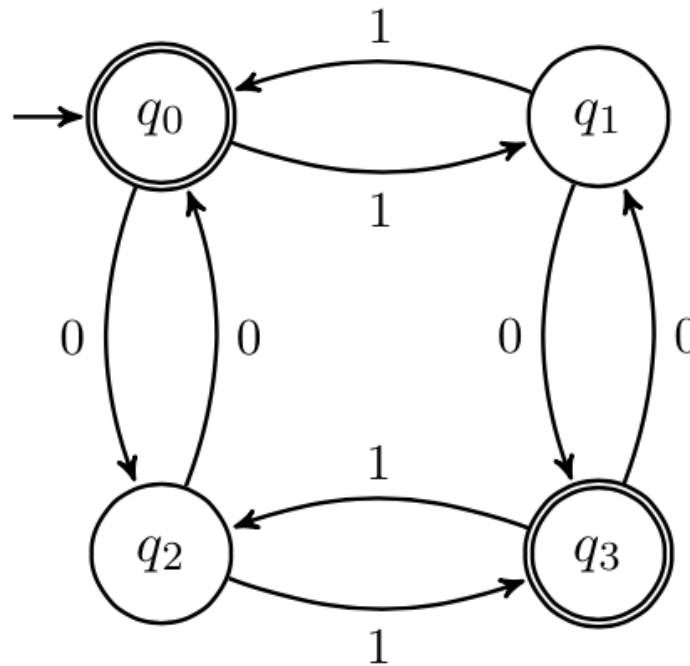


CSE 311: Foundations of Computing

Lecture 20: Directed Graphs, Closures, & Finite State Machines



Last Class: Relations & Composition

Let A and B be sets,

A **binary relation from A to B** is a subset of $A \times B$

Let A be a set,

A **binary relation on A** is a subset of $A \times A$

The **composition** of relation R and S , $S \circ R$ is the relation defined by:

$$S \circ R = \{ (a, c) \mid \exists b \text{ such that } (a,b) \in R \text{ and } (b,c) \in S \}$$

Last Class: Powers of a Relation

$R^0 = \{(a, a) \mid a \in A\}$ “the equality relation on A ”

$R^{n+1} = R^n \circ R$ for $n \geq 0$

Last class: Matrix Representation

Relation R on $A = \{a_1, \dots, a_n\}$

$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, a_j) \in R \\ 0 & \text{if } (a_i, a_j) \notin R \end{cases}$$

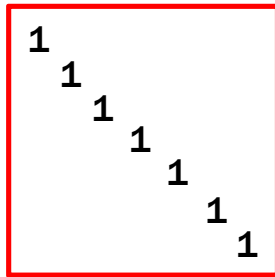
$\{(1, 1), (1, 2), (1, 4), (2, 1), (2, 3), (3, 2), (3, 3), (4, 2), (4, 3)\}$

	1	2	3	4
1	1	1	0	1
2	1	0	1	0
3	0	1	1	0
4	0	1	1	0

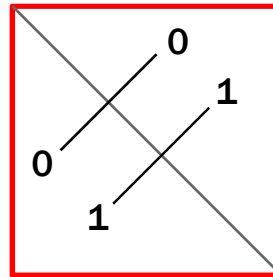
Last class: Matrix Representation

Relation R on $A = \{a_1, \dots, a_n\}$

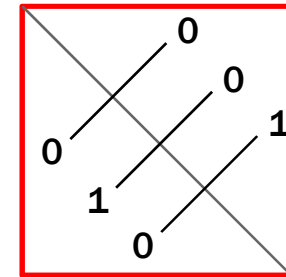
$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, a_j) \in R \\ 0 & \text{if } (a_i, a_j) \notin R \end{cases}$$



reflexive



symmetric

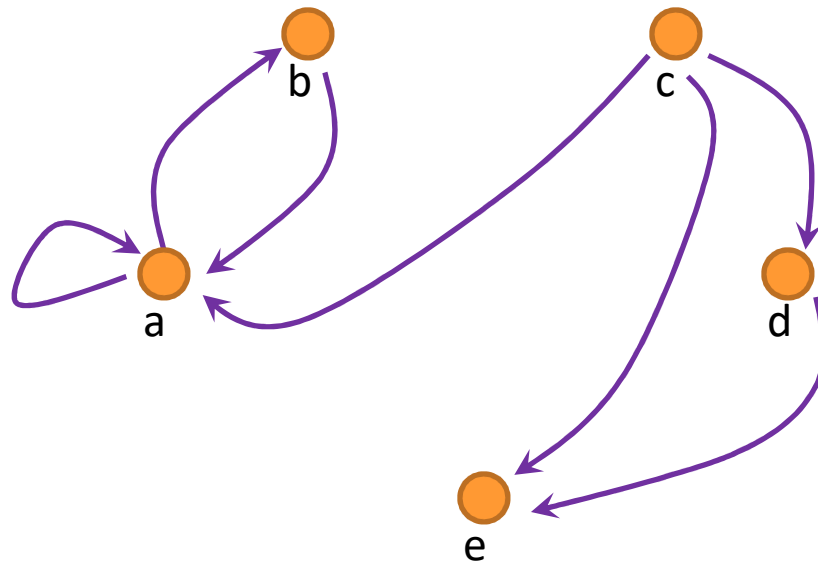


antisymmetric

Last Class: Representation of Relations

Directed Graph Representation (Digraph)

$\{(a, b), (a, a), (b, a), (c, a), (c, d), (c, e), (d, e)\}$



How Properties of Relations show up in Graphs

Let R be a relation on A .

R is **reflexive** iff $(a,a) \in R$ for every $a \in A$

R is **symmetric** iff $(a,b) \in R$ implies $(b,a) \in R$

R is **antisymmetric** iff $(a,b) \in R$ and $a \neq b$ implies $(b,a) \notin R$

R is **transitive** iff $(a,b) \in R$ and $(b,c) \in R$ implies $(a,c) \in R$

How Properties of Relations show up in Graphs

Let R be a relation on A .

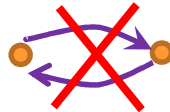
R is **reflexive** iff $(a,a) \in R$ for every $a \in A$

 at every node

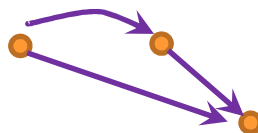
R is **symmetric** iff $(a,b) \in R$ implies $(b,a) \in R$

 or 

R is **antisymmetric** iff $(a,b) \in R$ and $a \neq b$ implies $(b,a) \notin R$



R is **transitive** iff $(a,b) \in R$ and $(b,c) \in R$ implies $(a,c) \in R$



Directed Graphs

$G = (V, E)$

V – vertices

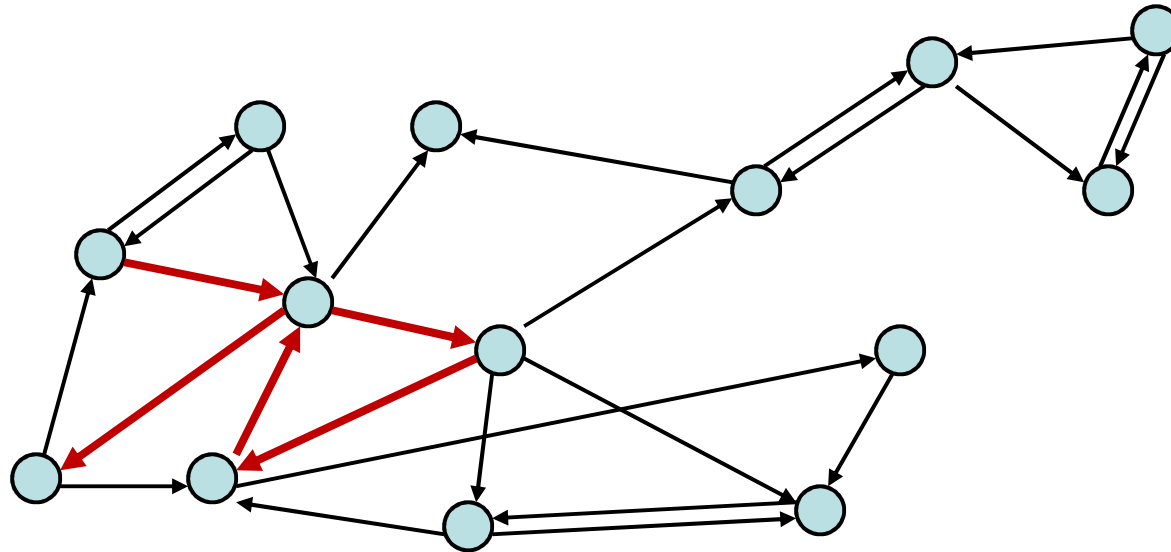
E – edges, ordered pairs of vertices

Path: v_0, v_1, \dots, v_k with each (v_i, v_{i+1}) in E

Simple Path: none of v_0, \dots, v_k repeated

Cycle: $v_0 = v_k$

Simple Cycle: $v_0 = v_k$, none of v_1, \dots, v_k repeated



Directed Graphs

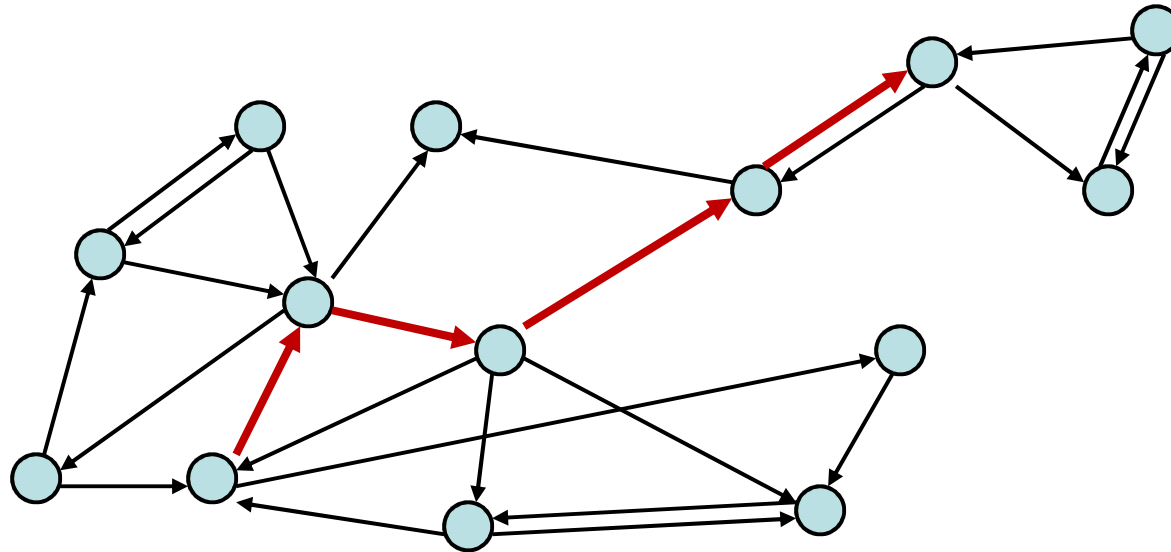
$G = (V, E)$ V – vertices
 E – edges, ordered pairs of vertices

Path: v_0, v_1, \dots, v_k with each (v_i, v_{i+1}) in E

Simple Path: none of v_0, \dots, v_k repeated

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Directed Graphs

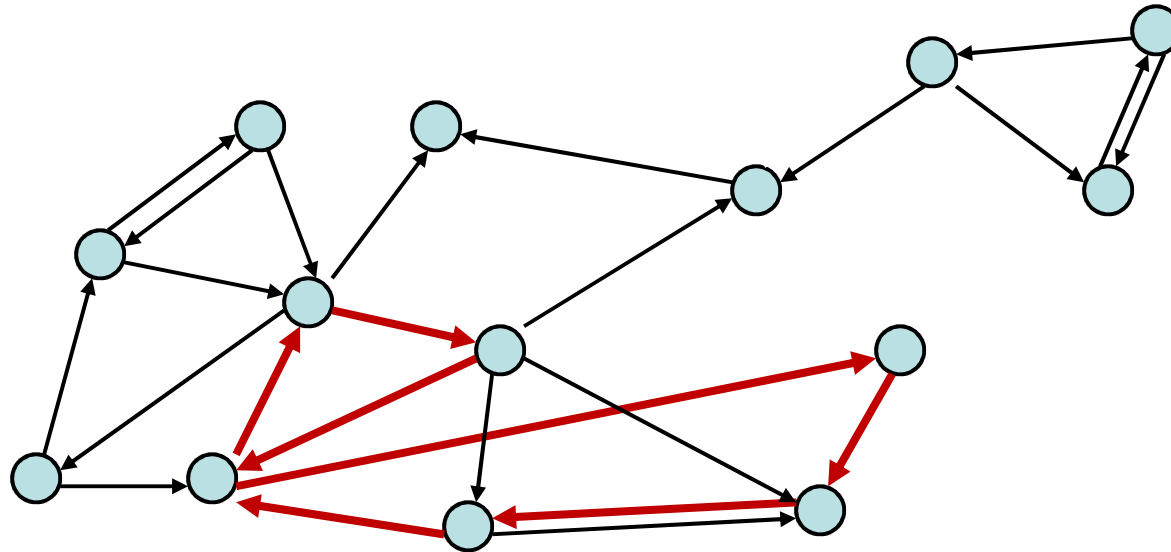
$G = (V, E)$ V – vertices
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Path: v_0, v_1, \dots, v_k with each (v_i, v_{i+1}) in E

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Directed Graphs

$G = (V, E)$

V – vertices

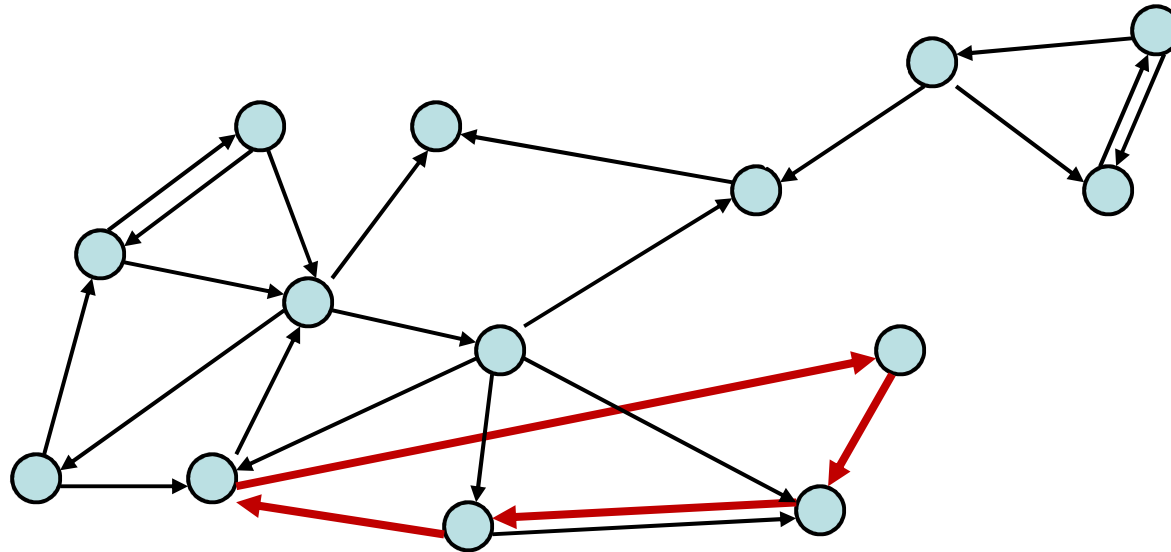
E – edges, ordered pairs of vertices

Path: v_0, v_1, \dots, v_k with each (v_i, v_{i+1}) in E

Simple Path: none of v_0, \dots, v_k repeated

Cycle: $v_0 = v_k$

Simple Cycle: $v_0 = v_k$, none of v_1, \dots, v_k repeated



Relational Composition using Digraphs

If $S = \{(2, 2), (2, 3), (3, 1)\}$ and $R = \{(1, 2), (2, 1), (1, 3)\}$

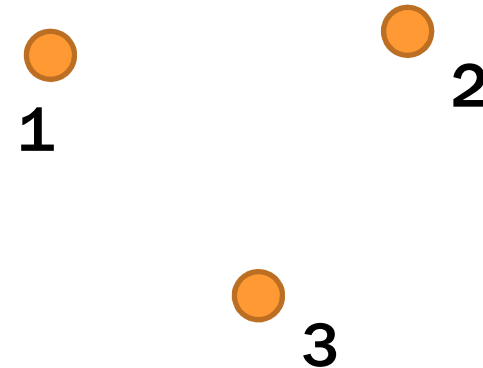
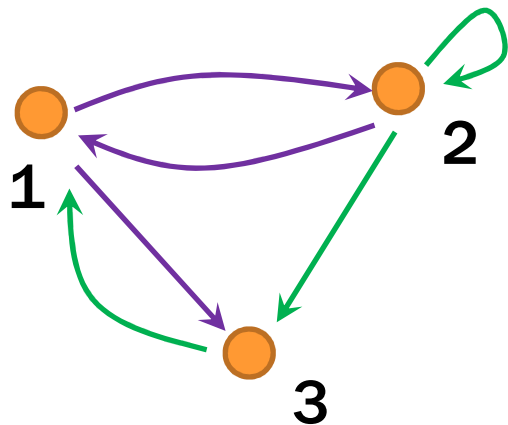
Compute $S \circ R$



Relational Composition using Digraphs

If $S = \{(2, 2), (2, 3), (3, 1)\}$ and $R = \{(1, 2), (2, 1), (1, 3)\}$

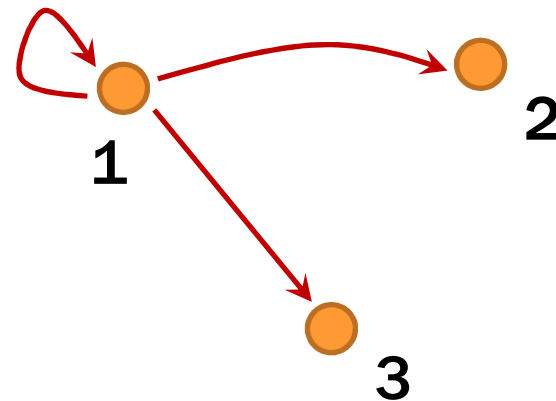
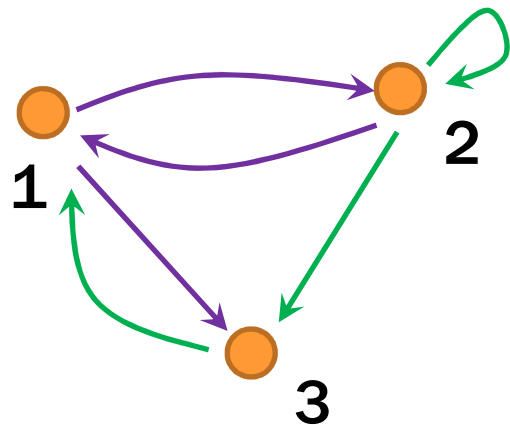
Compute $S \circ R$



Relational Composition using Digraphs

If $S = \{(2, 2), (2, 3), (3, 1)\}$ and $R = \{(1, 2), (2, 1), (1, 3)\}$

Compute $S \circ R$



Paths in Relations and Graphs

Defn: The **length** of a path in a graph is the number of edges in it (counting repetitions if edge used $>$ once).

Let R be a relation on a set A . There is a path of length n from a to b if and only if $(a,b) \in R^n$

Connectivity In Graphs

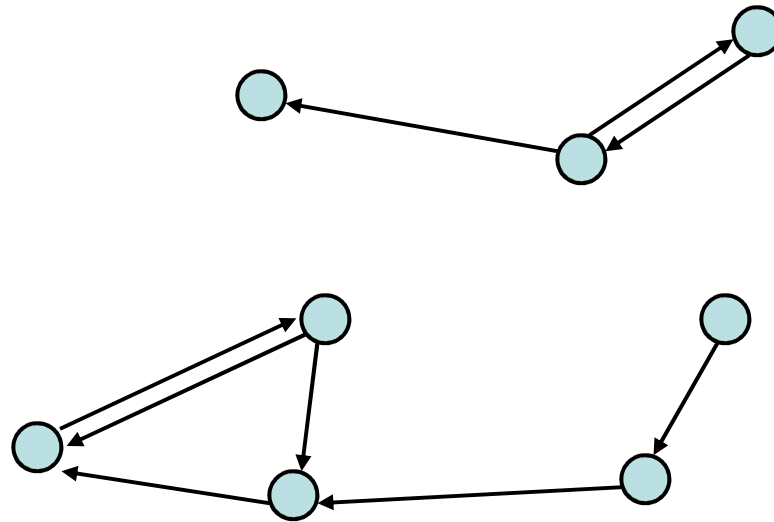
Defn: Two vertices in a graph are **connected** iff there is a path between them.

Let R be a relation on a set A . The **connectivity** relation R^* consists of the pairs (a,b) such that there is a path from a to b in R .

$$R^* = \bigcup_{k=0}^{\infty} R^k$$

Note: The text uses the wrong definition of this quantity. What the text defines (ignoring $k=0$) is usually called R^+

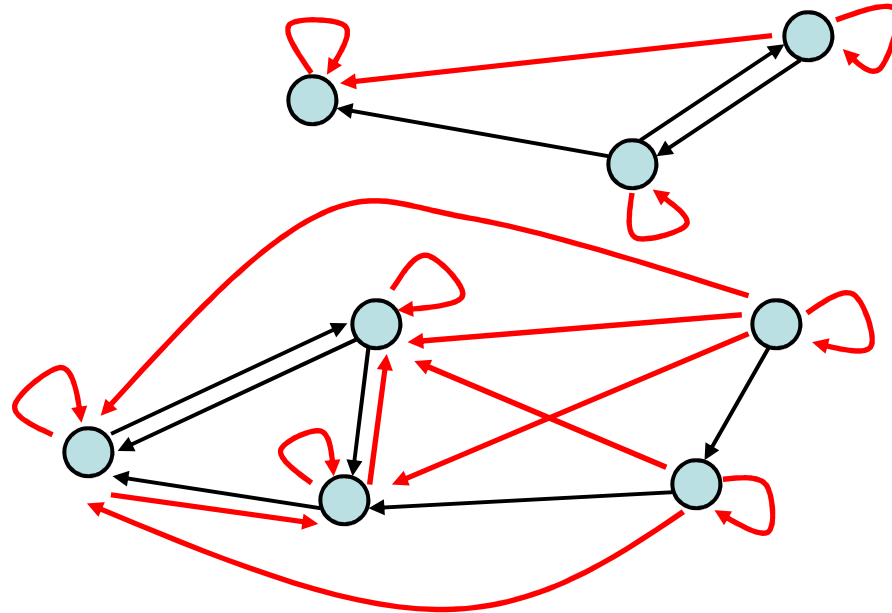
Transitive-Reflexive Closure



Add the **minimum possible** number of edges to make the relation transitive and reflexive.

The **transitive-reflexive closure** of a relation R is the connectivity relation R^*

Transitive-Reflexive Closure



Relation with the **minimum possible** number of **extra edges** to make the relation both transitive and reflexive.

The **transitive-reflexive closure** of a relation R is the connectivity relation R^*

n-ary Relations

Let A_1, A_2, \dots, A_n be sets. An ***n*-ary** relation on these sets is a subset of $A_1 \times A_2 \times \dots \times A_n$.

Relational Databases

STUDENT

Student_Name	ID_Number	Office	GPA
Knuth	328012098	022	4.00
Von Neuman	481080220	555	3.78
Russell	238082388	022	3.85
Einstein	238001920	022	2.11
Newton	1727017	333	3.61
Karp	348882811	022	3.98
Bernoulli	2921938	022	3.21

Relational Databases

STUDENT

Student_Name	ID_Number	Office	GPA	Course
Knuth	328012098	022	4.00	CSE311
Knuth	328012098	022	4.00	CSE351
Von Neuman	481080220	555	3.78	CSE311
Russell	238082388	022	3.85	CSE312
Russell	238082388	022	3.85	CSE344
Russell	238082388	022	3.85	CSE351
Newton	1727017	333	3.61	CSE312
Karp	348882811	022	3.98	CSE311
Karp	348882811	022	3.98	CSE312
Karp	348882811	022	3.98	CSE344
Karp	348882811	022	3.98	CSE351
Bernoulli	2921938	022	3.21	CSE351

What's not so nice?

Relational Databases

STUDENT

Student_Name	ID_Number	Office	GPA
Knuth	328012098	022	4.00
Von Neuman	481080220	555	3.78
Russell	238082388	022	3.85
Einstein	238001920	022	2.11
Newton	1727017	333	3.61
Karp	348882811	022	3.98
Bernoulli	2921938	022	3.21

TAKES

ID_Number	Course
328012098	CSE311
328012098	CSE351
481080220	CSE311
238082388	CSE312
238082388	CSE344
238082388	CSE351
1727017	CSE312
348882811	CSE311
348882811	CSE312
348882811	CSE344
348882811	CSE351
2921938	CSE351

Better

Database Operations: Projection

Find all offices: $\Pi_{\text{Office}}(\text{STUDENT})$

Office
022
555
333

Find offices and GPAs: $\Pi_{\text{Office,GPA}}(\text{STUDENT})$

Office	GPA
022	4.00
555	3.78
022	3.85
022	2.11
333	3.61
022	3.98
022	3.21

Database Operations: Selection

Find students with GPA > 3.9 : $\sigma_{\text{GPA}>3.9}(\text{STUDENT})$

Student_Name	ID_Number	Office	GPA
Knuth	328012098	022	4.00
Karp	348882811	022	3.98

Retrieve the name and GPA for students with GPA > 3.9:

$\Pi_{\text{Student_Name,GPA}}(\sigma_{\text{GPA}>3.9}(\text{STUDENT}))$

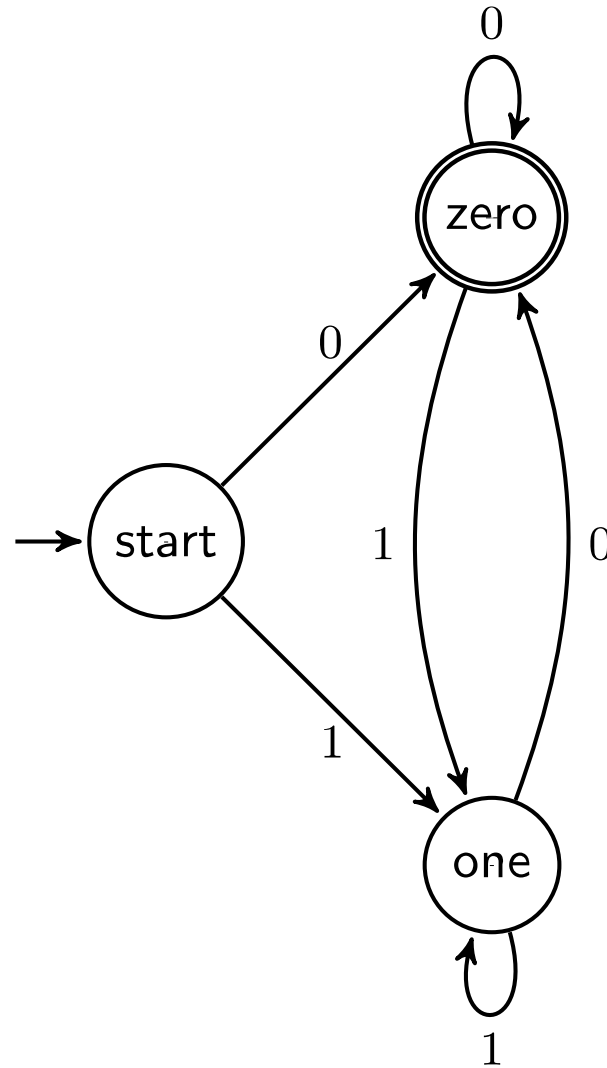
Student_Name	GPA
Knuth	4.00
Karp	3.98

Database Operations: Natural Join

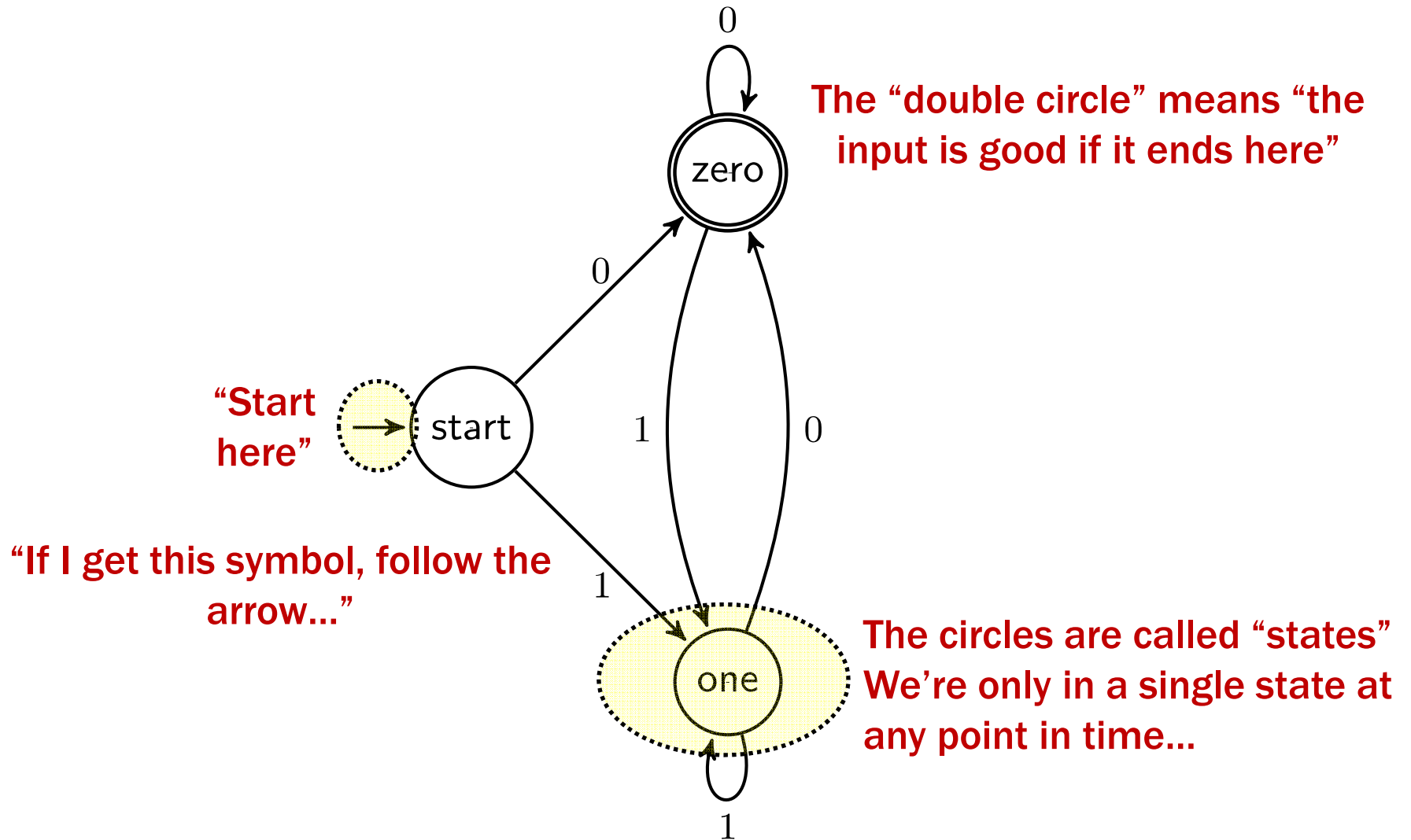
Student ⋈ Takes

Student_Name	ID_Number	Office	GPA	Course
Knuth	328012098	022	4.00	CSE311
Knuth	328012098	022	4.00	CSE351
Von Neuman	481080220	555	3.78	CSE311
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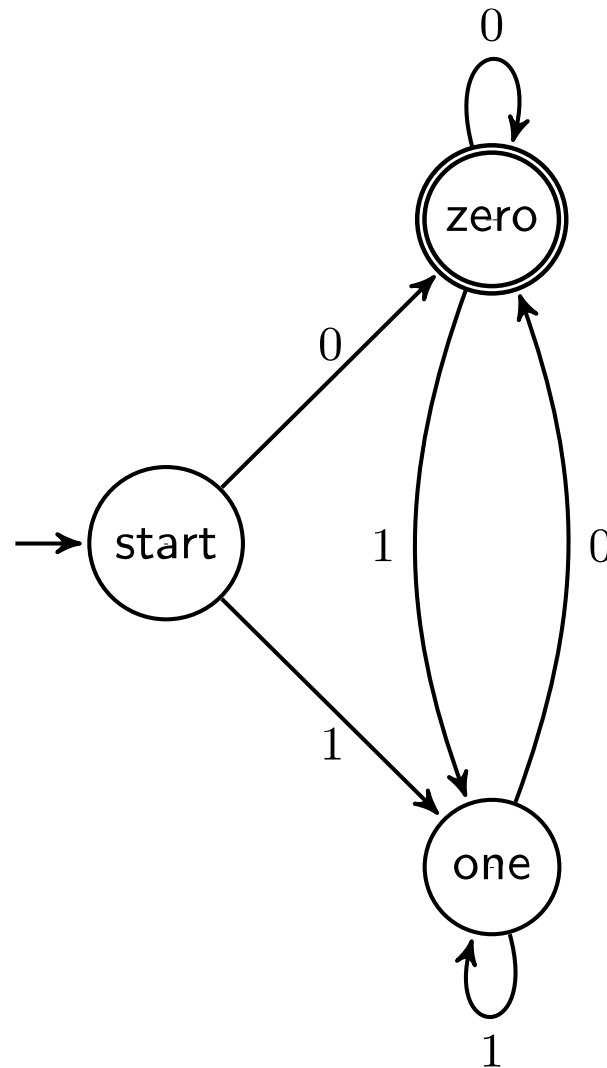
Selecting strings using labeled graphs as “machines”



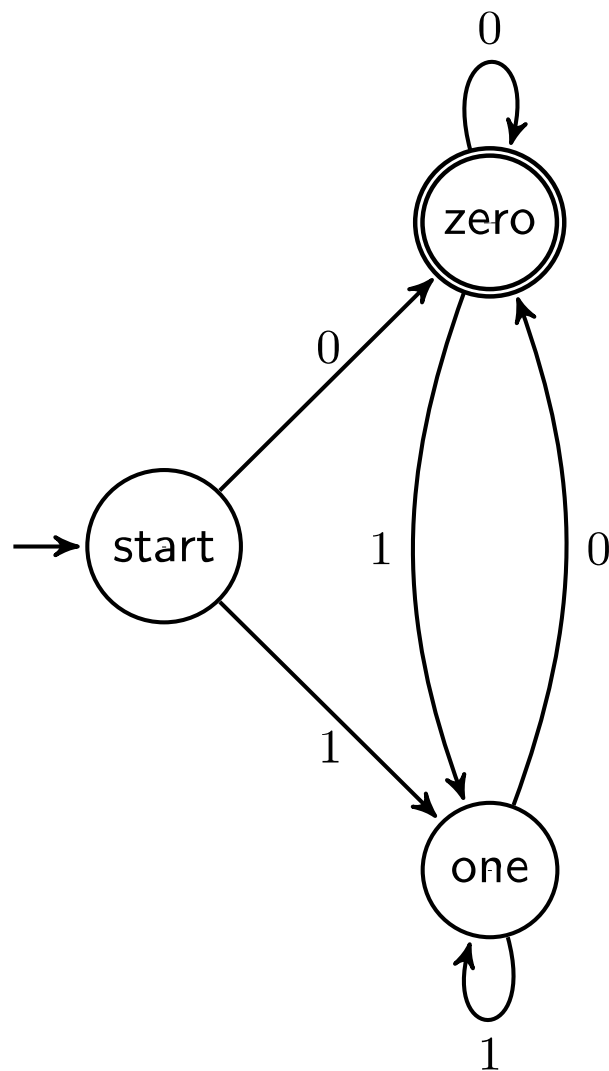
Finite State Machines



Which strings does this machine say are OK?



Which strings does this machine say are OK?

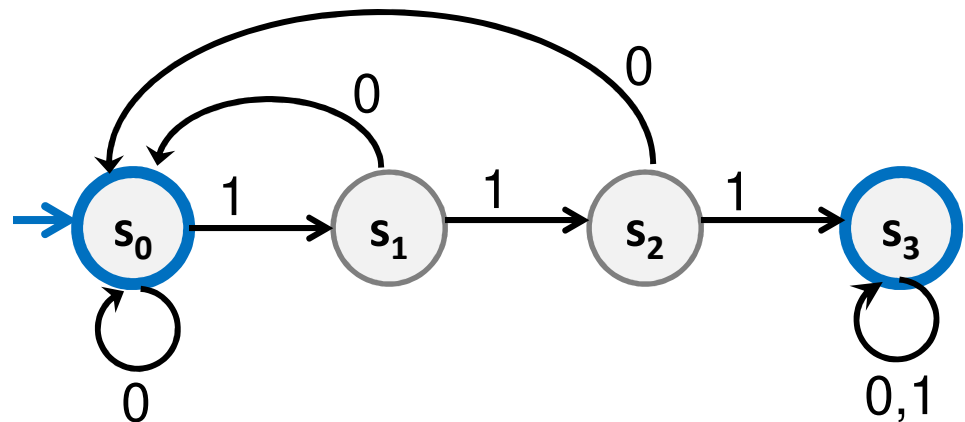


The set of all binary strings that end in 0

Finite State Machines

- States
- Transitions on input symbols
- Start state and final states
- The “language recognized” by the machine is the set of strings that reach a final state from the start

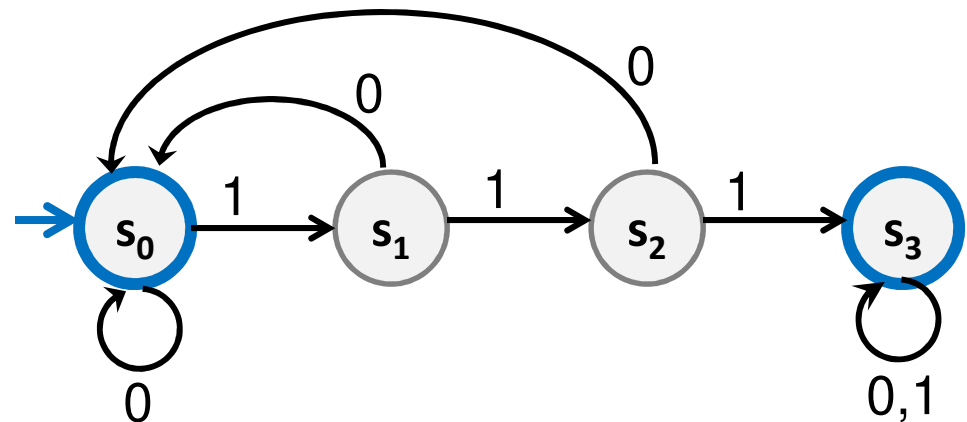
Old State	0	1
s_0	s_0	s_1
s_1	s_0	s_2
s_2	s_0	s_3
s_3	s_3	s_3



Finite State Machines

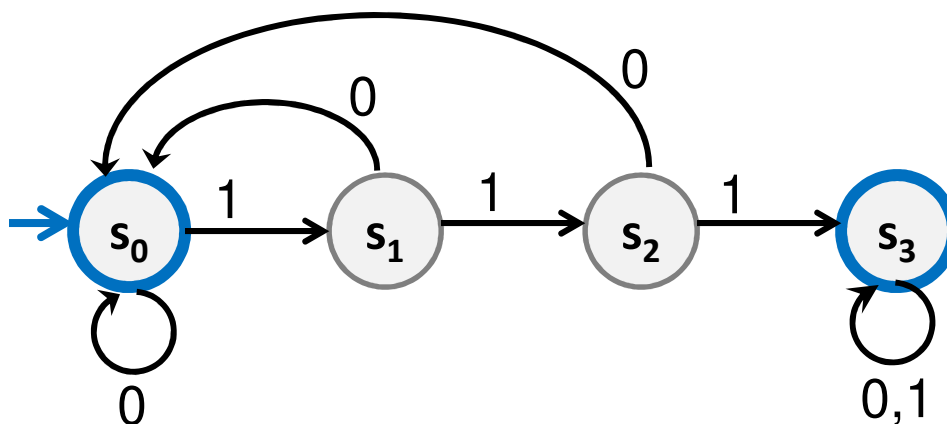
- Each machine designed for strings over some fixed alphabet Σ .
- Must have a transition defined from each state for every symbol in Σ .

Old State	0	1
s_0	s_0	s_1
s_1	s_0	s_2
s_2	s_0	s_3
s_3	s_3	s_3



What language does this machine recognize?

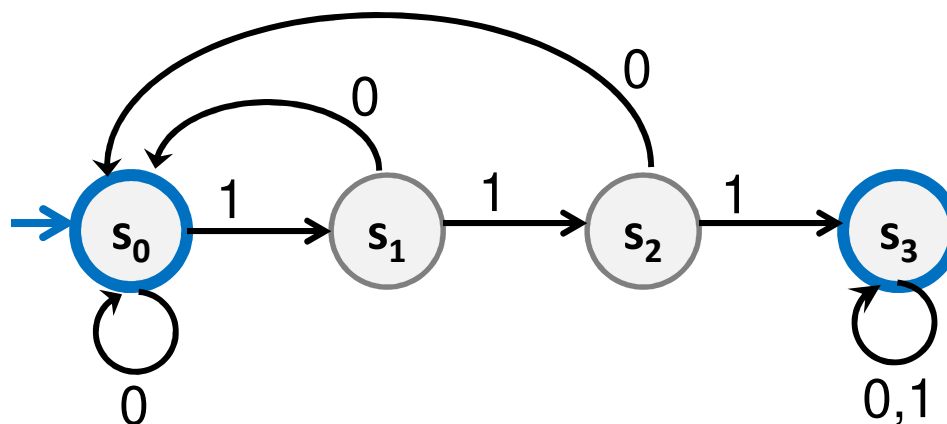
Old State	0	1
s_0	s_0	s_1
s_1	s_0	s_2
s_2	s_0	s_3
s_3	s_3	s_3



What language does this machine recognize?

The set of all binary strings that contain **111** or don't end in **1**

Old State	0	1
s_0	s_0	s_1
s_1	s_0	s_2
s_2	s_0	s_3
s_3	s_3	s_3



Applications of FSMs (a.k.a. Finite Automata)

- Implementation of regular expression matching in programs like `grep`
- Control structures for sequential logic in digital circuits
- Algorithms for communication and cache-coherence protocols
 - Each agent runs its own FSM
- Design specifications for reactive systems
 - Components are communicating FSMs

Applications of FSMs (a.k.a. Finite Automata)

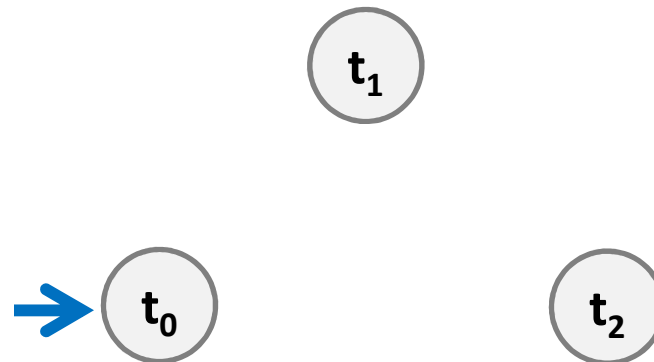
- **Formal verification of systems**
 - Is an unsafe state reachable?
- **Computer games**
 - FSMs provide worlds to explore
- **Minimization algorithms for FSMs can be extended to more general models used in**
 - Text prediction
 - Speech recognition

Strings over $\{0, 1, 2\}$

M_1 : Strings with an even number of 2's

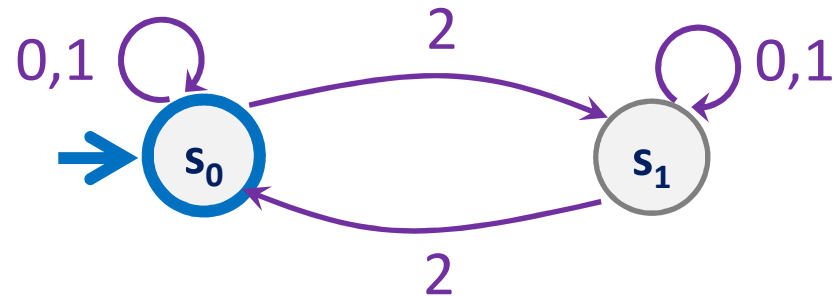


M_2 : Strings where the sum of digits mod 3 is 0

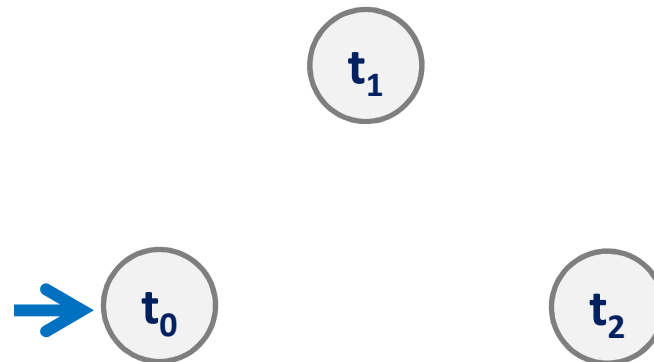


Strings over $\{0, 1, 2\}$

M_1 : Strings with an even number of 2's

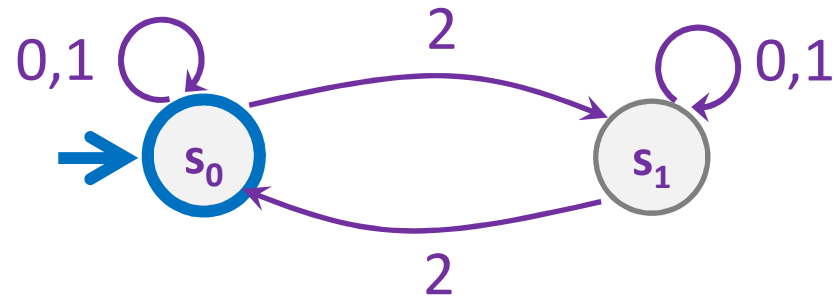


M_2 : Strings where the sum of digits mod 3 is 0

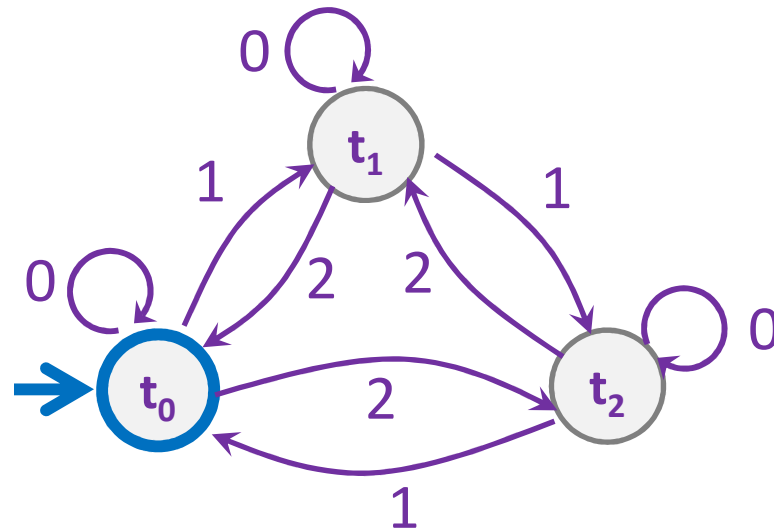


Strings over $\{0, 1, 2\}$

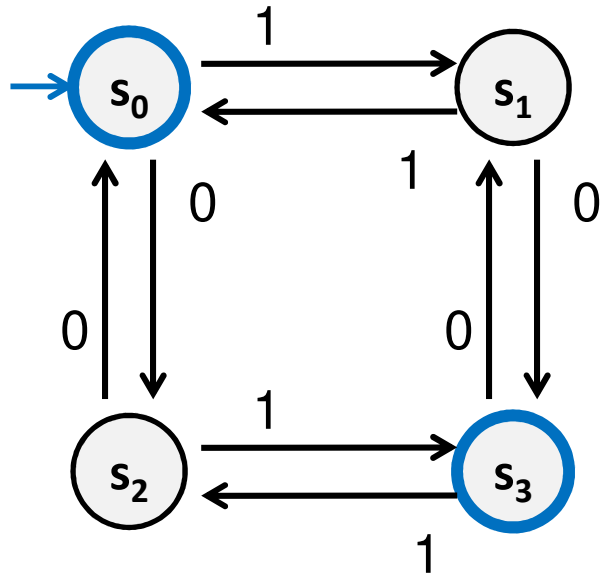
M_1 : Strings with an even number of 2's



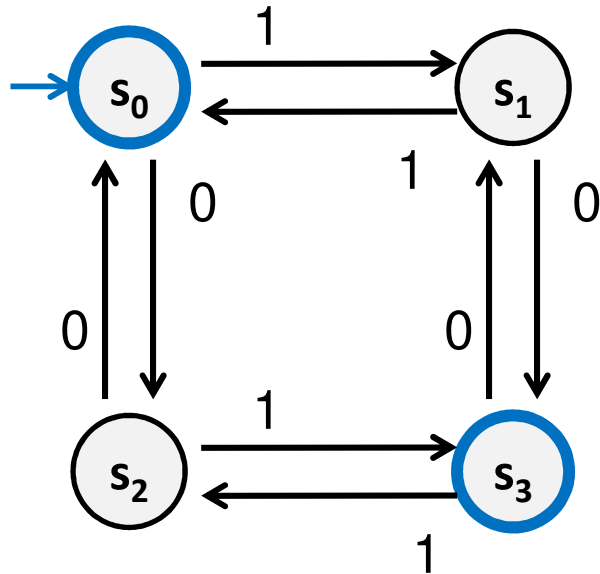
M_2 : Strings where the sum of digits mod 3 is 0



What language does this machine recognize?



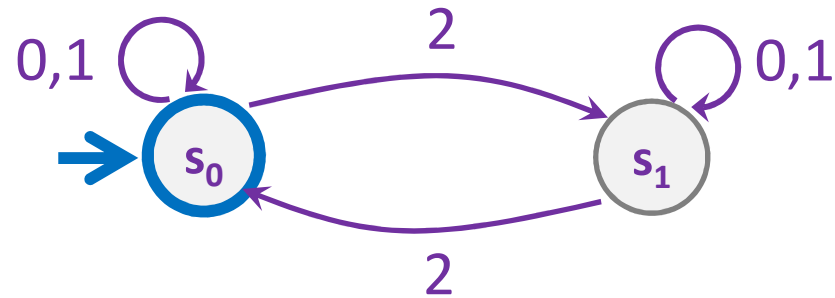
What language does this machine recognize?



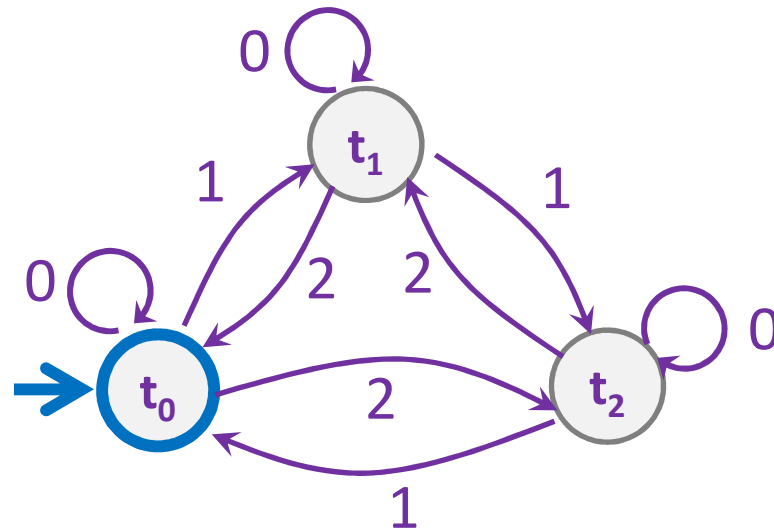
The set of all binary strings with # of 1's \equiv # of 0's (mod 2)
(both are even or both are odd).

Strings over $\{0, 1, 2\}$

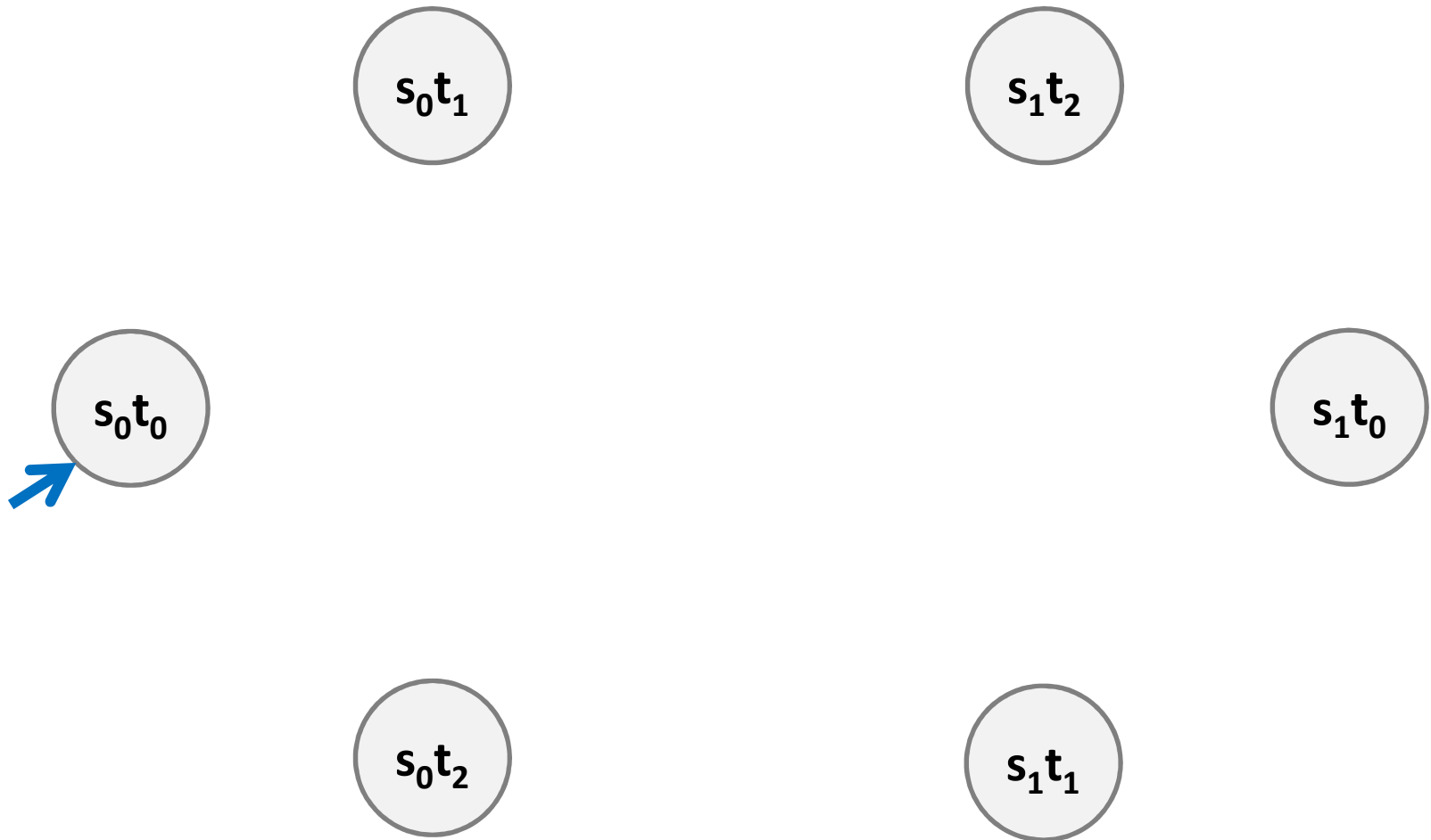
M_1 : Strings with an even number of 2's



M_2 : Strings where the sum of digits mod 3 is 0



Strings over $\{0,1,2\}$ w/ even number of 2's and mod 3 sum 0



Strings over $\{0,1,2\}$ w/ even number of 2's and mod 3 sum 0

