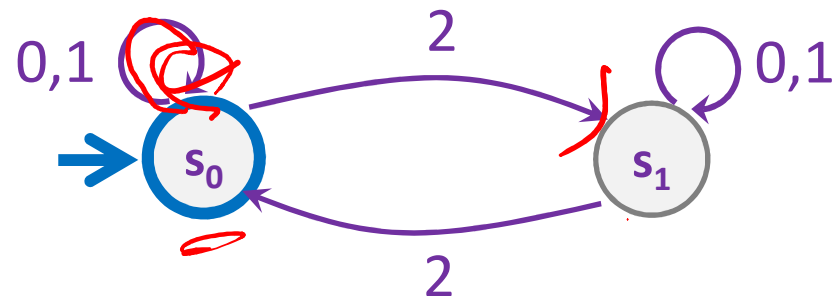
Can you think of a simpler description?



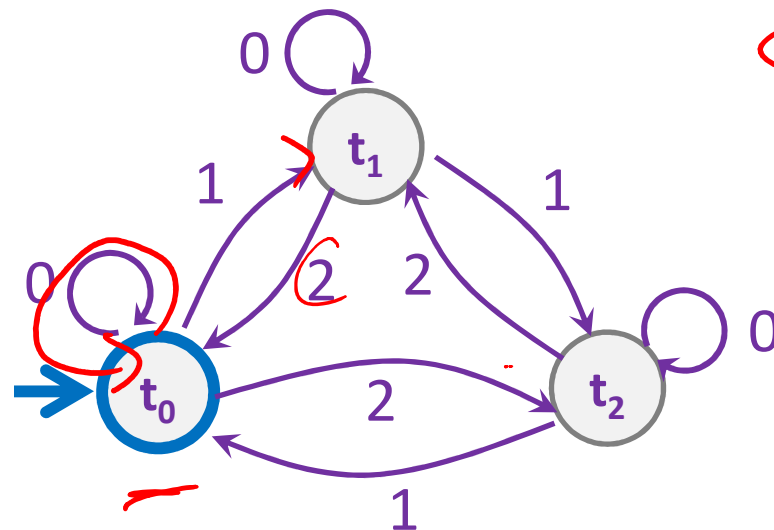
# Strings over $\{0, 1, 2\}$

---

$M_1$ : Strings with an even number of 2's



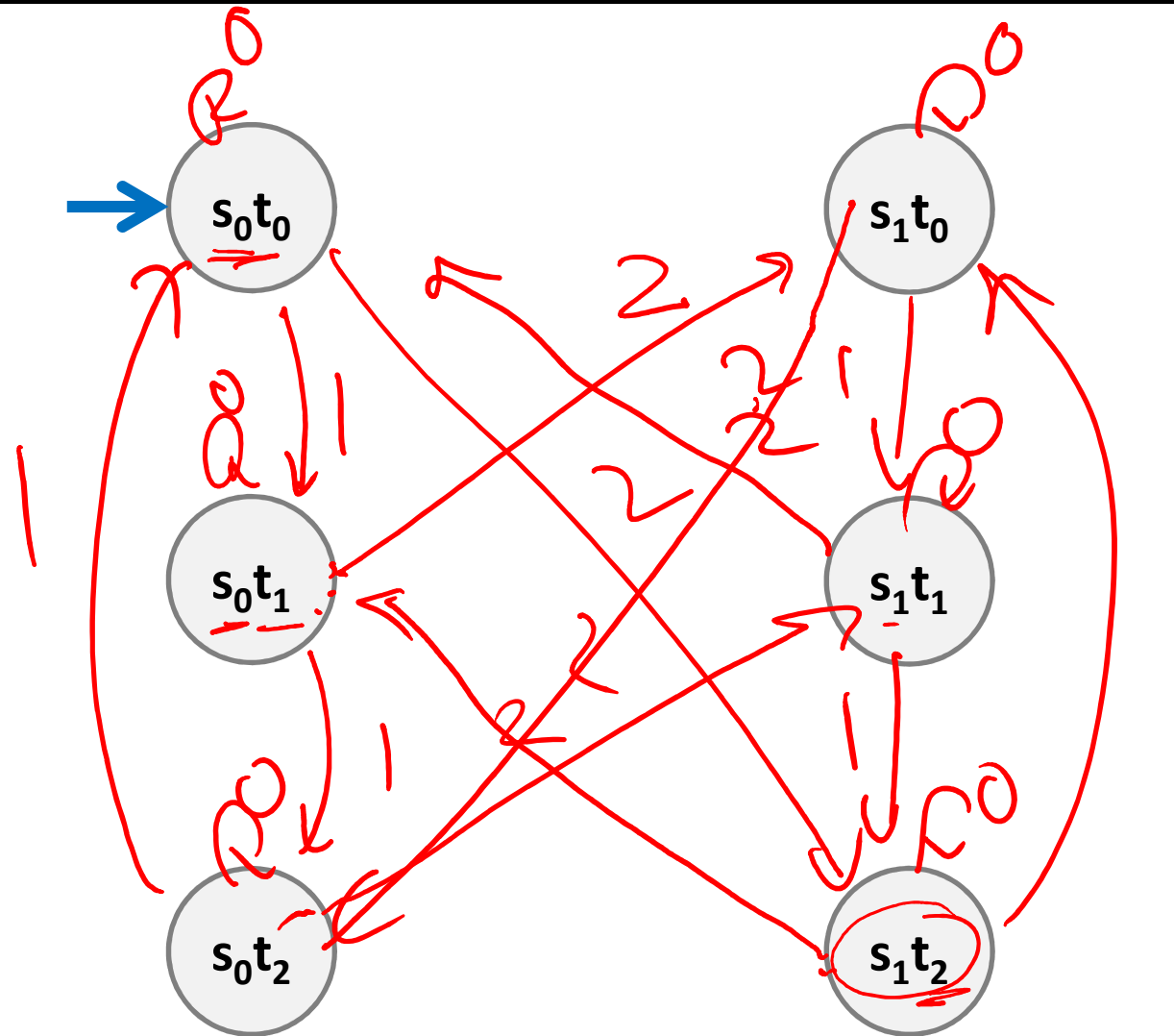
$M_2$ : Strings where the sum of digits mod 3 is 0



Handwritten red notes:  $012$  and  $222$  with underlines, circled together.

Strings over  $\{0,1,2\}$  w/ even number of 2's and mod 3 sum 0

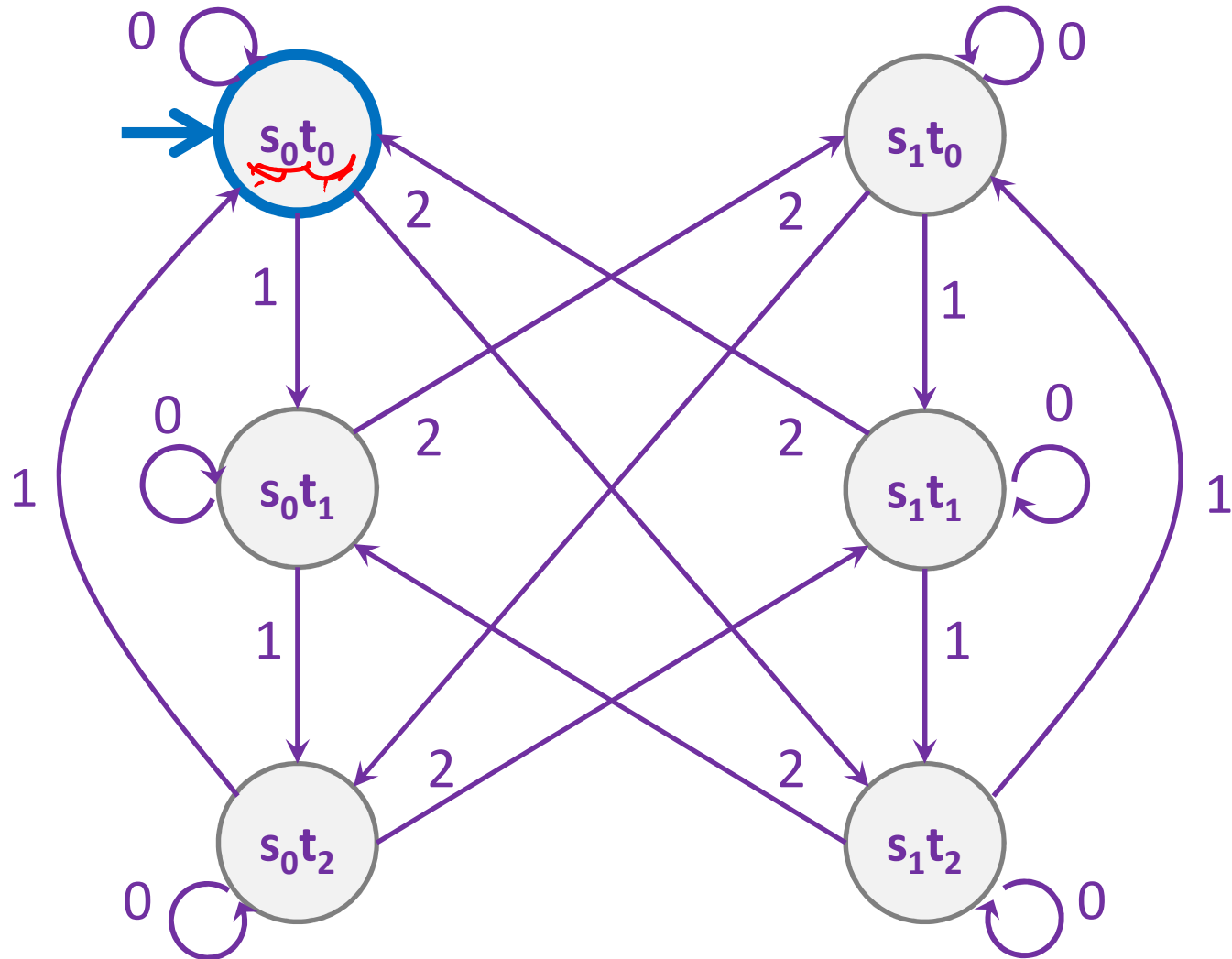
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Strings over  $\{0,1,2\}$  w/ even number of 2's and mod 3 sum 0

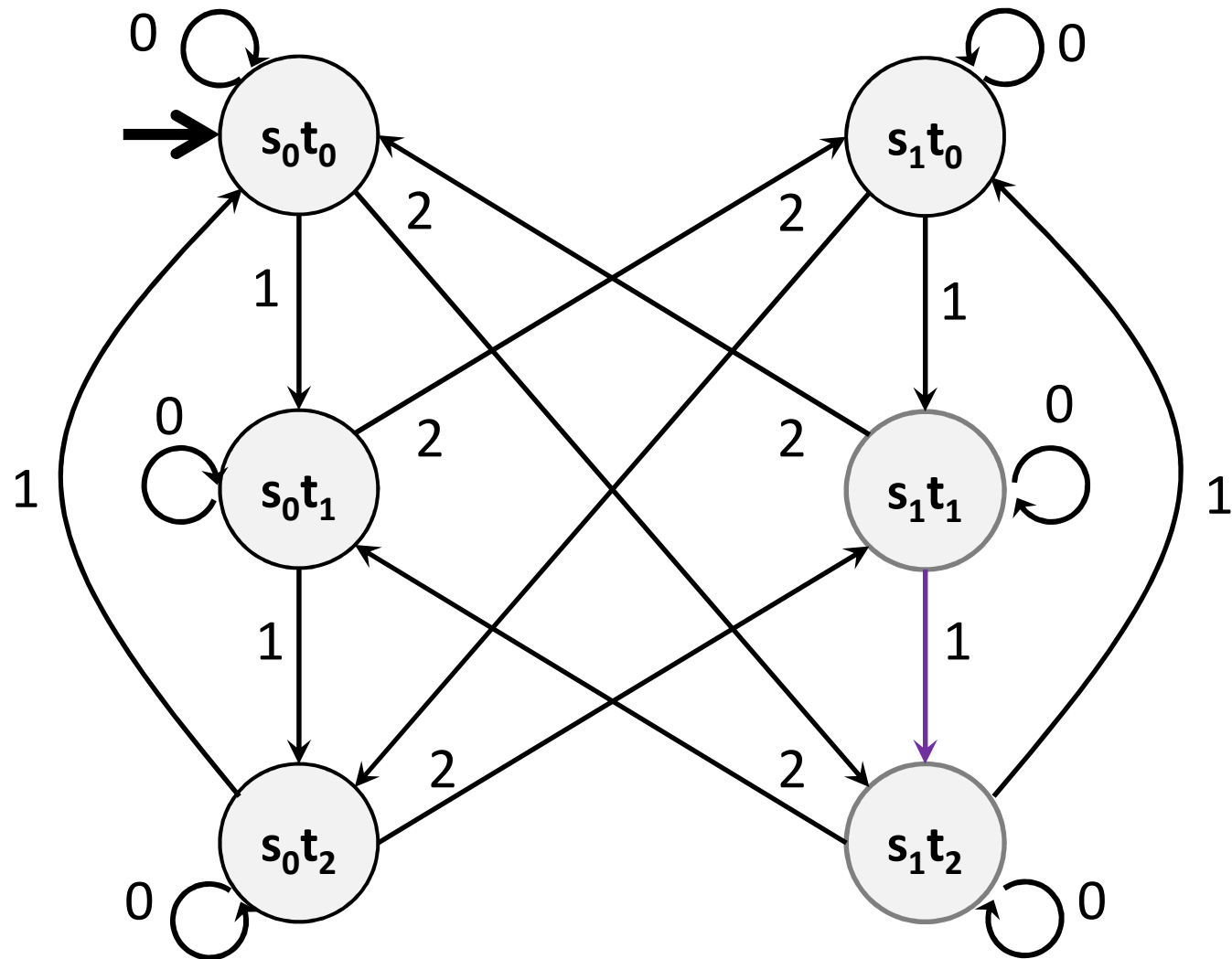
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or



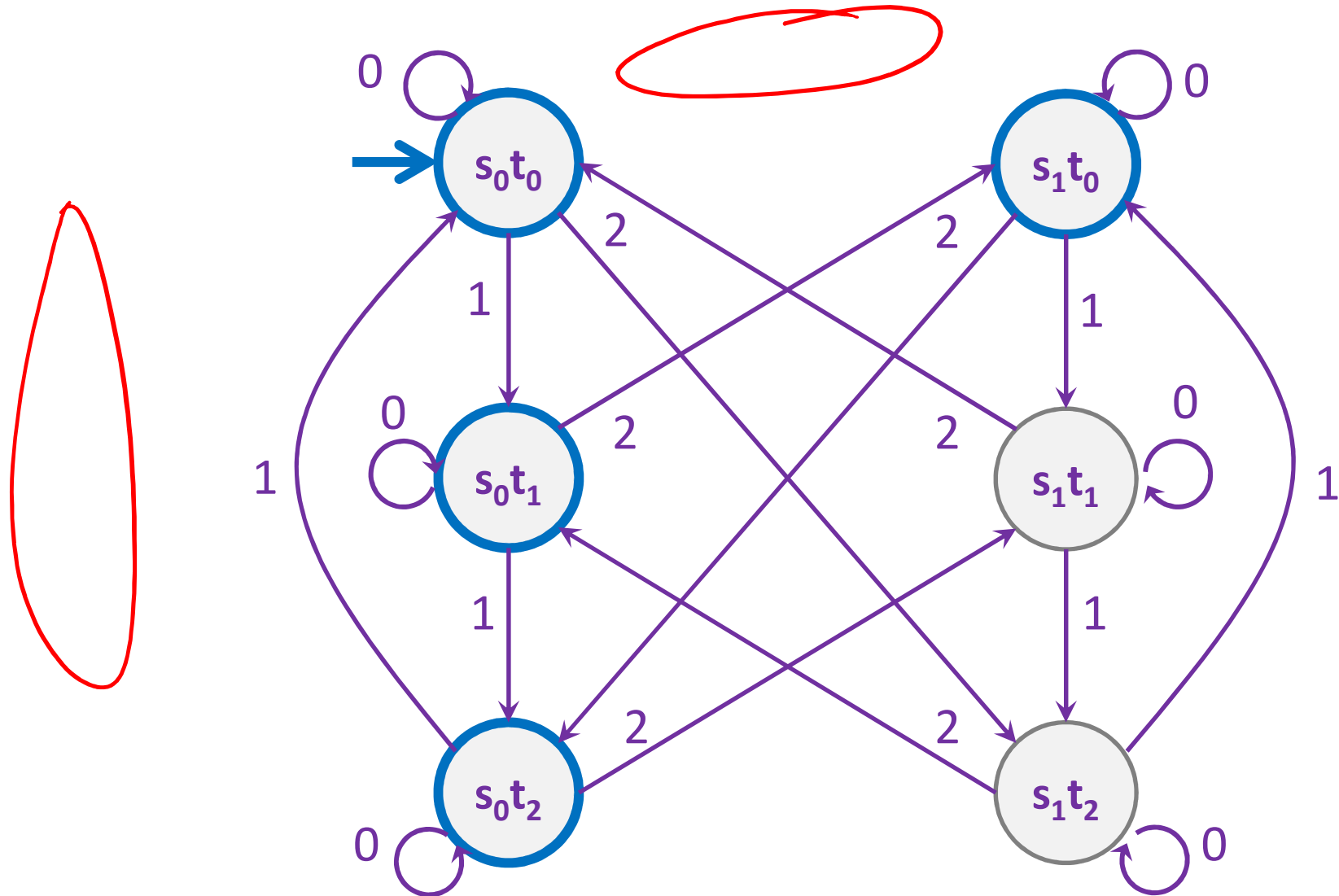
Strings over  $\{0,1,2\}$  w/ even number of 2's OR mod 3 sum 0?

---



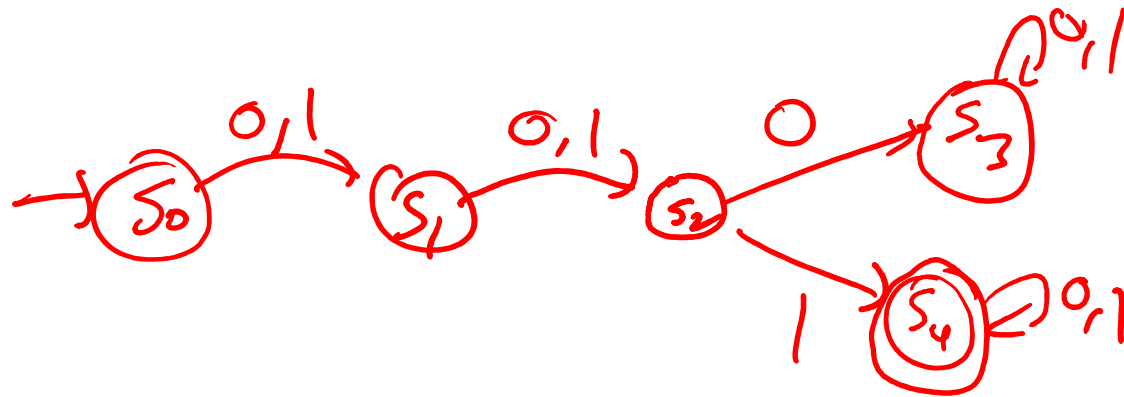
Strings over  $\{0,1,2\}$  w/ even number of 2's OR mod 3 sum 0

---



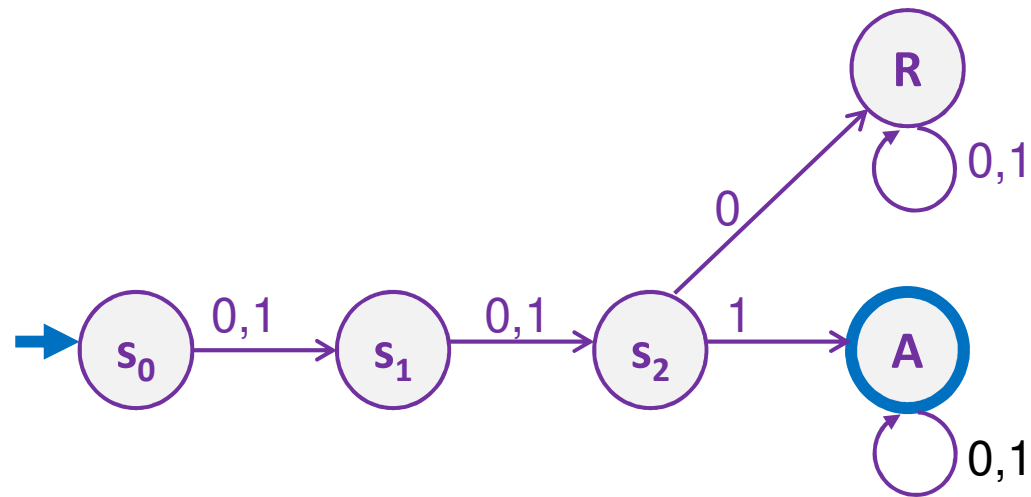
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the start

---



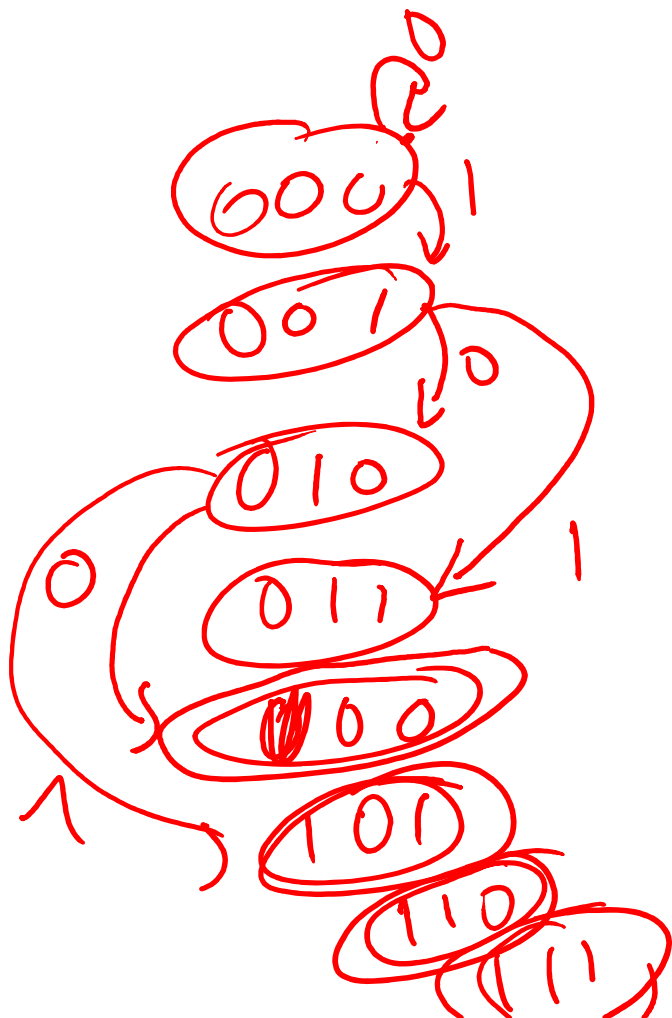
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the start

---



The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

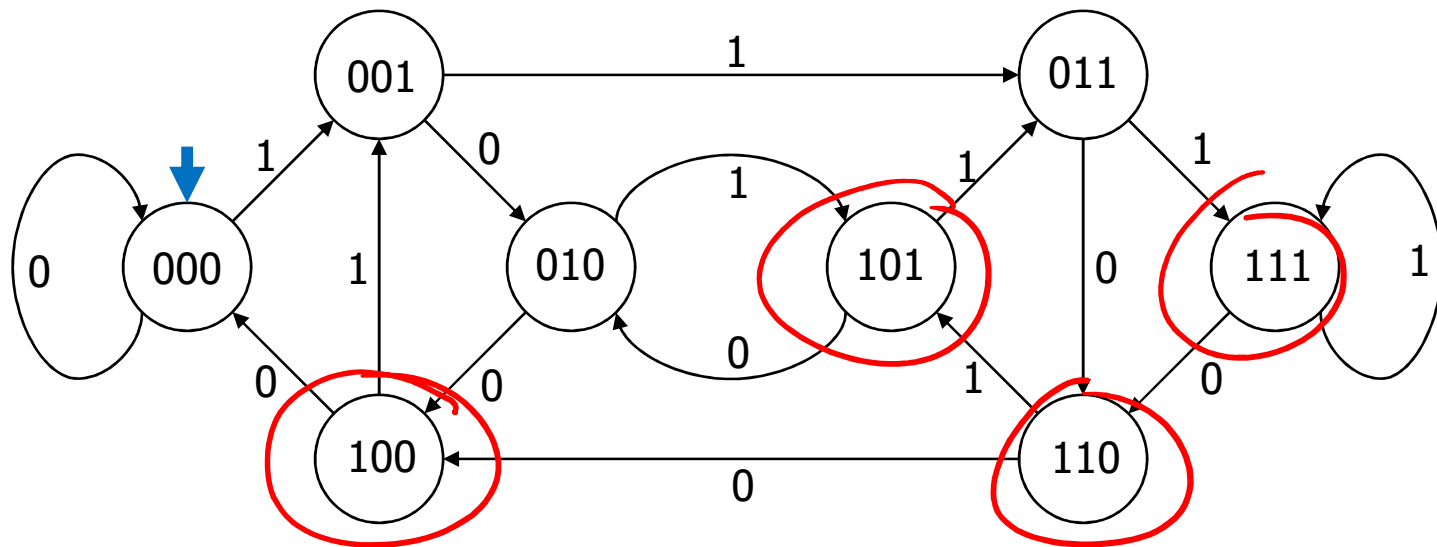
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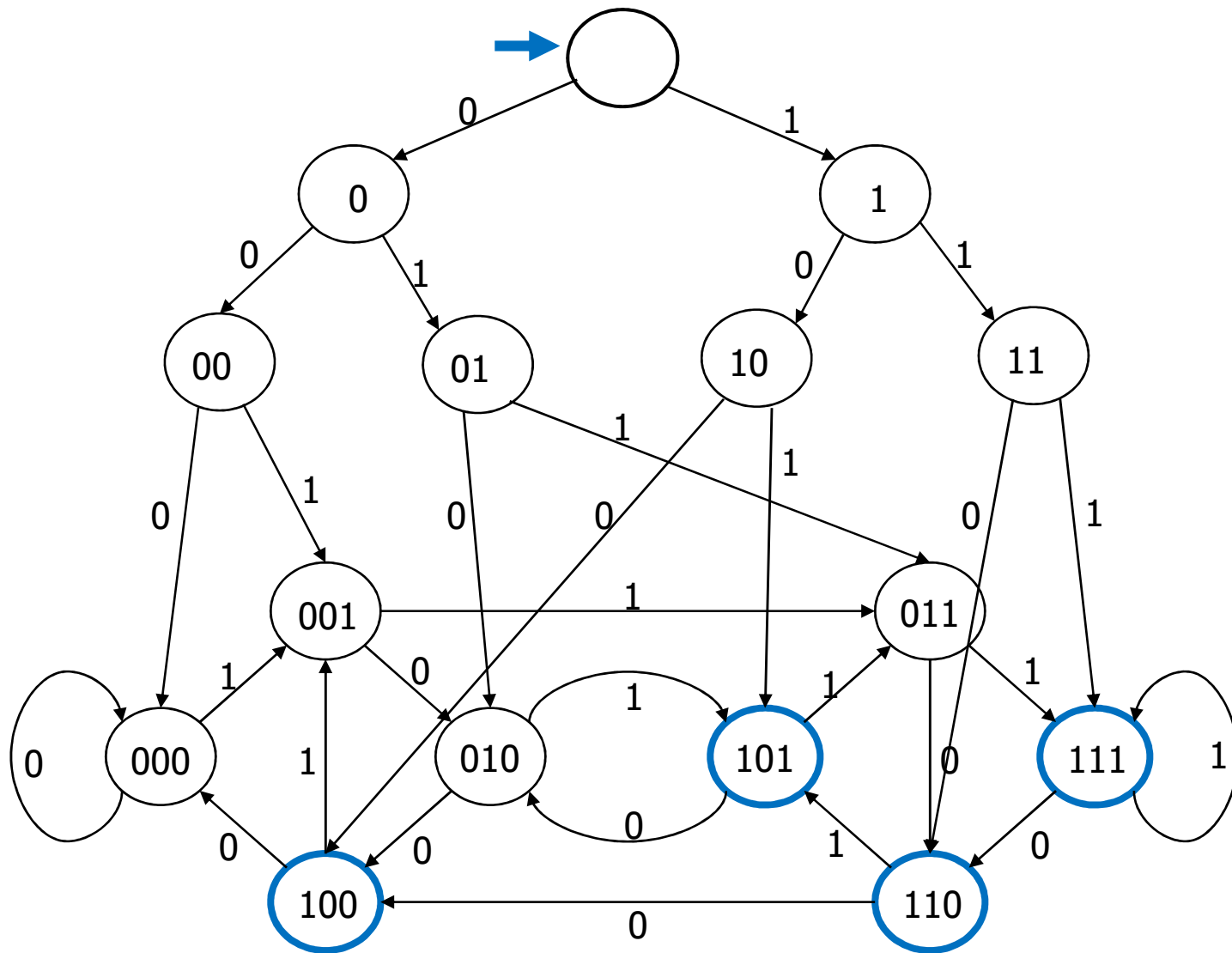
## 3 bit shift register “Remember the last three bits”

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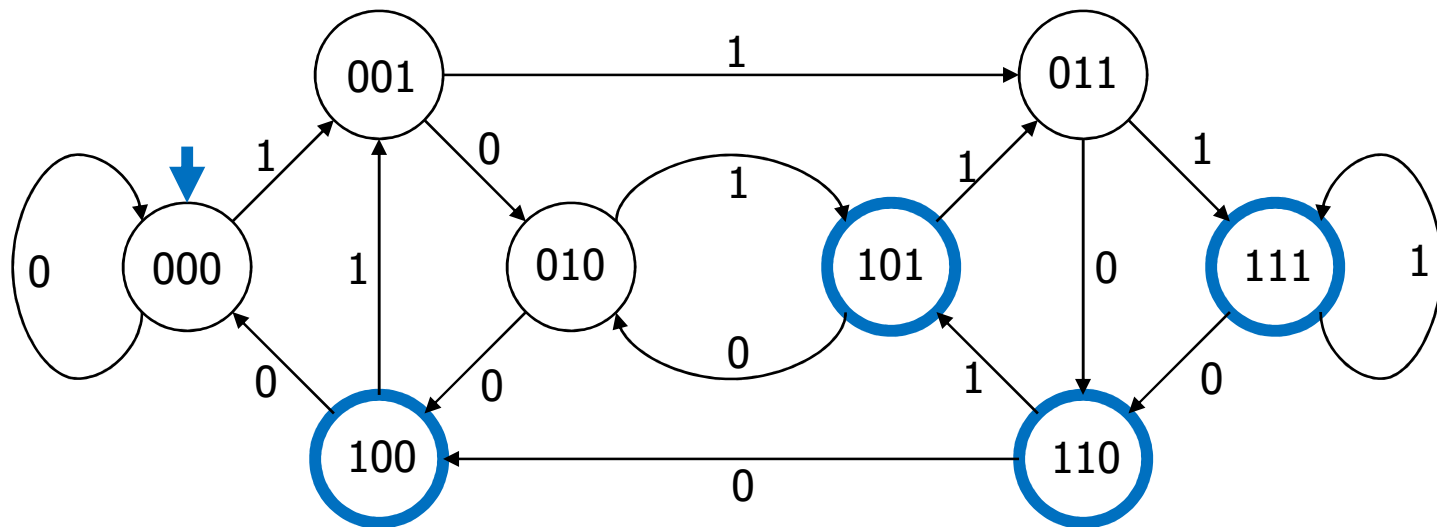
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

---



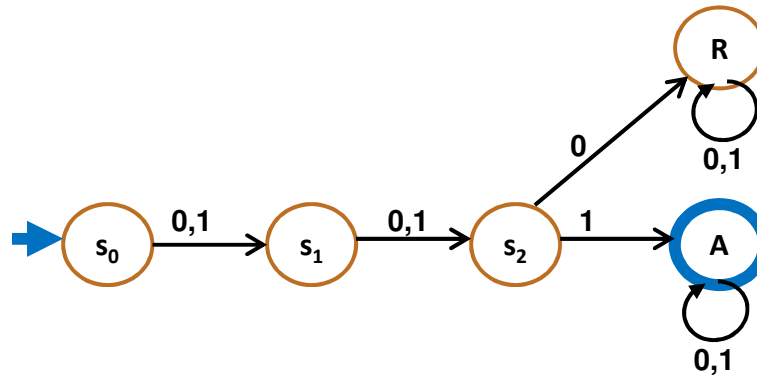
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

---

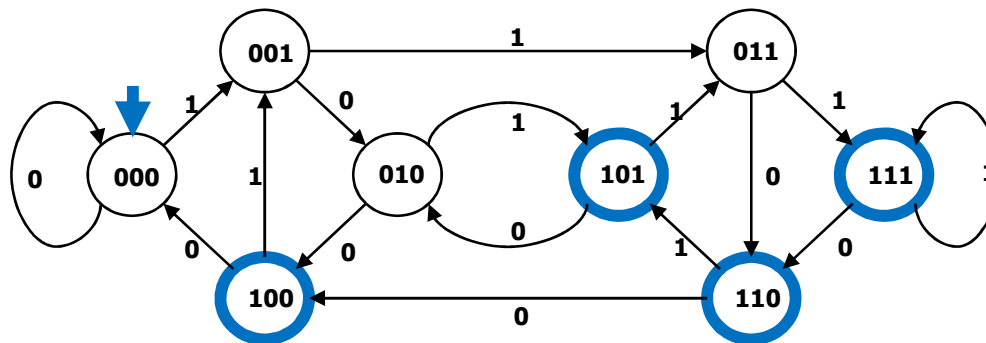


# The beginning versus the end

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*n+1 possible*  
*n+2 states*



*2<sup>n</sup>*

# Adding Output to Finite State Machines

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- So far we have considered finite state machines that just accept/reject strings
  - called “Deterministic Finite Automata” or DFAs
- Now we consider finite state machines that with output
  - These are the kinds used as controllers



# Vending Machine

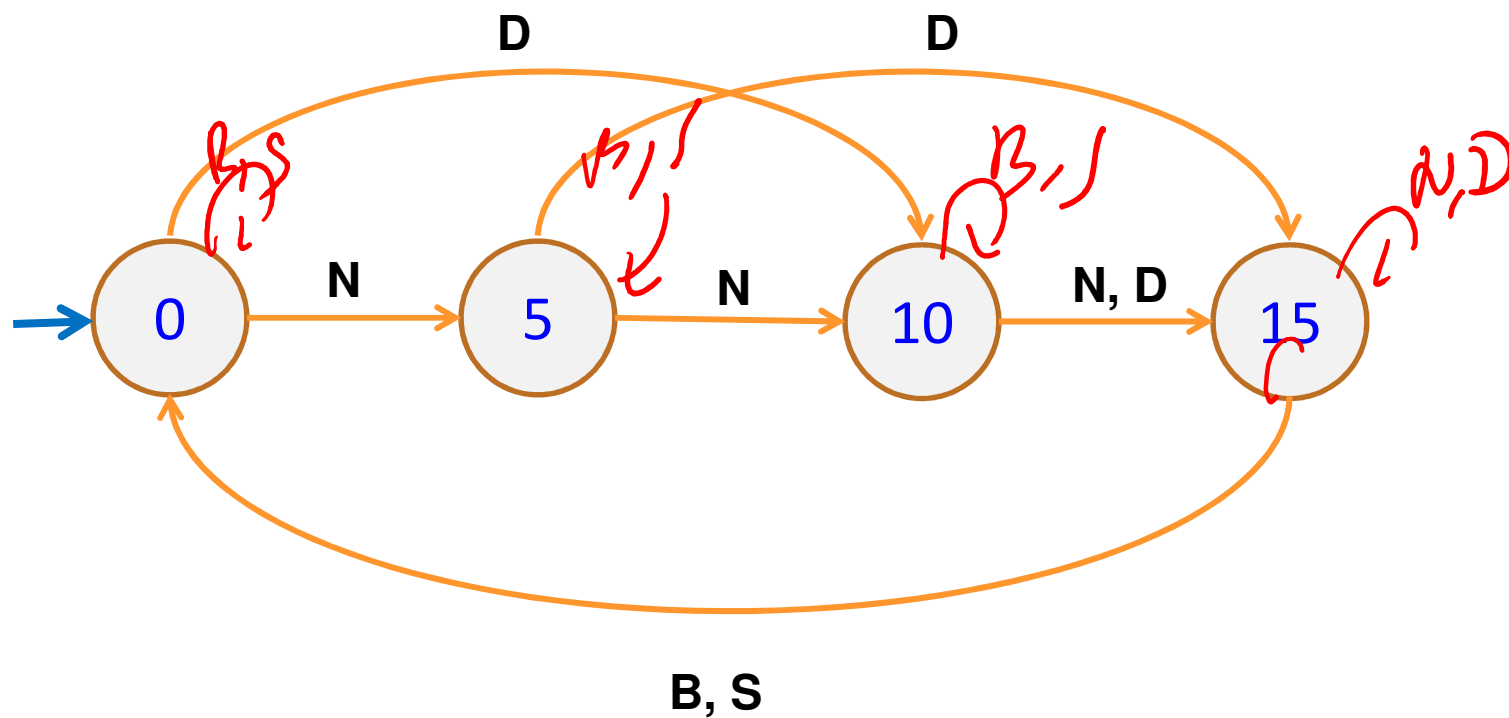


Enter 15 cents in dimes or nickels  
Press S or B for a candy bar



# Vending Machine, v0.1

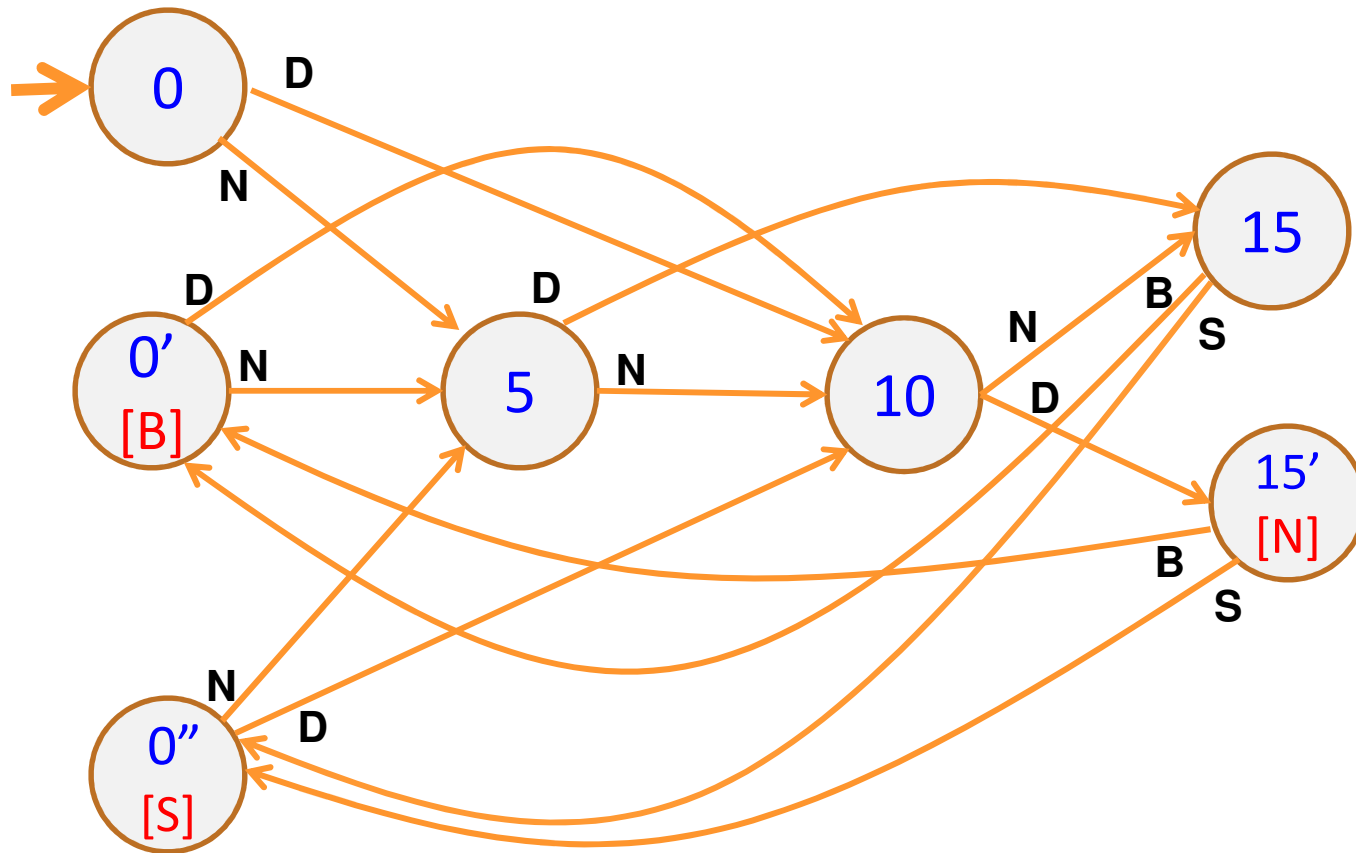
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Basic transitions on **N** (nickel), **D** (dime), **B** (butterfinger), **S** (snickers)

# Vending Machine, v0.2

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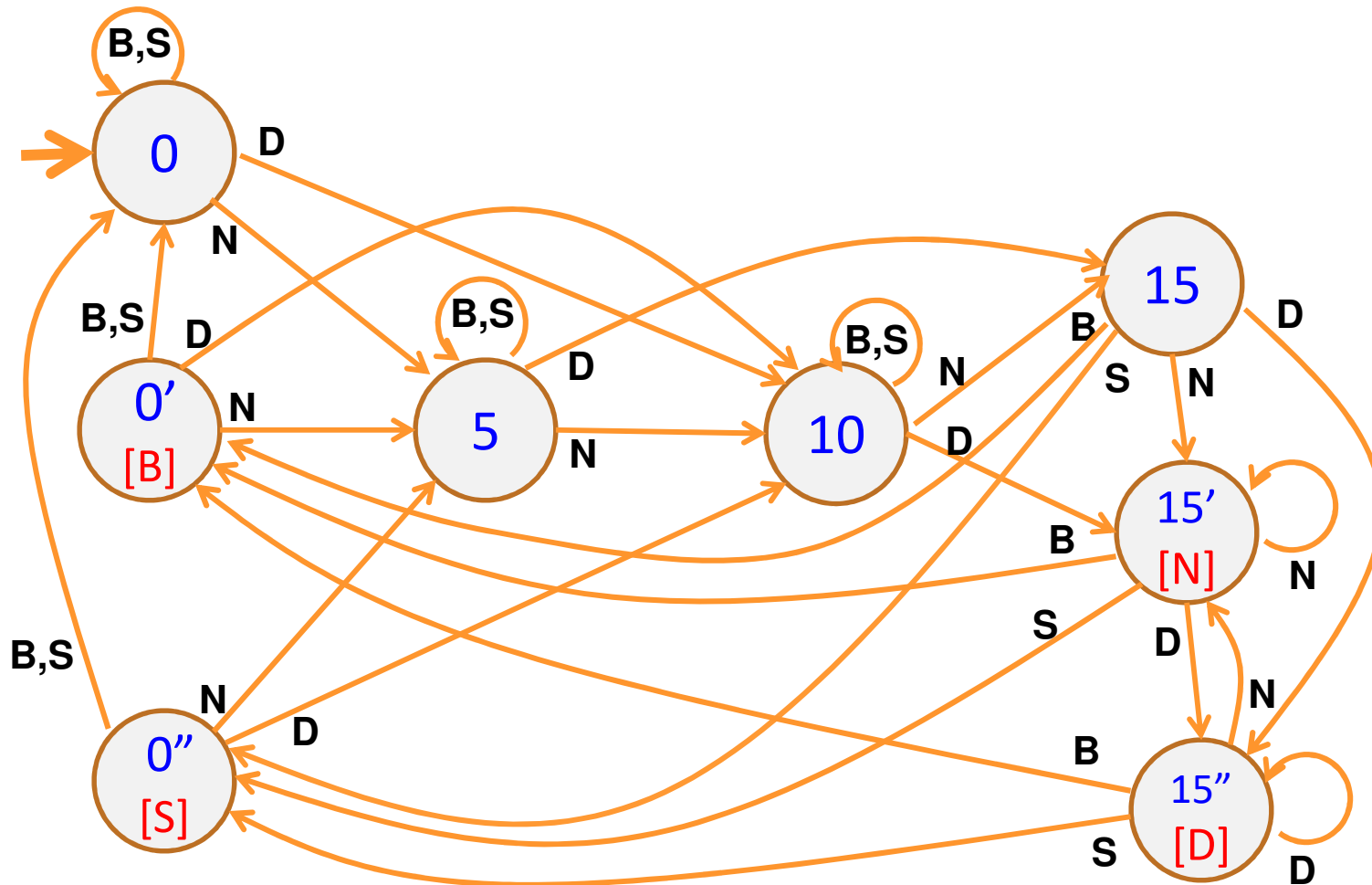


Adding output to states: **N** – Nickel, **S** – Snickers, **B** – Butterfinger



# Vending Machine, v1.0

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Adding additional “unexpected” transitions to cover all symbols for each state