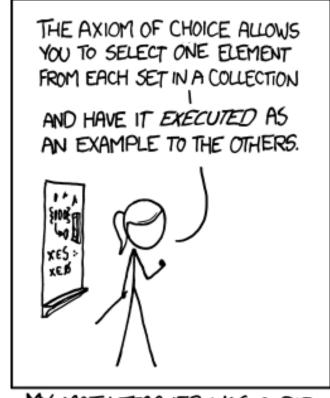
## **CSE 311: Foundations of Computing**

## **Lecture 8: Predicate Logic Proofs**



MY MATH TEACHER WAS A BIG BELIEVER IN PROOF BY INTIMIDATION.

## Last class: Propositional Inference Rules

Two inference rules per binary connective, one to eliminate it and one to introduce it

Elim 
$$\land$$
  $A \land B$ 
 $\therefore$   $A, B$ 

$$A \lor B; \neg A$$
 $\therefore$   $A \lor B; \neg A$ 
 $\therefore$   $A \lor B, B \lor A$ 

Modus Ponens
$$A; A \to B$$

$$\therefore B$$

Direct Proof
$$Rule$$

$$A \Rightarrow B$$

$$A \Rightarrow B$$

$$A \Rightarrow B$$

$$A \Rightarrow B$$

Not like other rules

## **One General Proof Strategy**

- Look at the rules for introducing connectives to see how you would build up the formula you want to prove from pieces of what is given
- 2. Use the rules for eliminating connectives to break down the given formulas so that you get the pieces you need to do 1.
- 3. Write the proof beginning with what you figured out for 2 followed by 1.

## Example

Prove:  $((p \rightarrow q) \land (q \rightarrow r)) \rightarrow (p \rightarrow r)$ 

## Last class: Example

Prove: 
$$((p \rightarrow q) \land (q \rightarrow r)) \rightarrow (p \rightarrow r)$$

1.1. 
$$(p \rightarrow q) \land (q \rightarrow r)$$
 Assumption  
1.2.  $p \rightarrow q$   $\land$  Elim: 1.1  
1.3.  $q \rightarrow r$   $\land$  Elim: 1.1  
1.4.1.  $p$  Assumption  
1.4.2.  $q$  MP: 1.2, 1.4.1  
1.4.3.  $r$  MP: 1.3, 1.4.2  
1.4.  $p \rightarrow r$  Direct Proof Rule

1.  $((p \rightarrow q) \land (q \rightarrow r)) \rightarrow (p \rightarrow r)$  Direct Proof Rule

## Inference Rules for Quantifiers: First look

P(c) for some c
$$\therefore \exists x P(x)$$

$$\begin{array}{c|c}
 & \forall x P(x) \\
 & \therefore P(a) \text{ for any } a
\end{array}$$

Let a be arbitrary\*"...P(a)
$$\therefore \forall x P(x)$$

$$\Rightarrow P(c) \text{ for some } special** c$$
\* in the domain of P

\*\* By special, we mean that c is a name for a value where P(c) is true. We can't use anything else about that value, so c has to be a NEW name!

## **Predicate Logic Proofs**

- Can use
  - Predicate logic inference rules whole formulas only
  - Predicate logic equivalences (De Morgan's)
     even on subformulas
  - Propositional logic inference rules whole formulas only
  - Propositional logic equivalences
     even on subformulas

$$\begin{array}{c}
P(c) \text{ for some } c \\
\therefore \quad \exists x P(x)
\end{array}$$

$$\begin{array}{c|c}
 & \forall x \ P(x) \\
 & \therefore \ P(a) \ \text{for any } a
\end{array}$$

Prove 
$$\forall x P(x) \rightarrow \exists x P(x)$$

$$5. \quad \forall x P(x) \rightarrow \exists x P(x)$$

The main connective is implication so Direct Proof Rule seems good

$$\begin{array}{c}
 P(c) \text{ for some c} \\
 \vdots \quad \exists x P(x)
\end{array}$$

$$\frac{\forall x P(x)}{\therefore P(a) \text{ for any } a}$$

Prove 
$$\forall x P(x) \rightarrow \exists x P(x)$$

1.1.  $\forall x P(x)$  Assumption

We need an ∃ we don't have so "intro ∃" rule makes sense

1.5. 
$$\exists x P(x)$$

1. 
$$\forall x P(x) \rightarrow \exists x P(x)$$
 Direct Proof Rule

$$\frac{\forall x P(x)}{\therefore P(a) \text{ for any } a}$$

Prove 
$$\forall x P(x) \rightarrow \exists x P(x)$$

**1.1.**  $\forall x P(x)$  Assumption

We need an ∃ we don't have so "intro ∃" rule makes sense

**1.5.** 
$$\exists x P(x)$$

That requires P(c) for some c.

1. 
$$\forall x P(x) \rightarrow \exists x P(x)$$
 Direct Proof Rule

$$\frac{\forall x P(x)}{\therefore P(a) \text{ for any } a}$$

Prove 
$$\forall x P(x) \rightarrow \exists x P(x)$$

1.1. 
$$\forall x P(x)$$
 Assumption   
1.2. Let  $a$  be an object.  
1.3.  $P(a)$  Elim  $\forall$ : 1.1

**Elim** ∀: **1.1** 

We could have picked any name or domain expression here.

**1.5.** 
$$\exists x P(x)$$
 Intro  $\exists$ : ?

1.  $\forall x P(x) \rightarrow \exists x P(x)$  Direct Proof Rule

$$\begin{array}{c}
P(c) \text{ for some c} \\
\therefore \quad \exists x P(x)
\end{array}$$

$$\begin{array}{c|c}
 & \forall x \ P(x) \\
 & \therefore \ P(a) \ \text{for any } a
\end{array}$$

Prove 
$$\forall x P(x) \rightarrow \exists x P(x)$$

No holes. Just need to clean up.

1.1. 
$$\forall x P(x)$$
 Assumption

1.2. Let  $\alpha$  be an object.

**1.3.** P(a) Elim  $\forall$ : **1.1** 

**1.5.** 
$$\exists x P(x)$$
 Intro  $\exists$ : **1.3**

1.  $\forall x P(x) \rightarrow \exists x P(x)$  Direct Proof Rule

P(c) for some c
$$\therefore \exists x P(x)$$

$$\frac{\forall x P(x)}{\therefore P(a) \text{ for any } a}$$

Prove  $\forall x P(x) \rightarrow \exists x P(x)$ 

1.1.  $\forall x P(x)$  Assumption

1.2. Let  $\alpha$  be an object.

**1.3.** P(a) Elim  $\forall$ : **1.1** 

**1.4.**  $\exists x P(x)$  Intro  $\exists$ : **1.3** 

1.  $\forall x P(x) \rightarrow \exists x P(x)$  Direct Proof Rule

Working forwards as well as backwards:

In applying "Intro ∃" rule we didn't know what expression we might be able to prove P(c) for, so we worked forwards to figure out what might work.

## **Predicate Logic Proofs with more content**

- In propositional logic we could just write down other propositional logic statements as "givens"
- Here, we also want to be able to use domain knowledge so proofs are about something specific
- Example: Domain of Discourse Integers
- Given the basic properties of arithmetic on integers, define:

  Predicate Definitions

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
Odd(x)  $\equiv \exists y (x = 2 \cdot y + 1)$ 

## A Not so Odd Example

# Domain of Discourse Integers

### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
Odd(x)  $\equiv \exists y (x = 2 \cdot y + 1)$ 

Prove "There is an even number"

Formally: prove  $\exists x \; Even(x)$ 

## A Not so Odd Example

### **Domain of Discourse**

Integers

### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$

$$Odd(x) \equiv \exists y (x = 2 \cdot y + 1)$$

Prove "There is an even number"

Formally: prove  $\exists x \; \text{Even}(x)$ 

1. 
$$2 = 2 \cdot 1$$
 Arithmetic

**2.** 
$$\exists y (2 = 2 \cdot y)$$
 Intro  $\exists : 1$ 

4. 
$$\exists x \, \text{Even}(x)$$
 Intro  $\exists : 3$ 

## A Prime Example

# Domain of Discourse Integers

#### **Predicate Definitions**

```
Even(x) \equiv \exists y (x = 2 \cdot y)

Odd(x) \equiv \exists y (x = 2 \cdot y + 1)

Prime(x) \equiv "x > 1 and x \neq a \cdot b for

all integers a, b with 1 < a < x"
```

Prove "There is an even prime number"

## A Prime Example

# Domain of Discourse Integers

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
Odd(x)  $\equiv \exists y (x = 2 \cdot y + 1)$   
Prime(x)  $\equiv "x > 1$  and  $x \neq a \cdot b$  for  
all integers a, b with  $1 < a < x$ "

Prove "There is an even prime number"

Formally: prove  $\exists x (Even(x) \land Prime(x))$ 

- 1.  $2 = 2 \cdot 1$
- **2.** Prime(**2**)\*

**Arithmetic** 

**Property of integers** 

<sup>\*</sup> Later we will further break down "Prime" using quantifiers to prove statements like this

## A Prime Example

# Domain of Discourse Integers

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
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all integers a, b with  $1 < a < x$ "

**Arithmetic** 

Prove "There is an even prime number"

2 - 2.1

Formally: prove  $\exists x (Even(x) \land Prime(x))$ 

	2 - 2.1	Antimietic
2.	Prime( <b>2</b> )*	<b>Property of integers</b>
3.	$\exists y (2 = 2 \cdot y)$	Intro ∃: 1
4.	Even( <b>2</b> )	Defn of Even: 3
5.	Even( <b>2</b> ) ∧ Prime( <b>2</b> )	Intro ∧: 2, 4

6.  $\exists x (Even(x) \land Prime(x))$  Intro  $\exists : 5$ 

<sup>\*</sup> Later we will further break down "Prime" using quantifiers to prove statements like this

## Inference Rules for Quantifiers: First look

P(c) for some c
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$$\begin{array}{c|c}
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\* in the domain of P

\*\* By special, we mean that c is a name for a value where P(c) is true. We can't use anything else about that value, so c has to be a NEW name!

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers



Prove: "The square of every even number is even."

Formal proof of:  $\forall x (Even(x) \rightarrow Even(x^2))$ 

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1. Let a be an arbitrary integer

- 2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)
- 3.  $\forall x (Even(x) \rightarrow Even(x^2))$



Intro  $\forall$ : 1,2

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**2.1** Even(a)

**Assumption** 

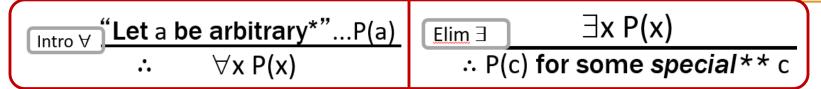
- 2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)
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Direct proof rule

Intro  $\forall$ : 1,2

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Prove: "The square of every even number is even."

Formal proof of:  $\forall x (Even(x) \rightarrow Even(x^2))$ 

- 1. Let a be an arbitrary integer
  - **2.1** Even(a)

Assumption

2.2  $\exists y (a = 2y)$ 

**Definition of Even** 

**2.5** 
$$\exists y (a^2 = 2y)$$

?

2.6 Even(a<sup>2</sup>)

**Definition of Even** 

2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)

Direct proof rule

3.  $\forall x (Even(x) \rightarrow Even(x^2))$ 

Intro ∀: 1,2

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers



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Assumption

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**Definition of Even** 

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$$\exists y (a^2 = 2y)$$

Intro∃rule: 🕐

Need  $a^2 = 2c$  for some c

2.6 Even(a<sup>2</sup>)

**Definition of Even** 

2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)

Direct proof rule

3.  $\forall x (Even(x) \rightarrow Even(x^2))$ 

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1. Let a be an arbitrary integer

2.2 
$$\exists y (a = 2y)$$
 Definition of Even

2.3 
$$a = 2b$$
 Elim  $\exists$ : b special depends on a

**2.5** 
$$\exists y (a^2 = 2y)$$

2. Even(a)
$$\rightarrow$$
Even(a<sup>2</sup>)

3. 
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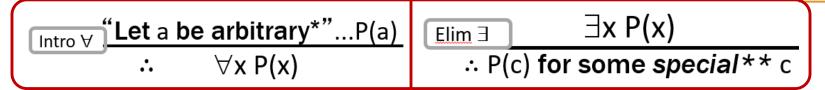
Intro 
$$\exists$$
 rule: Properties Prope

Direct proof rule

Intro 
$$\forall$$
: 1,2

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers

Used  $a^2 = 2c$  for  $c=2b^2$ 



Prove: "The square of every even number is even."

Formal proof of:  $\forall x (Even(x) \rightarrow Even(x^2))$ 

1. Let a be an arbitrary integer

21	Even(a)	Assumption
<b>Z. _</b>	LVCII(a)	Assumption

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$$\exists y (a = 2y)$$
 Definition of Even

2.3 
$$a = 2b$$
 Elim  $\exists$ : b special depends on a

**2.4** 
$$a^2 = 4b^2 = 2(2b^2)$$
 Algebra

**2.5** 
$$\exists y (a^2 = 2y)$$
 Intro  $\exists$  rule

2. Even(a)
$$\rightarrow$$
Even(a<sup>2</sup>) Direct proof rule

3. 
$$\forall x \text{ (Even(x)} \rightarrow \text{Even(x}^2\text{))}$$
 Intro  $\forall : 1,2$ 

## Why did we need to say that **b** depends on **a**?

There are extra conditions on using these rules:

"Let a be arbitrary\*"...P(a)
$$\therefore \forall x \ P(x)$$

$$\vdots \ P(c) \ \text{for some } special** c$$

$$* \text{ in the domain of P}$$

$$** c \text{ has to be a NEW name.}$$

Over integer domain:  $\forall x \exists y (y \ge x)$  is True but  $\exists y \forall x (y \ge x)$  is False

### **BAD "PROOF"**

- **1.**  $\forall x \exists y (y \ge x)$  Given
- 2. Let a be an arbitrary integer
- 3.  $\exists y (y \ge a)$  Elim  $\forall : 1$
- 4.  $b \ge a$  Elim  $\exists$ : b special depends on a
- 5.  $\forall x (b \ge x)$  Intro  $\forall : 2,4$
- 6.  $\exists y \forall x (y \ge x)$  Intro  $\exists : 5$

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There are extra conditions on using these rules:

Let a be arbitrary\*"...P(a)

∴ 
$$\forall x P(x)$$

\* in the domain of P

Elim∃  $\exists x P(x)$ 

∴  $P(c)$  for some special\*\* c

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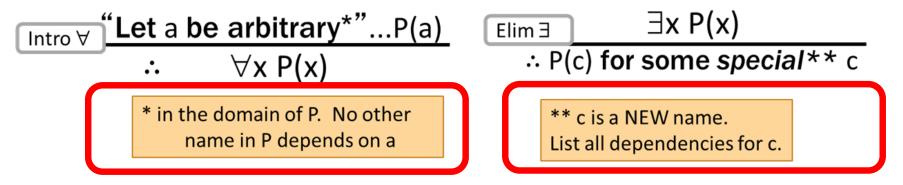
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Can't get rid of a since another name in the same line, b, depends on it!

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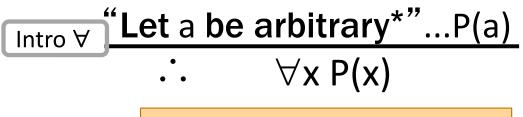
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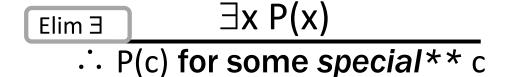
## Inference Rules for Quantifiers: Full version

P(c) for some c
$$\therefore \exists x P(x)$$

$$\begin{array}{c|c}
 & \forall x \ P(x) \\
 & \therefore \ P(a) \ \text{for any } a
\end{array}$$



\* in the domain of P. No other name in P depends on a



\*\* c is a NEW name. List all dependencies for c.

## **English Proofs**

- We often write proofs in English rather than as fully formal proofs
  - They are more natural to read

- English proofs follow the structure of the corresponding formal proofs
  - Formal proof methods help to understand how proofs really work in English...
    - ... and give clues for how to produce them.