CSE 311 Lecture 23: Finite State Machines

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Topics

Finite state machines (FSMs)
  Definition and examples.

Finite state machines with output
  Definition and examples.
Finite state machines (FSMs)

Definition and examples.
Finite state machines by example

An FSM recognizes (or accepts) a set of strings.
Finite state machines by example

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(1) Given a string \( s \), begin at the start state, denoted by an incoming arrow with no source.
Finite state machines by example

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2. If $s = aw$, take the edge labeled $a$ to get to the next state.
3. Otherwise ($s$ is empty), stop.
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4. Let $s$ be $w$ and repeat (2)-(4) until the machine stops.

The FSM *accepts* $s$ iff it stops in a *final state*, denoted by a double circle.
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<thead>
<tr>
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<th>Recognized?</th>
</tr>
</thead>
<tbody>
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<td>no</td>
</tr>
<tr>
<td>0</td>
<td>yes</td>
</tr>
<tr>
<td>0110</td>
<td>yes</td>
</tr>
<tr>
<td>1011</td>
<td>no</td>
</tr>
</tbody>
</table>

All binary strings that end in 0.
Defining an FSM

Finite state machine (FSM)

A finite state machine (FSM) $M = (S, \Sigma_{\text{in}}, f, s_0, F)$ consists of
a finite set of states $S$, a finite input alphabet $\Sigma_{\text{in}}$, a transition function $f$ that maps each state in $S$ and input in $\Sigma_{\text{in}}$ to a state in $S$, a start state $s_0 \in S$, and a set of final states $F \subseteq S$. 
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<table>
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<tr>
<th>state</th>
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<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>$s_0$</td>
<td>$s_0$</td>
<td>$s_1$</td>
</tr>
<tr>
<td>$s_1$</td>
<td>$s_0$</td>
<td>$s_2$</td>
</tr>
<tr>
<td>$s_2$</td>
<td>$s_0$</td>
<td>$s_3$</td>
</tr>
<tr>
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What language does this FSM accept?
Defining an FSM

Finite state machine (FSM) or deterministic finite automaton (DFA)

A finite state machine (FSM) \( M = (S, \Sigma_{\text{in}}, f, s_0, F) \) consists of an input alphabet \( \Sigma_{\text{in}} \), a set of states \( S \), an initial state \( s_0 \in S \), and a set of final states \( F \subseteq S \).

The set of all binary strings that contain 111 or end in 0.
Applications of FSMs

Implementation of string matching.
   For example, in grep.

Algorithms for communication and cache-coherence protocols.
   Each agent runs its own FSM.

Specifications for controllers.
   For example, device drivers.

Specifications for reactive systems.
   Components are communicating FSMs.

Verification.
   Is an unsafe state reachable?
Example: an FSM over binary strings

What language does this FSM recognize?
Example: an FSM over binary strings

What language does this FSM recognize?

The set of all binary strings where the numbers of 1’s and 0’s are congruent mod 2. That is, both numbers are even or both are odd.
Example: FSMs over \$\{0, 1, 2\}\$

\(M_1\): strings with an even number of 2’s

\(M_2\): strings where the sum of digits mod 3 is 0
Example: FSMs over \{0, 1, 2\}

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\(M_1\): strings with an even number of 2’s

\(M_2\): strings where the sum of digits mod 3 is 0
Example: a combined FSM over \{0, 1, 2\}

\(M_1\) and \(M_2\): strings with an even number of 2’s where digit sum mod 3 is 0
Example: a combined FSM over \{0, 1, 2\}

\(M_1\) and \(M_2\): strings with an even number of 2’s where digit sum mod 3 is 0
Example: another combined FSM over \( \{0, 1, 2\} \)

\( M_1 \) or \( M_2 \): strings with an even number of 2’s or digit sum mod 3 is 0
Finite state machines with output

Definition and examples.
Adding output to FSMs

FSMs can do more than just accept or reject strings.
   So DFAs aren’t the only kind of FSM!

Another useful kind of FSM can also return output.
   Whenever you enter specific states, the FSM produces an output.
   These FSMs (Moore machines) are used as controllers.
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Finite state machine (FSM) with output

A finite state machine (FSM) \( M = (S, \Sigma_{\text{in}}, \Sigma_{\text{out}}, f, g, s_0) \) consists of:
- a finite set of states \( S \),
- a finite input alphabet \( \Sigma_{\text{in}} \),
- a finite output alphabet \( \Sigma_{\text{out}} \),
- a transition function \( f \) that maps each state in \( S \) and input in \( \Sigma_{\text{in}} \) to a state in \( S \),
- an output function \( g \) that maps each state in \( S \) to an output symbol in \( \Sigma_{\text{out}} \),
- and start state \( s_0 \in S \).
Example: vending machine

Consider a simple vending machine.
- Enter 15 cents in dimes or nickels.
- Press S (Snickers) or B (Butterfinger) for a candy bar.
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Basic transitions on N (nickel), D (dime), B (Butterfinger), and S (Snickers).
Adding output to states: N (nickel), B (Butterfinger), and S (Snickers).
Example: complete vending machine with output

Add transitions to cover all symbols for each state.
Summary

A finite state machine (FSM) consists of states and transitions.
  States include the start state and final states.
  Transitions map states and input symbols to states.
  Also known as deterministic finite automata (DFAs).

An FSM recognizes a set of strings (language).
  These are all strings that reach a final state from the start.

FSMs have many applications.
  Regular expression matching, cache coherence protocols.
  Verification (e.g., of OS kernels and devices).
  Controllers (FSMs with output).