

CSE 311 Autumn 20 Lecture 14

```
//Assume i is a nonnegative integer
//returns 2^i.
public int CalculatesTwoToTheI(int
          return 1;
     else
          return 2*CaclulatesTwoToTheI(i-1);
```

Why does Calculates Two To The I (4) calculate 2^4? Convince the other people in your room

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The the base casaishit the function returns I bell up the steek and multiplies by 2 at each step "

Bose cure makes sonse,

Recursive core also, it is o multiply by 2 and subtract (from i.

Something like this:

```
Well, as long as CalculatesTwoToTheI(3) = 8, we get 16... Which happens as long as CalculatesTwoToTheI(2) = 4 Which happens as long as CalculatesTwoToTheI(1) = 2 Which happens as long as CalculatesTwoToTheI(0) = 1 And it is! Because that's what the base case says.
```

There's really only two cases.

The Base Case is Correct

THEN our value is correct.

```
CalculatesTwoToTheI(0) = 1 (which it should!)

And that means CalculatesTwoToTheI(1) = 2, (like it should)

And that means CalculatesTwoToTheI(2) = 4, (like it should)

And that means CalculatesTwoToTheI(3) = 8, (like it should)

And that means CalculatesTwoToTheI(4) = 16, (like it should)

IF the recursive call we make is correct
```

The code has two big cases,
So our proof had two big cases

"The base case of the code produces the correct output"

"IF the calls we rely on produce the correct output THEN the current call produces the right output"

A bit more formally...

"The base case of the code produces the correct output"

"IF the calls we rely on produce the correct output THEN the current call produces the right output"

Let P(i) be "Calculates Two To The I (i)" returns 2^i .

How do we know P(4)?

P(0) is true.

And $P(0) \rightarrow P(1)$ so P(1).

And $P(1) \rightarrow P(2)$, so P(2).

And $P(2) \rightarrow P(3)$, so P(3).

And $P(3) \rightarrow P(4)$, so P(4).

A bit more formally...

This works alright for P(4).

What about P(1000)? P(100000000)?

At this point, we'd need to show that implication $P(k) \rightarrow P(k+1)$ for A BUNCH of values of k.

But the code is the same each time.

And so was the argument!

We should instead show $\forall k[P(k) \rightarrow P(k+1)]$.

Induction

Your new favorite proof technique!

How do we show $\forall n, P(n)$?

 $\langle Show P(0) \rangle$

Show $\forall k(P(k) \rightarrow P(k+1))$

```
Induction

//Assume i is a nonnegative integer
public int CalculatesTwoToTheI(int i) {
    if (i == 0)
        return 1;
    else
    return 2*CaclulatesTwoToTheI(i-1);
}
```

Let P(i) be "CalculatesTwoToTheI(i)" returns 2^i .

Note that if the input i is 0, then the if-statement evaluates to true, and

 $1 = 2^0$ is returned, so P(0) is true.

Suppose P(k) holds for an arbitrary $k \ge 0$.

So P(k+1) holds.

Therefore P(n) holds for all $n \geq 0$ by the principle of induction.

Making Induction Proofs Pretty

(Let P(i) be "CalculatesTwoToTheI(i)" returns 2^i .

Base Case (i = 0) Note that if the input i is 0, then the if-statement evaluates to true, and $1 = 2^0$ is returned, so P(0) is true.

(Inductive Hypothesis: Suppose P(k) holds for an arbitrary $k \geq 0$.

Inductive Step: Since $k \ge 0$, $k \ge 1$, so the code goes to the recursive case. We will return $2 \cdot CalculatesTwoToTheI(k)$. By Inductive Hypothesis,

CalculatesTwoToTheI(k) = 2^k . Thus we return $2 \cdot 2^k = 2^{k+1}$.

So P(k + 1) holds.

Therefore P(n) holds for all $n \ge 0$ by the principle of induction.

Making Induction Proofs Pretty

All of our induction proofs will come in 5 easy(?) steps!

- 1. Define P(n). State that your proof is by induction on n.
- 2. Show P(0) i.e. show the base case
- 3. Suppose P(k) for an arbitrary k.
- 4. Show P(k + 1) (i.e. get $P(k) \rightarrow P(k + 1)$)
- 5. Conclude by saying P(n) is true for all n by induction.

Some Other Notes

Always state where you use the inductive hypothesis when you're using it in the inductive step.

It's usually the key step, and the reader really needs to focus on it.

Be careful about what values you're assuming the Inductive Hypothesis for – the smallest possible value of k should assume the base case but nothing more.

The Principle of Induction (formally)

Principle of Induction
$$P(0); \forall k(P(k) \rightarrow P(k+1))$$
$$\therefore \qquad \forall n(P(n))$$

Informally: if you knock over one domino, and every domino knocks over the next one, then all your dominoes fell over.

What's Induction Good For

Induction lets us rigorously prove statements that hold for all natural numbers (or all positive integers, or all positive integers greater than 7, or all even numbers, or...)

It's the formal way of saying "and this pattern continues."

We've occasionally done this in hand-wavy ways before – we know have the tool to do it carefully.

The extra structure lets us detect mistakes more easily.

More Induction

Induction doesn't **only** work for code!

Show that
$$\sum_{i=0}^{n} 2^i = 1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1$$
.

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Let
$$P(n) = \sum_{i=0}^{n} 2^{i} = 2^{n+1} - 1$$
."

We show P(n) holds for all n by induction on n.

Base Case ()

Inductive Hypothesis:

Inductive Step:

P(n) holds for all $n \ge 0$ by the principle of induction.

More Induction

Induction doesn't only work for code!

Show that
$$\sum_{i=0}^{n} 2^i = 1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1$$
.

Let
$$P(n) = \sum_{i=0}^{n} 2^i = 2^{n+1} - 1$$
."

We show P(n) holds for all n by induction on n.

Base Case
$$(n = 0) \sum_{i=0}^{0} 2^i = 1 = 2 - 1 = 2^{0+1} - 1$$
.

Inductive Hypothesis: Suppose P(k) holds for an arbitrary $k \geq 0$.

Inductive Step: We show P(k+1). Consider the summation $\sum_{i=0}^{k+1} 2^i = 2^{k+1} + \sum_{i=0}^k 2^i = 2^{k+1} + 2^{k+1} - 1$, where the last step is by IH.

Simplifying, we get:
$$\sum_{i=0}^{k+1} 2^i = 2^{k+1} + 2^{k+1} - 1 = 2 \cdot 2^{k+1} - 1 = 2^{(k+1)+1} - 1$$
.

P(n) holds for all $n \geq 0$ by the principle of induction.