

CSE 311: Foundations of Computing I

Section 6: Induction Solutions

1. Extended Euclidean Algorithm

- (a) Find the multiplicative inverse y of 7 mod 33. That is, find y such that $7y \equiv 1 \pmod{33}$. You should use the extended Euclidean Algorithm. Your answer should be in the range $0 \leq y < 33$.

Solution:

First, we find the gcd:

$$\begin{aligned} \gcd(33, 7) &= \gcd(7, 5) & 33 &= \boxed{7} \cdot 4 + 5 & (1) \\ &= \gcd(5, 2) & 7 &= \boxed{5} \cdot 1 + 2 & (2) \\ &= \gcd(2, 1) & 5 &= \boxed{2} \cdot 2 + 1 & (3) \\ &= \gcd(1, 0) & 2 &= 1 \cdot 2 + 0 & (4) \\ &= 1 & & & (5) \end{aligned}$$

Next, we re-arrange equations (1) - (3) by solving for the remainder:

$$\begin{aligned} 1 &= 5 - \boxed{2} \cdot 2 & (6) \\ 2 &= 7 - \boxed{5} \cdot 1 & (7) \\ 5 &= 33 - \boxed{7} \cdot 4 & (8) \\ & & (9) \end{aligned}$$

Now, we backward substitute into the boxed numbers using the equations:

$$\begin{aligned} 1 &= 5 - \boxed{2} \cdot 2 \\ &= 5 - (7 - \boxed{5} \cdot 1) \cdot 2 \\ &= 3 \cdot \boxed{5} - 7 \cdot 2 \\ &= 3 \cdot (33 - \boxed{7} \cdot 4) - 7 \cdot 2 \\ &= 33 \cdot 3 + 7 \cdot -14 \end{aligned}$$

So, $1 = 33 \cdot 3 + \boxed{7} \cdot -14$. Thus, $33 - 14 = 19$ is the multiplicative inverse of 7 mod 33.

- (b) Now, solve $7z \equiv 2 \pmod{33}$.

Solution:

If z is a solution to that equation, then multiplying both sides by 19, we have $z = 1z \equiv 19 \cdot 7z \equiv 19 \cdot 2 \equiv 5 \pmod{33}$. Hence, every solution must be of the form $z = 5 + 33k$ for some $k \in \mathbb{Z}$.

Furthermore, we can see that every number of this form is a solution since $(7(5 + 33k)) \pmod{33} = (35 + 7 \cdot 33k) \pmod{33} = 35 \pmod{33} = 2 = 2 \pmod{33}$.

2. A Strict Inequality

Prove that $6n + 6 < 2^n$ for all $n \geq 6$.

Solution:

Let $P(n)$ be " $6n + 6 < 2^n$ ". We will prove $P(n)$ for all integers $n \geq 6$ by induction.

Base Case ($n = 6$): $6 \cdot 6 + 6 = 42 < 64 = 2^6$, so $P(6)$ holds.

Induction Hypothesis: Assume that $6j + 6 < 2^j$ for an arbitrary integer $j \geq 6$.

Induction Step: Goal: Show $6(j + 1) + 6 < 2^{j+1}$

$$\begin{aligned} 6(j + 1) + 6 &= 6j + 6 + 6 \\ &< 2^j + 6 && \text{[Induction Hypothesis]} \\ &< 2^j + 2^j && \text{[Since } 2^j > 6, \text{ since } j \geq 6\text{]} \\ &< 2 \cdot 2^j \\ &< 2^{j+1}, \end{aligned}$$

which shows that $P(j + 1)$ is true.

Conclusion: $P(n)$ holds for all integers $n \geq 6$ by induction.

3. Divisibility by Induction

Prove that $9 \mid n^3 + (n + 1)^3 + (n + 2)^3$ for all $n > 1$ by induction.

Solution:

Let $P(n)$ be " $9 \mid n^3 + (n + 1)^3 + (n + 2)^3$ ". We will prove $P(n)$ for all integers $n > 1$ by induction.

Base Case ($n = 2$): $2^3 + (2 + 1)^3 + (2 + 2)^3 = 8 + 27 + 64 = 99 = 9 \cdot 11$, so $9 \mid 2^3 + (2 + 1)^3 + (2 + 2)^3$.

Induction Hypothesis: Assume that $9 \mid j^3 + (j + 1)^3 + (j + 2)^3$ for an arbitrary integer $j > 1$. Note that this is equivalent to assuming that $j^3 + (j + 1)^3 + (j + 2)^3 = 9k$ for some integer k .

Induction Step: Goal: Show $9 \mid (j + 1)^3 + (j + 2)^3 + (j + 3)^3$

$$\begin{aligned} (j + 1)^3 + (j + 2)^3 + (j + 3)^3 &= (j + 3)^3 + 9k - j^3 && \text{[Induction Hypothesis]} \\ &= j^3 + 9j^2 + 27j + 27 + 9k - j^3 \\ &= 9j^2 + 27j + 27 + 9k \\ &= 9(j^2 + 3j + 3 + k) \end{aligned}$$

So $9 \mid (j + 1)^3 + (j + 2)^3 + (j + 3)^3$, which is $P(j + 1)$.

Conclusion: $P(n)$ holds for all integers $n > 1$ by induction.

4. Another Inequality

Prove that, for all integers $n \geq 1$, if you have numbers a_1, \dots, a_n and b_1, \dots, b_n , with $\forall i \in [n]. a_i \leq b_i$, then:

$$\sum_{i=1}^n a_i \leq \sum_{i=1}^n b_i$$

Solution:

Let $P(n)$ be the statement “if $a_1 \leq b_1, a_2 \leq b_2, \dots, a_n \leq b_n$, then $\sum_{i=1}^n a_i \leq \sum_{i=1}^n b_i$ ”. We prove that $P(n)$ is true for all integers $n \geq 1$ by induction on n :

Base Case ($n = 1$). Suppose $a_1 \leq b_1$. Using the definition of summation, we can see that

$$\sum_{i=1}^n a_i = \sum_{i=1}^1 a_i = a_1 \leq b_1 = \sum_{i=1}^1 b_i = \sum_{i=1}^n b_i,$$

so the claim is true for $n = 1$.

Induction Hypothesis. Suppose that $P(k)$ holds for an arbitrary integer $k \geq 1$.

Induction Step. Suppose that $a_1 \leq b_1, a_2 \leq b_2, \dots, a_{k+1} \leq b_{k+1}$. Then, we can calculate

$$\begin{aligned} \sum_{i=1}^{k+1} a_i &= \sum_{i=1}^k a_i + a_{k+1} && \text{[Splitting the summation]} \\ &\leq \sum_{i=1}^k b_i + a_{k+1} && \text{[By IH]} \\ &\leq \sum_{i=1}^k b_i + b_{k+1} && \text{[By Assumption]} \\ &\leq \sum_{i=1}^{k+1} b_i && \text{[Algebra]} \end{aligned}$$

This shows $P(k + 1)$.

Therefore, we have shown the claim for all $n \in \mathbb{N}$ by induction.