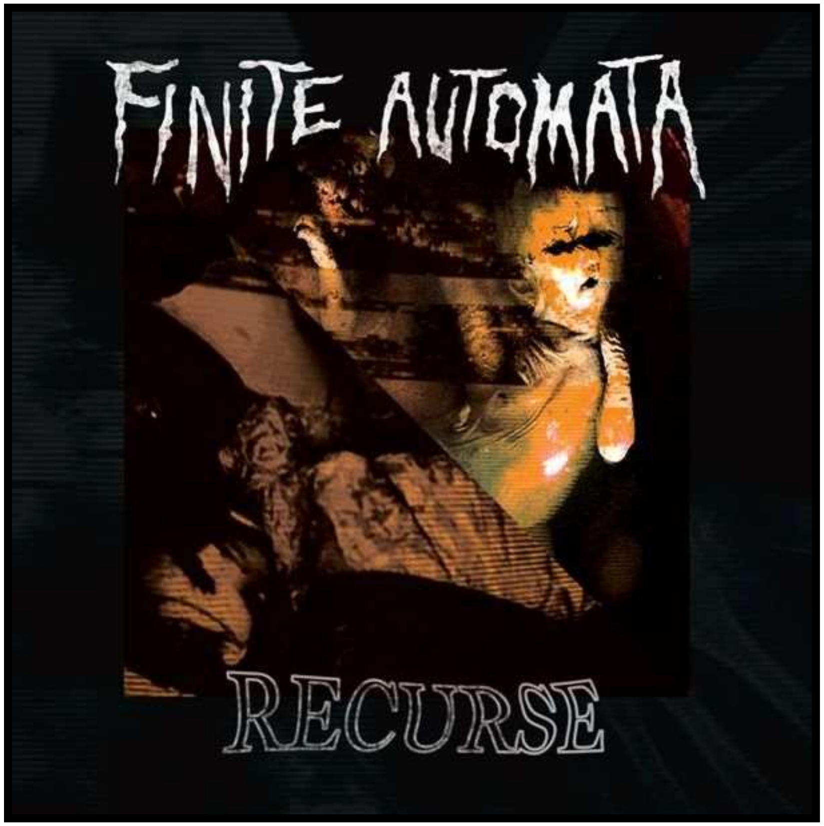
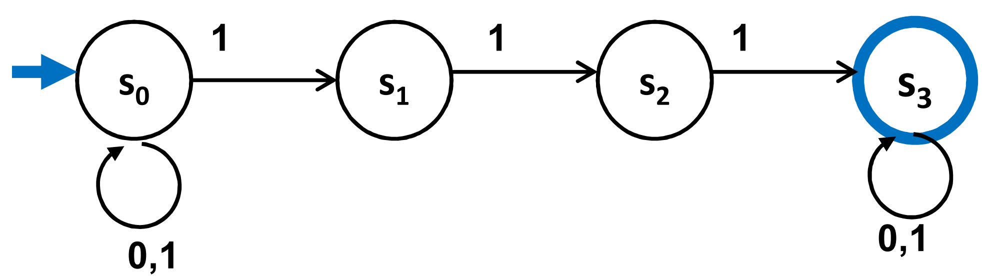
CSE 311: Foundations of Computing

Lecture 24: NFAs, Regular expressions, and NFA→DFA



Last time: Nondeterministic Finite Automata (NFA)

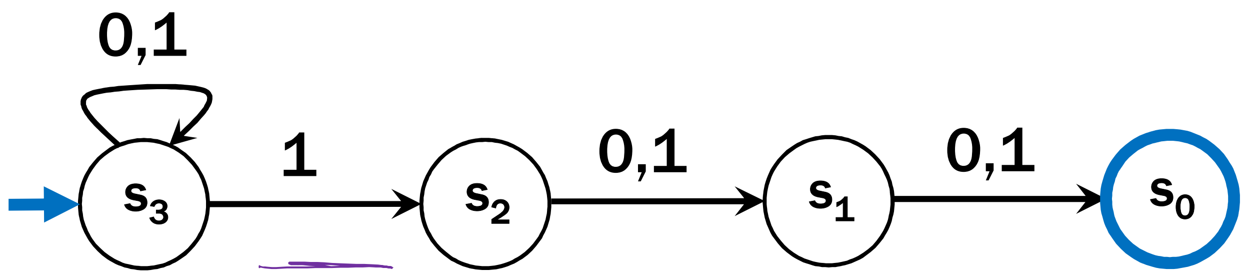
- Graph with start state, final states, edges labeled by symbols (like DFA) but
  - Not required to have exactly 1 edge out of each state labeled by each symbol— can have 0 or >1
  - Also can have edges labeled by empty string ε
- **Defn:** x is in the language recognized by an NFA if and only if x labels a path from the start state to some final state

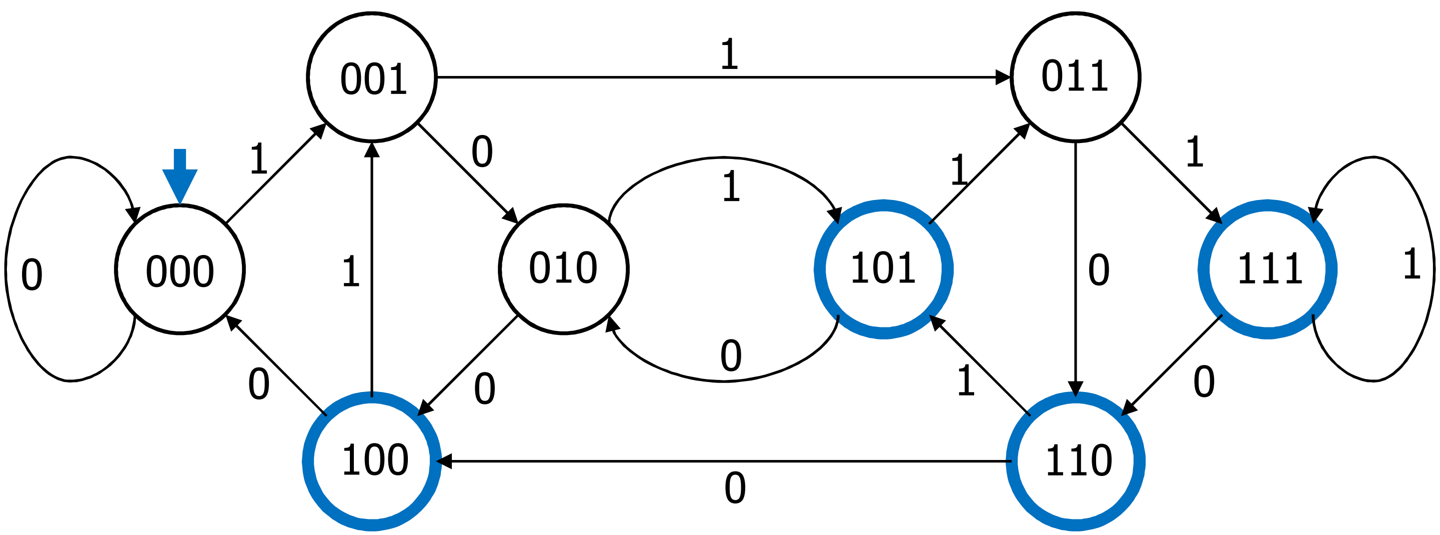


Last time: Three ways of thinking about NFAs

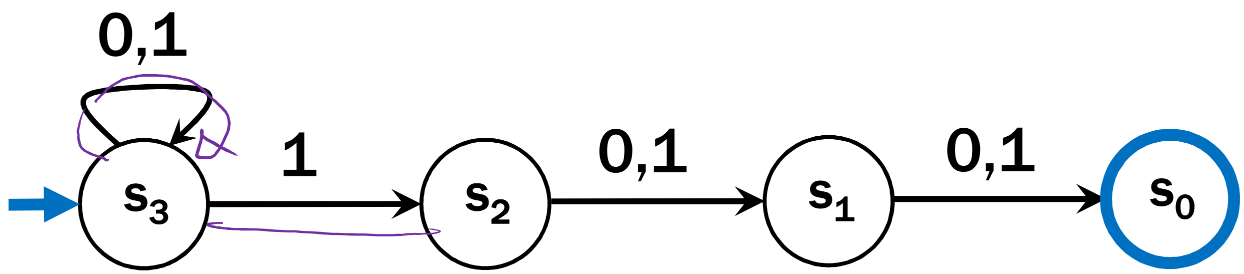
- Outside observer: Is there a path labeled by x from the start state to some final state?
- Perfect guesser: The NFA has input x and whenever there is a choice of what to do it magically guesses a good one (if one exists)
- Parallel exploration: The NFA computation runs all possible computations on x step-by-step at the same time in parallel

Last time: Compare with the smallest DFA

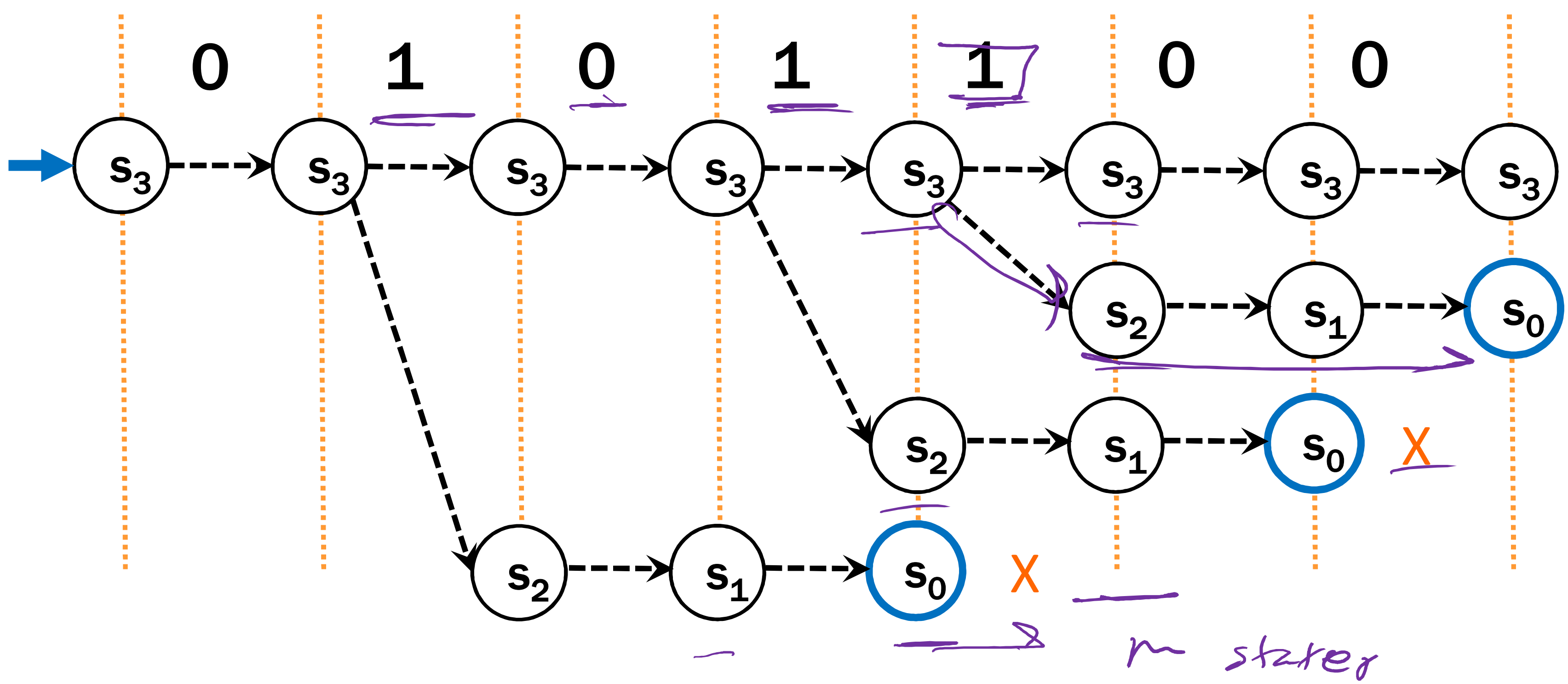




Last time: Parallel Exploration view of an NFA



Input string 0101100



NFAs and regular expressions

Theorem: For any set of strings (language) A described by a regular expression, there is an NFA that recognizes A.

Proof idea: Structural induction based on the recursive definition of regular expressions...

Regular Expressions over Σ

- **Basis:**
  - $\varnothing$, ε are regular expressions
  - *a* is a regular expression for any $a \in \Sigma$
- **Recursive step:**
  - If A and B are regular expressions then so are:

    $(A \cup B)$

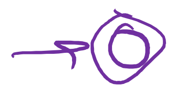
    (AB)

    A^*

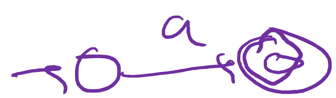
Base Case

- **Case ∅:**

- **Case ε:**

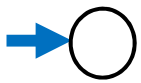


- **Case *a*:**

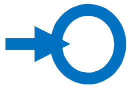


Base Case

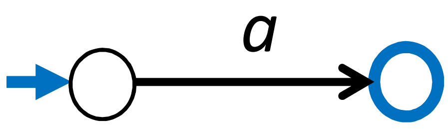
- **Case ∅:**



- **Case ε:**

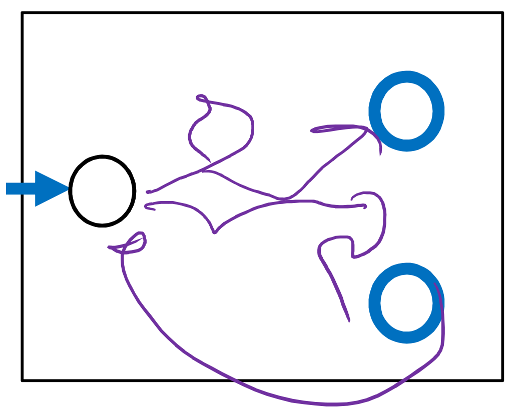


- **Case *a*:**

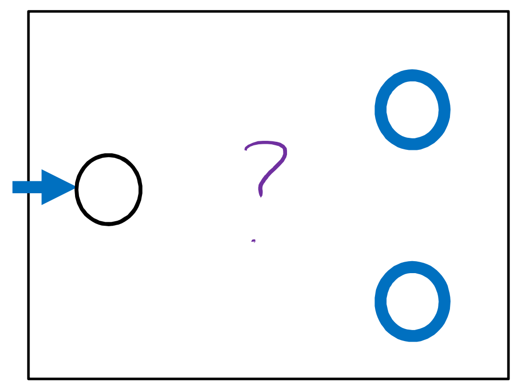


Inductive Hypothesis

- Suppose that for some regular expressions A and B there exist NFAs N_A and N_B such that N_A recognizes the language given by A and N_B recognizes the language given by B



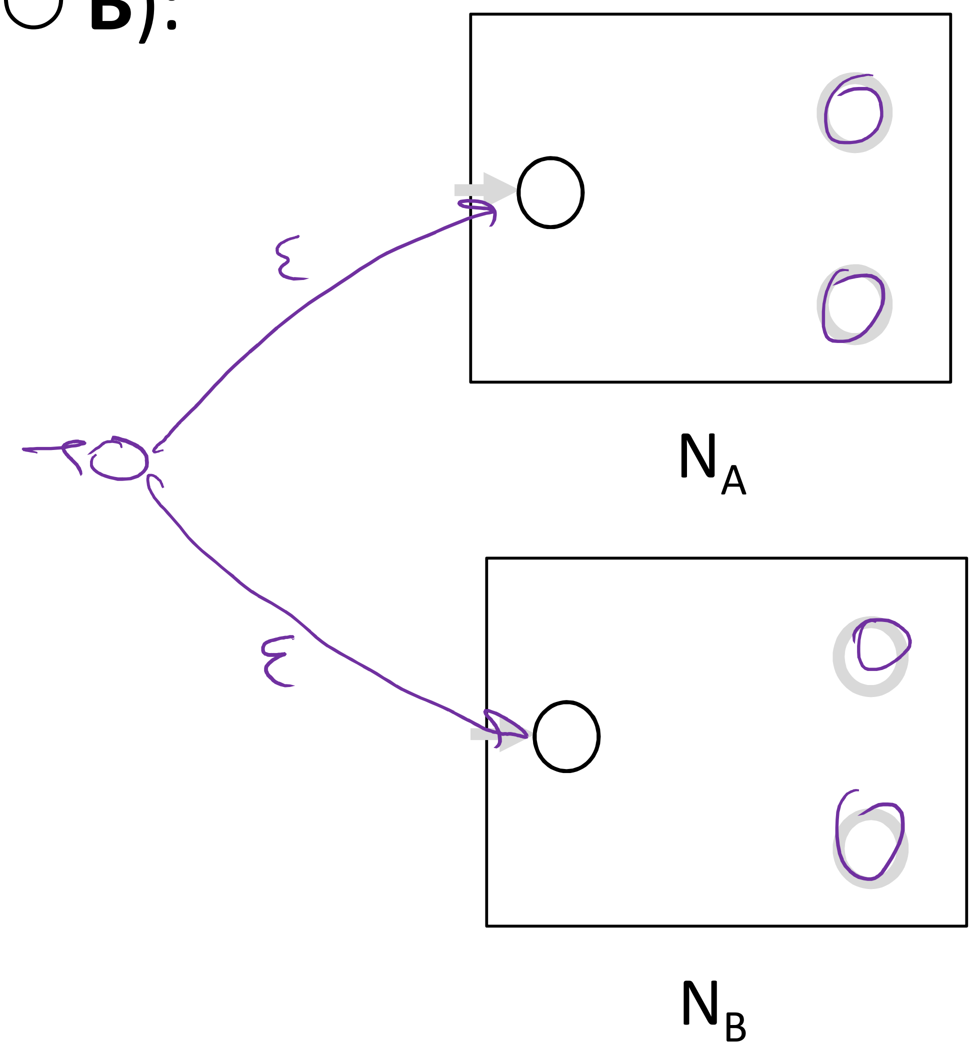
N_A



N_B

Inductive Step

Case (A ∪ B):



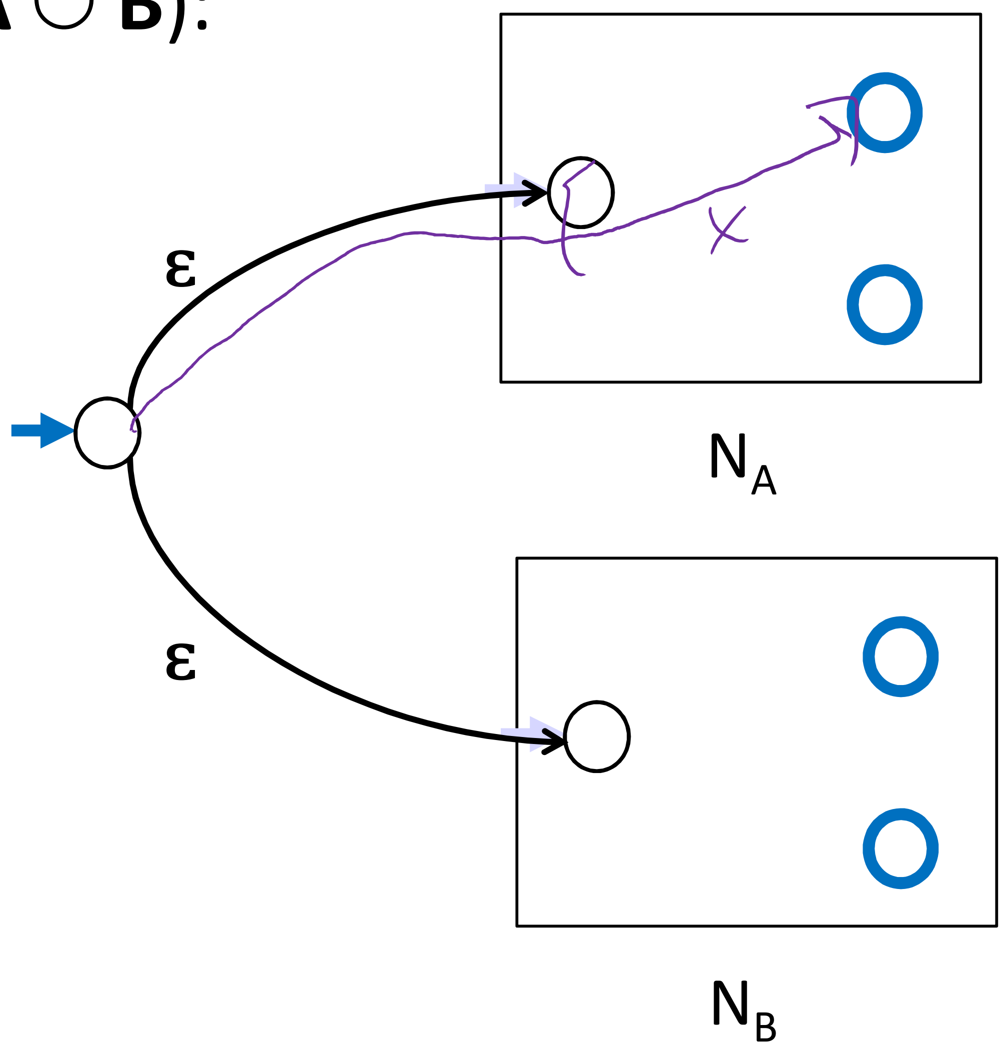
Inductive Step

Case (A ∪ B):

$x \in A \cup B$

Suppose $x \in A$

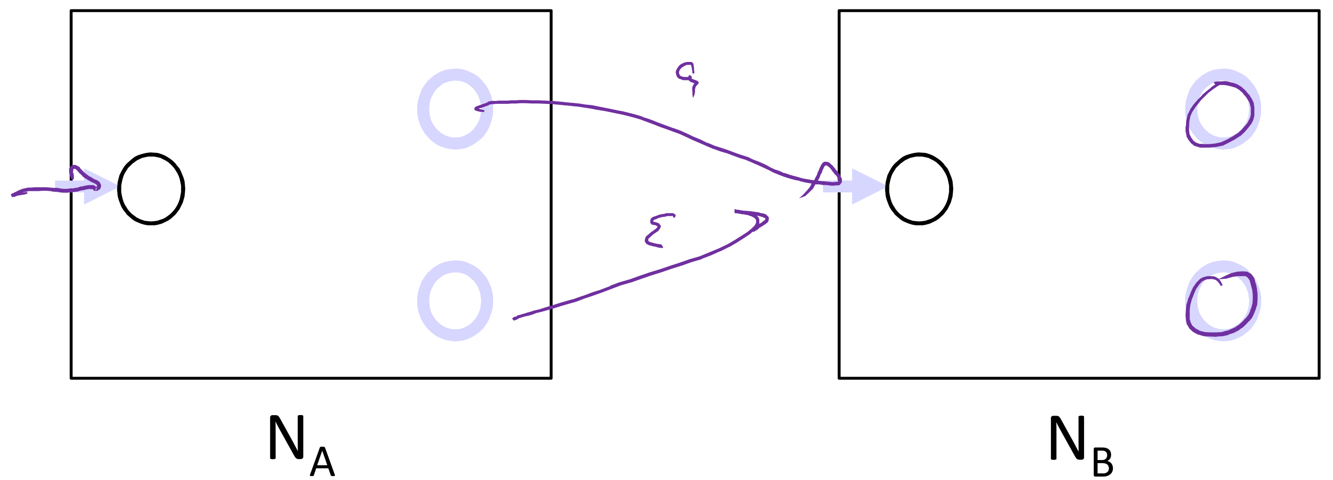
$\rightarrow$ x recognized by N.



$\boxed{N}$

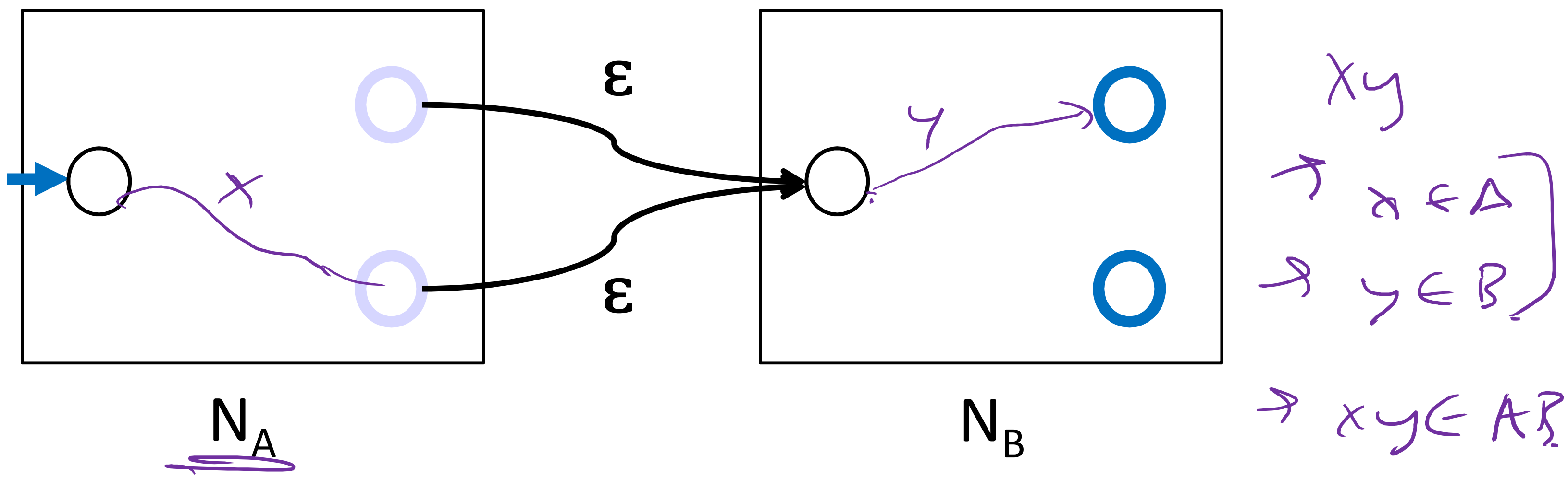
Inductive Step

Case (AB):



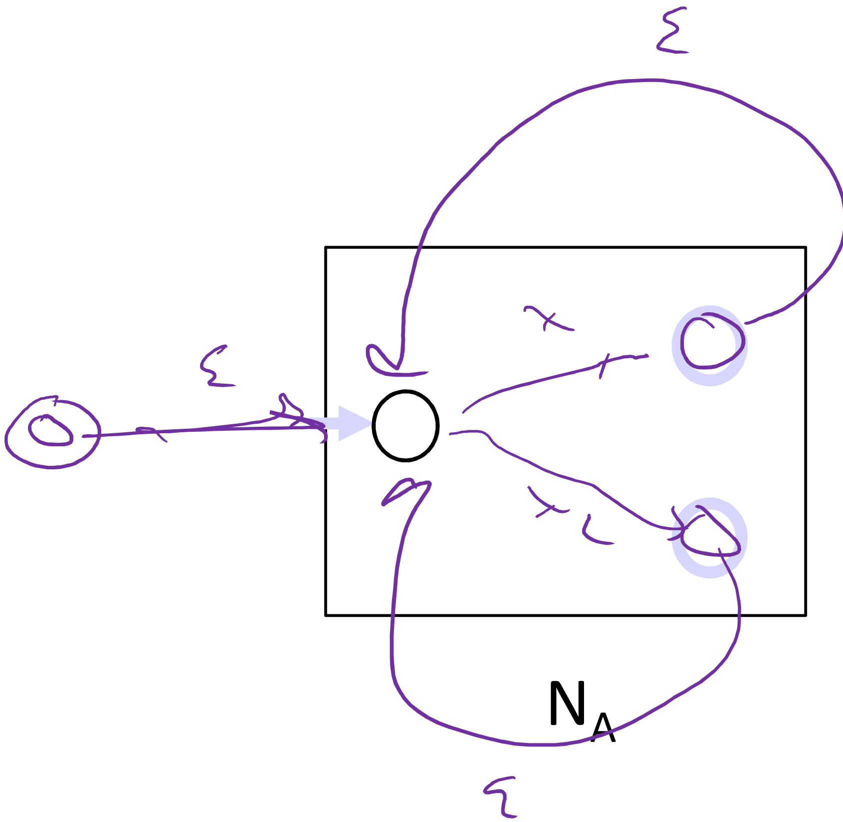
Inductive Step

Case (AB):



Inductive Step

Case A*



$x = x_1 x_2 \cdots x_k$

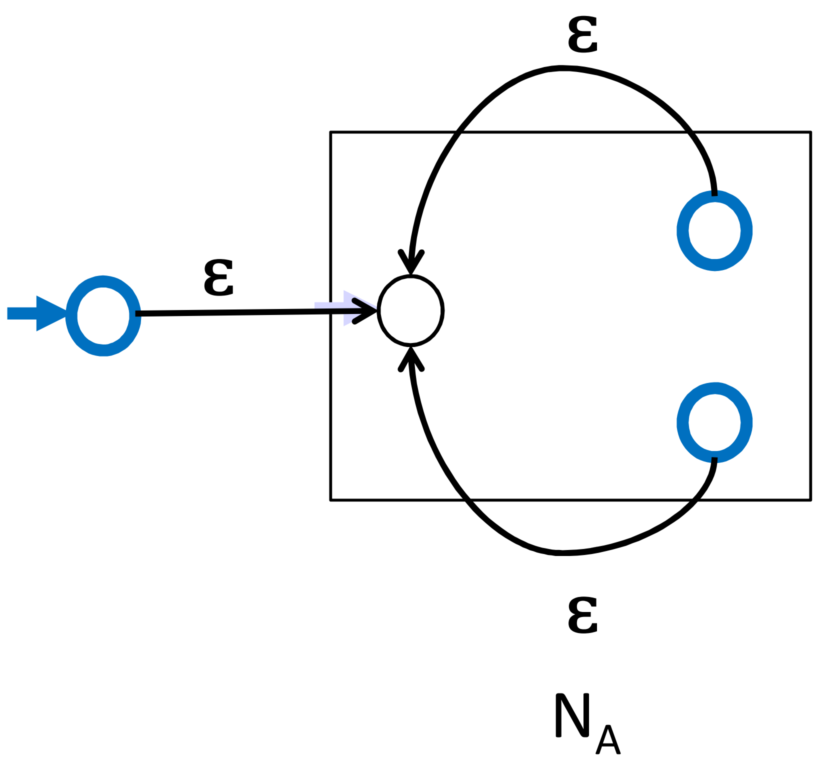
$\Rightarrow x_i \in A$

$\Rightarrow x \in A^*$.

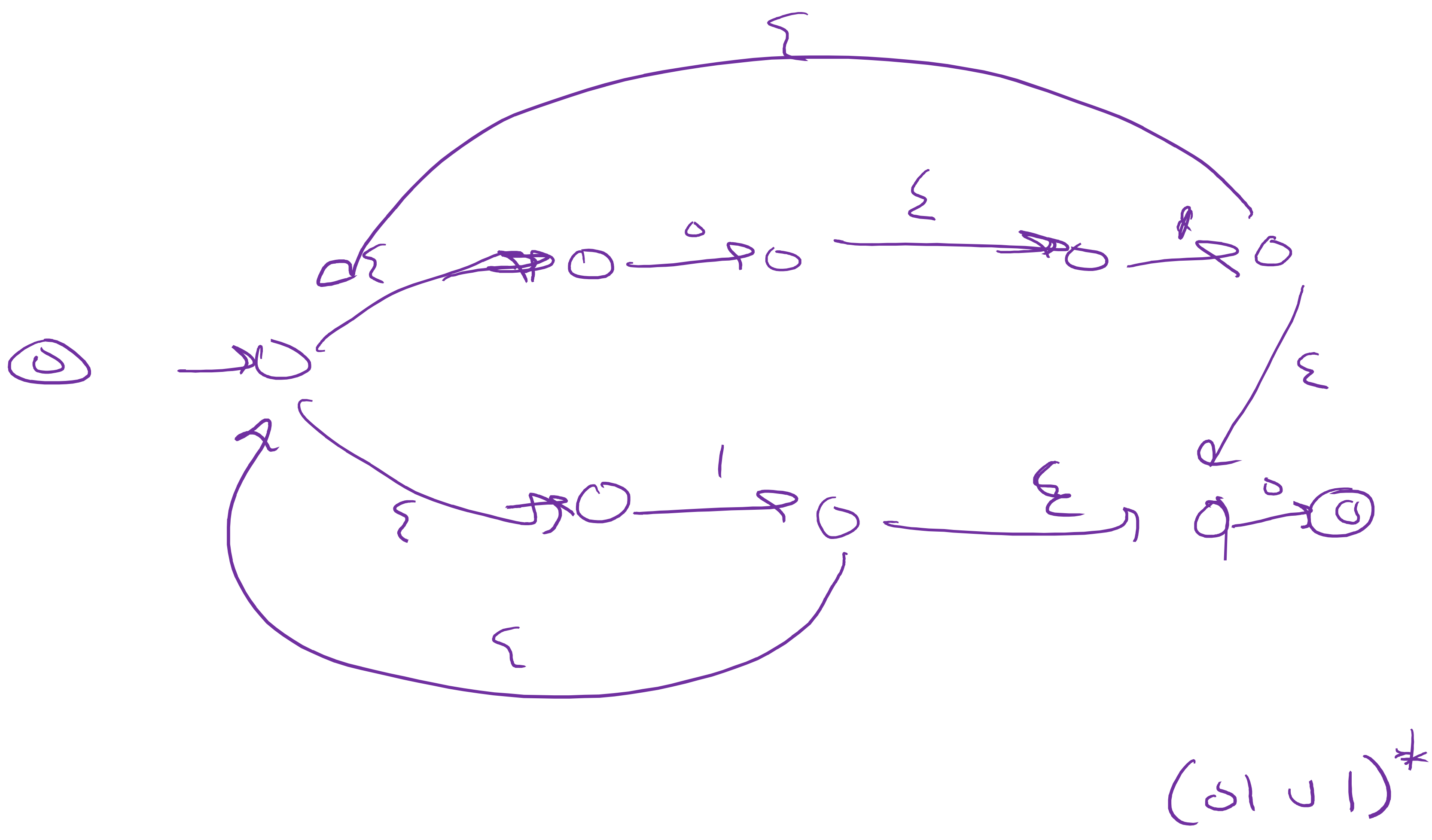
$\varepsilon \cup A$

Inductive Step

Case A*

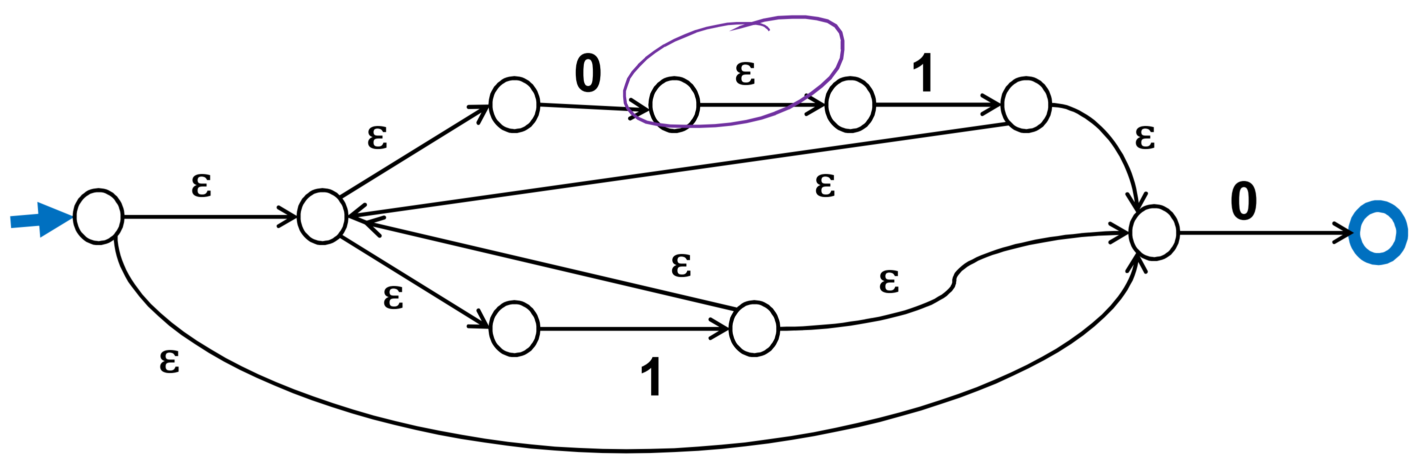


Build an NFA for (01 ∪1)*0



Solution

$(01 \cup 1)^*0$



NFAs and DFAs

$RE \subset NFA$

$DFA \subset NFA$.

Every DFA is an NFA

– DFAs have requirements that NFAs don't have

Can NFAs recognize more languages?

NFAs and DFAs

Every DFA is an NFA

– DFAs have requirements that NFAs don't have

Can NFAs recognize more languages? No!

Theorem: For every NFA there is a DFA that recognizes exactly the same language

Three ways of thinking about NFAs

- **Outside observer: Is there a path labeled by** x **from the start state to some final state?**

- **Perfect guesser: The NFA has input** x **and whenever there is a choice of what to do it magically guesses a good one (if one exists)**

- **Parallel exploration: The NFA computation runs all possible computations on** x **step-by-step at the same time in parallel**

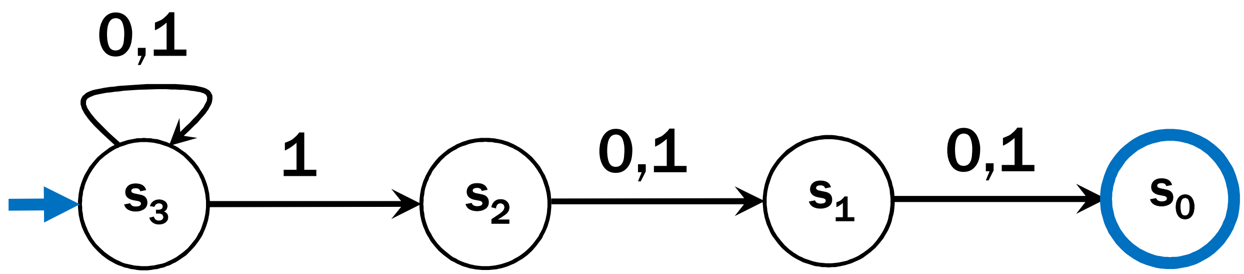
Conversion of NFAs to a DFAs

- Proof Idea:
  - The DFA keeps track of ALL the states that the part of the input string read so far can reach in the NFA
  - There will be one state in the DFA for each *subset* of states of the NFA that can be reached by some string

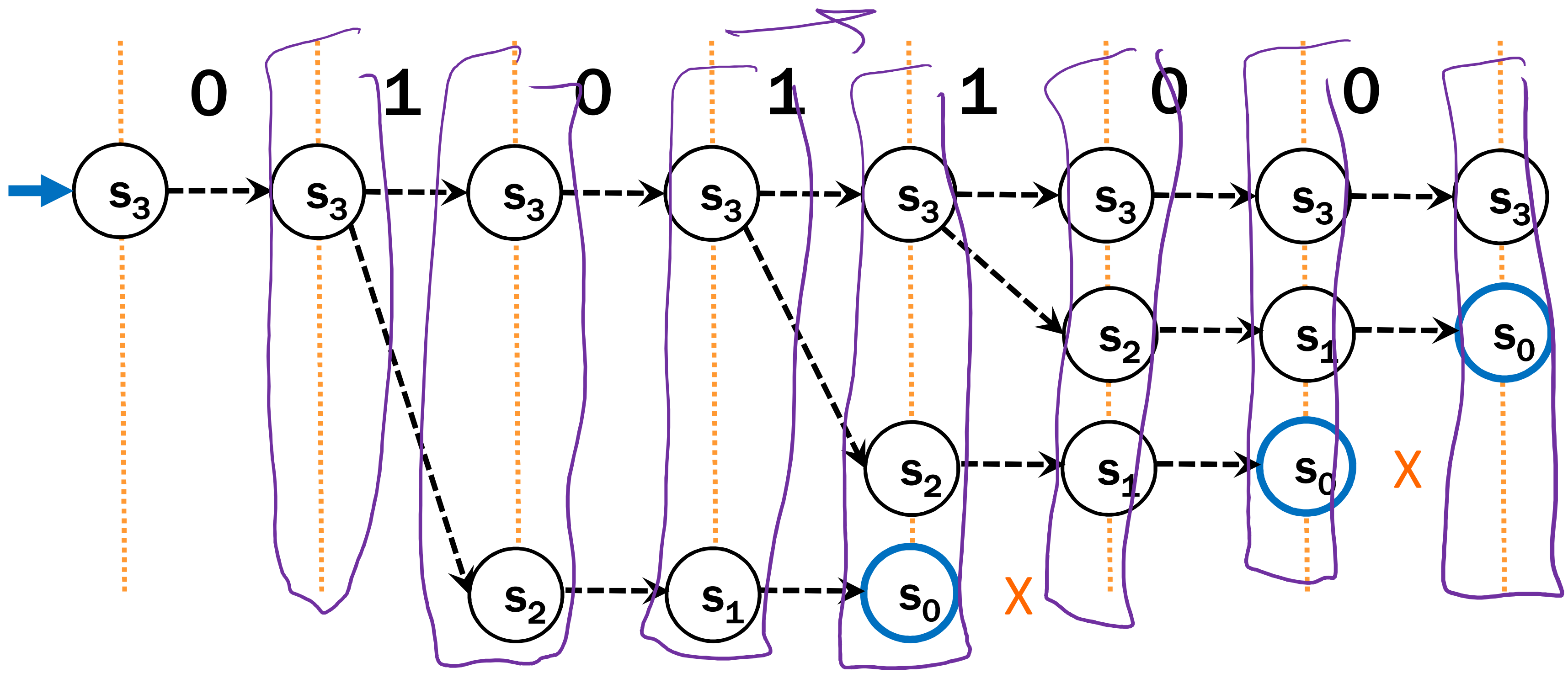
n

2^n

Parallel Exploration view of an NFA



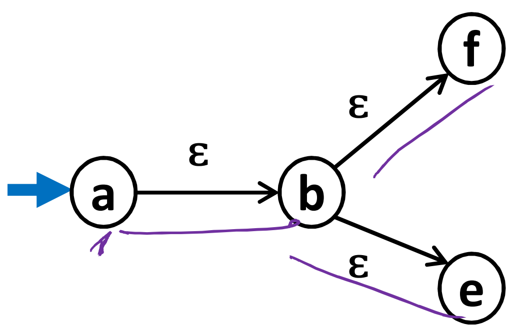
Input string 0101100



Conversion of NFAs to a DFAs

New start state for DFA

– The set of all states reachable from the start state of the NFA using only edges labeled ε



NFA

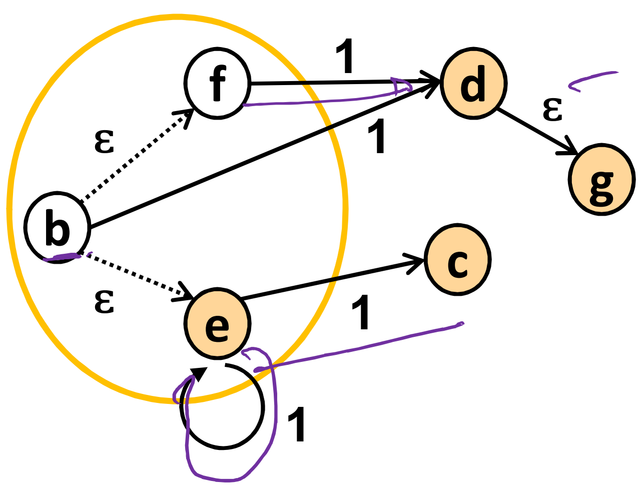
DFA

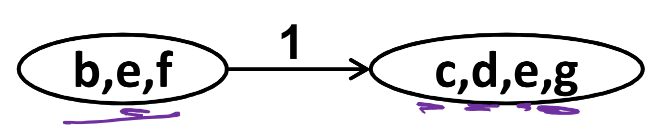
Conversion of NFAs to a DFAs

For each state of the DFA corresponding to a set S of states of the NFA and each symbol s

- **Add an edge labeled s to state corresponding to T, the set of states of the NFA reached by**
  - starting from some state in S, then
  - following one edge labeled by s, and

    then following some number of edges labeled by ε
- T **will be** ∅ **if no edges from** S **labeled** s **exist**

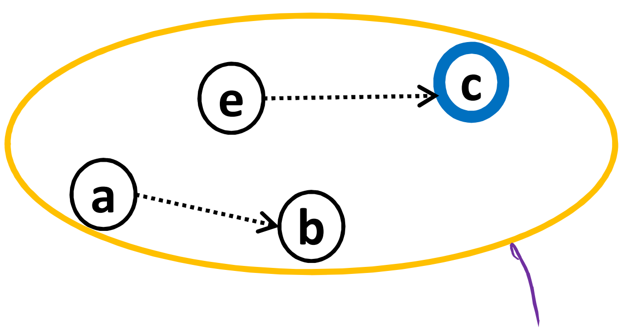




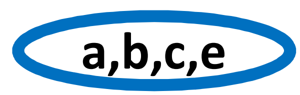
Conversion of NFAs to a DFAs

Final states for the DFA

– All states whose set contain some final state of the NFA

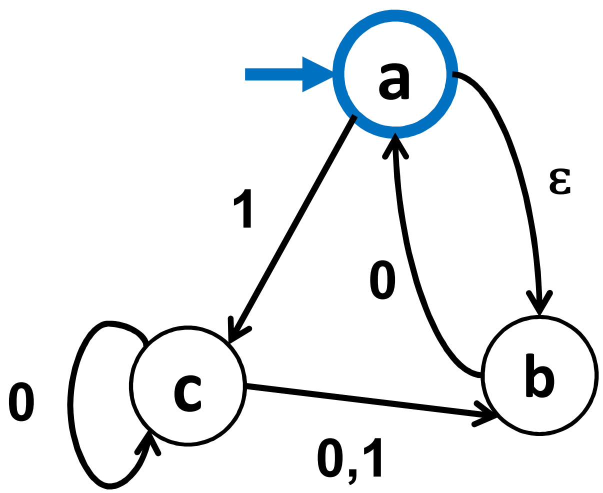


NFA

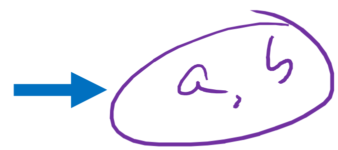


DFA

Example: NFA to DFA

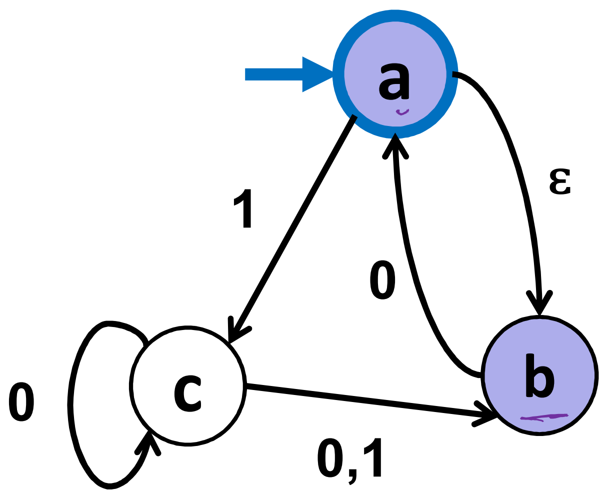


NFA

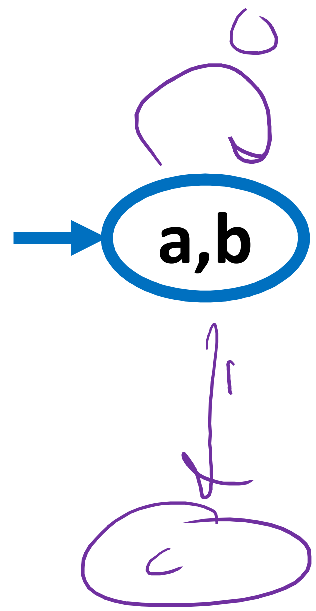


DFA

Example: NFA to DFA

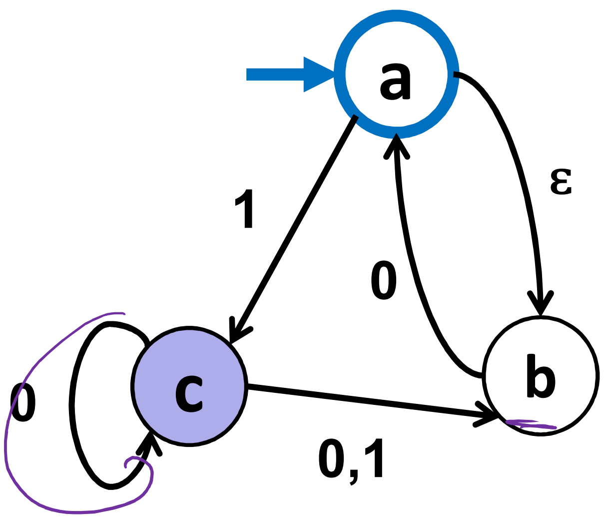


NFA

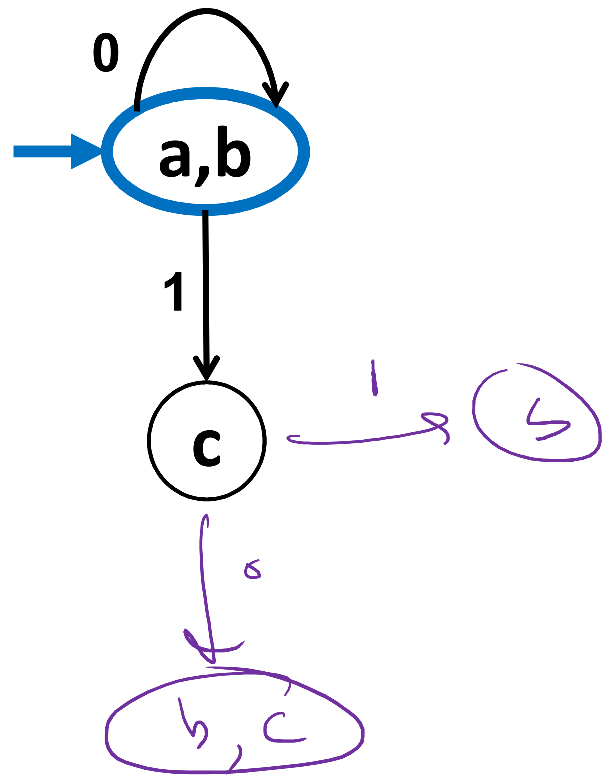


DFA

Example: NFA to DFA

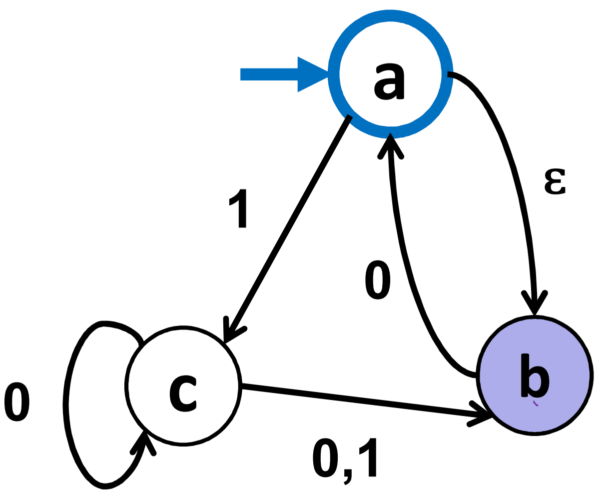


NFA

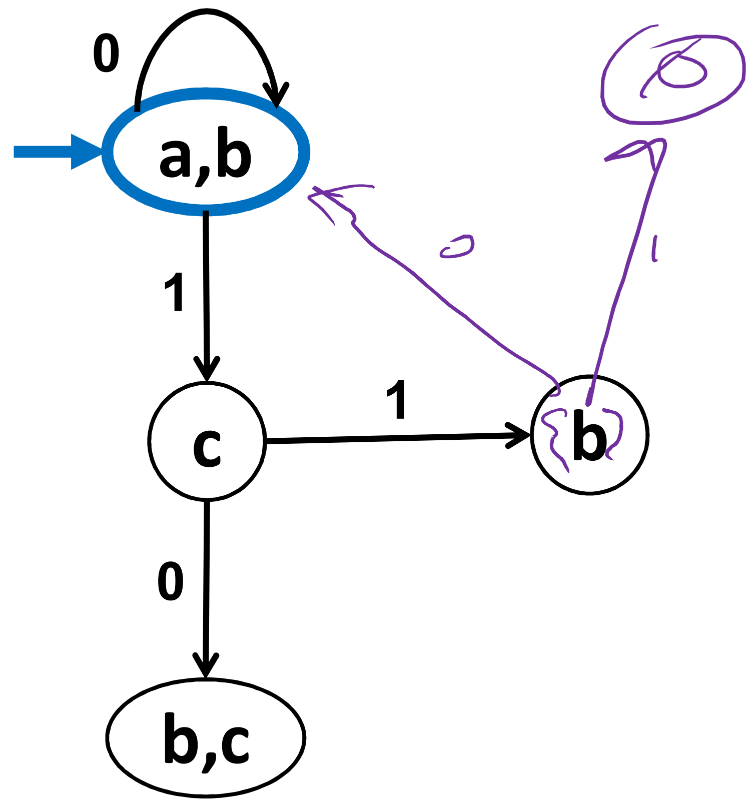


DFA

Example: NFA to DFA

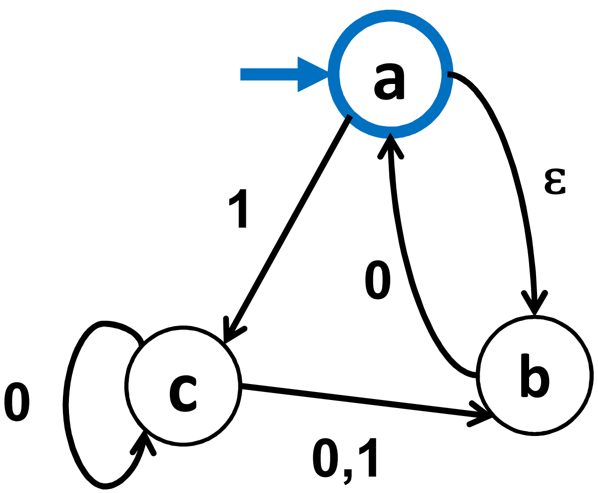


NFA

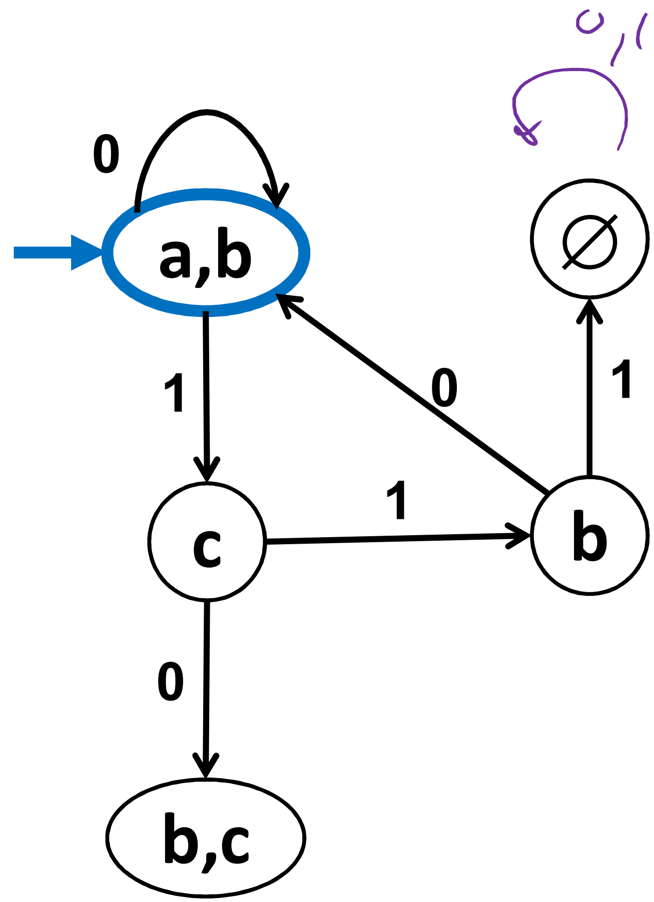


DFA

Example: NFA to DFA

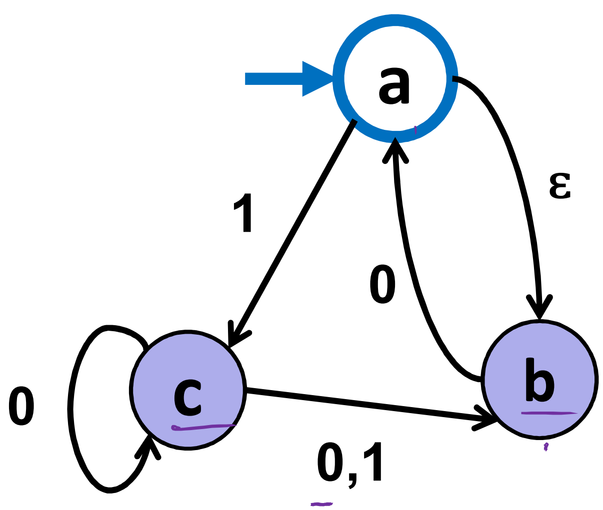


NFA

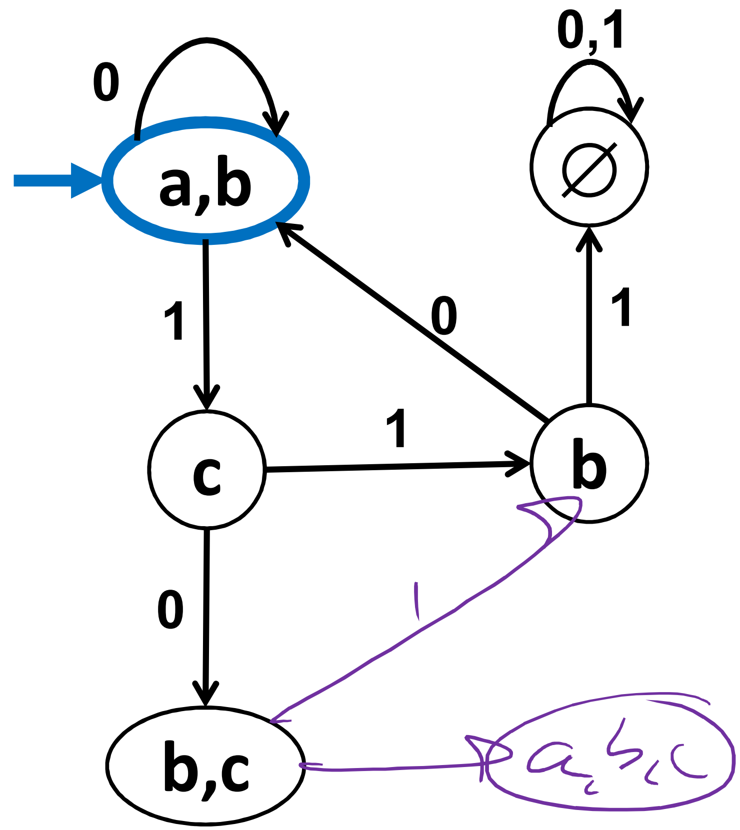


DFA

Example: NFA to DFA



NFA

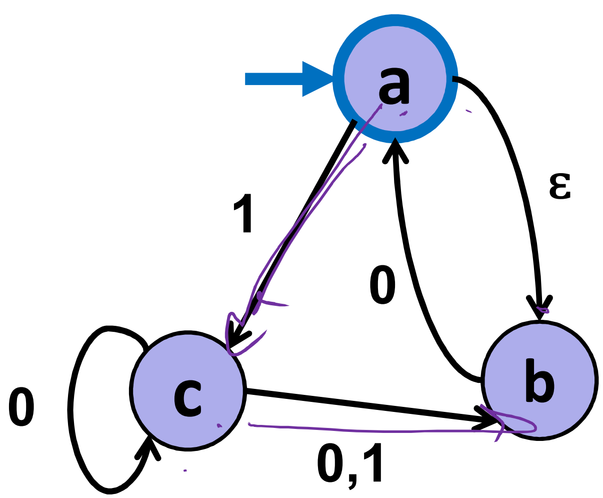


DFA

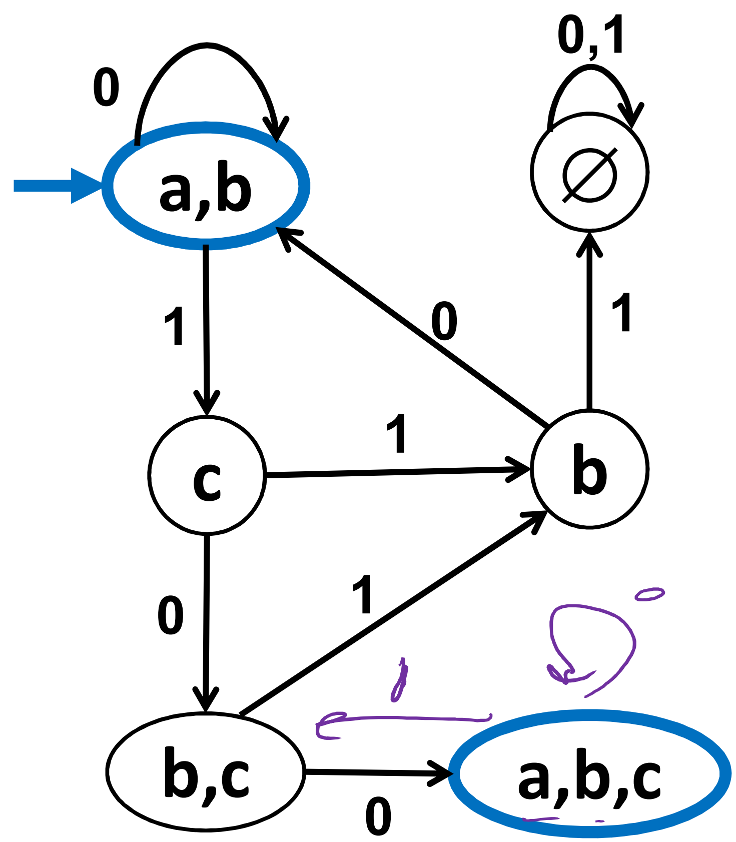
0 b, a, c

1 b

Example: NFA to DFA



NFA

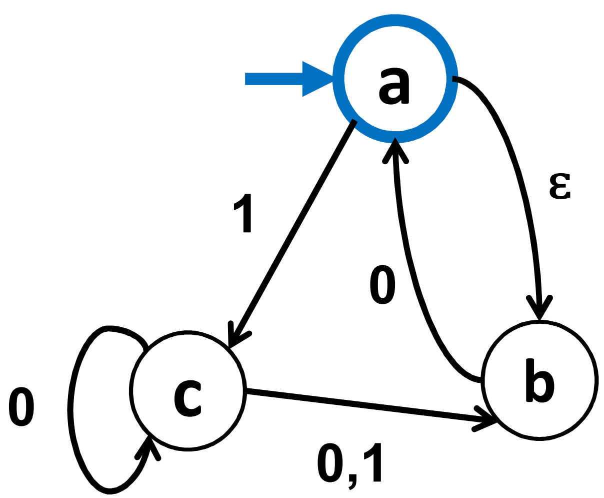


DFA

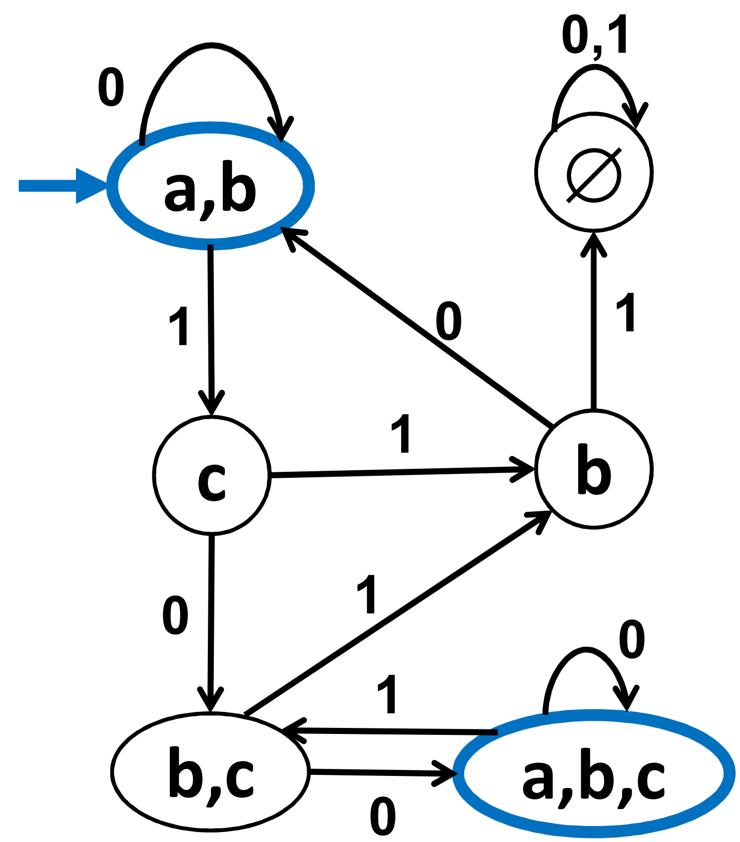
0 a,b,c

1 c,b

Example: NFA to DFA



NFA



DFA

Exponential Blow-up in Simulating Nondeterminism

- In general the DFA might need a state for every subset of states of the NFA
  - Power set of the set of states of the NFA
  - n-state NFA yields DFA with at most 2^n states
  - We saw an example where roughly 2^n is necessary

    "Is the n^{th} char from the end a 1?"

- The famous "$\boxed{\text{P=NP?}}$" question asks whether a similar blow-up is always necessary to get rid of nondeterminism for polynomial-time algorithms

DFAs ≡ NFAs ≡ Regular expressions

We have shown how to build an optimal DFA for every regular expression

– Build NFA
– Convert NFA to DFA using subset construction
– Minimize resulting DFA

Theorem: A language is recognized by a DFA (or NFA) if and only if it has a regular expression

You need to know this fact but we won't ask you anything about the "only if" direction from DFA/NFA to regular expression. For fun, we sketch the idea.

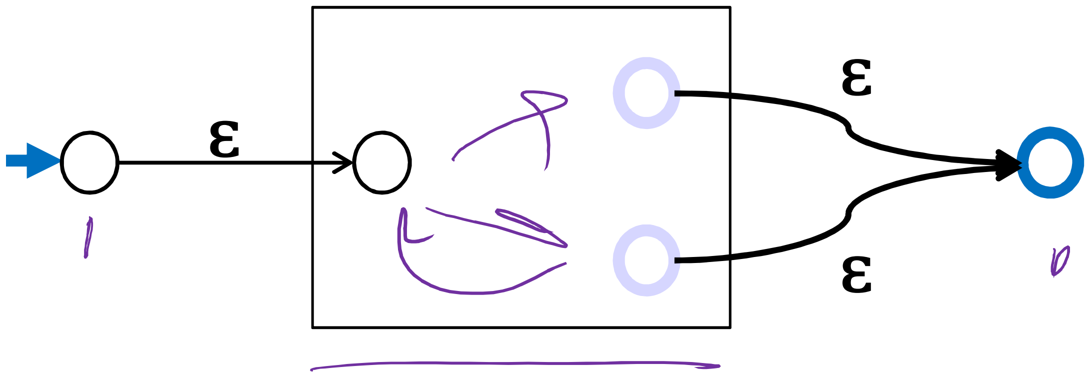
Generalized NFAs

- Like NFAs but allow
  - Regular Expressions as edge labels

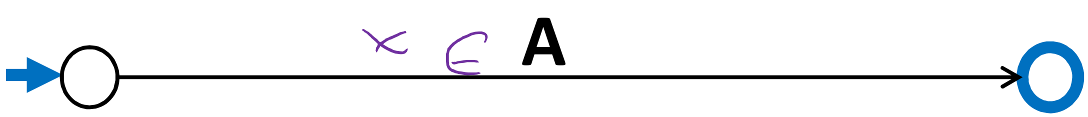
    NFAs already have edges labeled ε or ***a***
- An edge labeled by A can be followed by reading a string of input chars that is in the language represented by A
- Defn: A string x is accepted iff there is a *path* from start to final state *labeled by a regular expression* whose language contains x

Starting from an NFA

Add new start state and final state



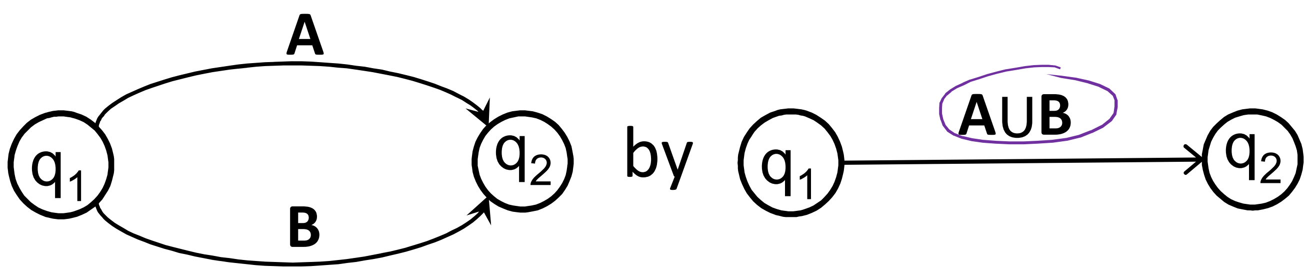
Then <u>eliminate original states</u> one by one, keeping the same language, until it looks like:



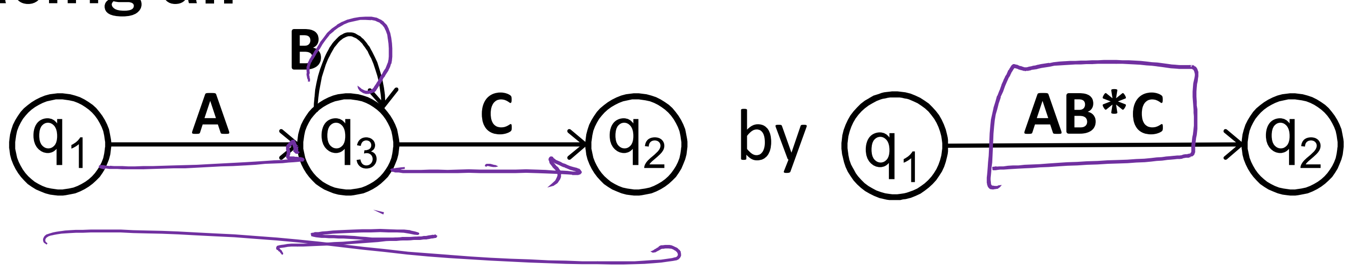
Final regular expression will be A

Only two simplification rules

- **Rule 1:** For any two states q_1 and q_2 with parallel edges (possibly $q_1=q_2$), replace



- **Rule 2:** Eliminate non-start/final state q_3 by replacing all

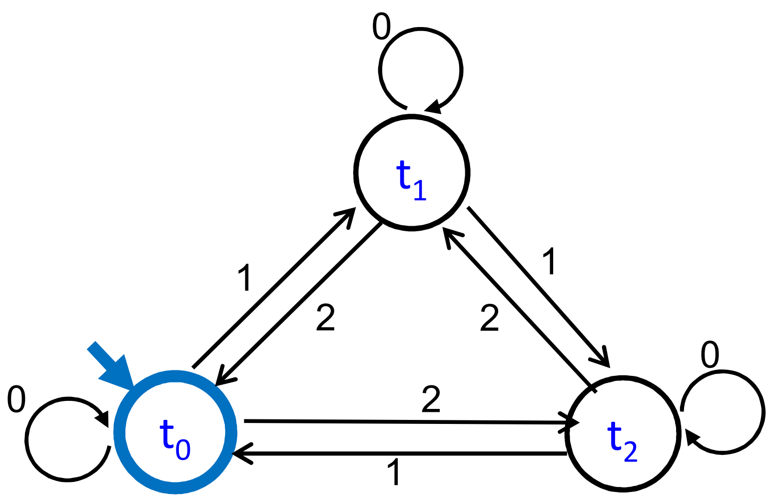


for *every* pair of states q_1, q_2 (even if $q_1=q_2$)

Converting an NFA to a regular expression

Consider the DFA for the mod 3 sum

- Accept strings from {0,1,2}* where the digits mod 3 sum of the digits is 0



Splicing out a state t_1

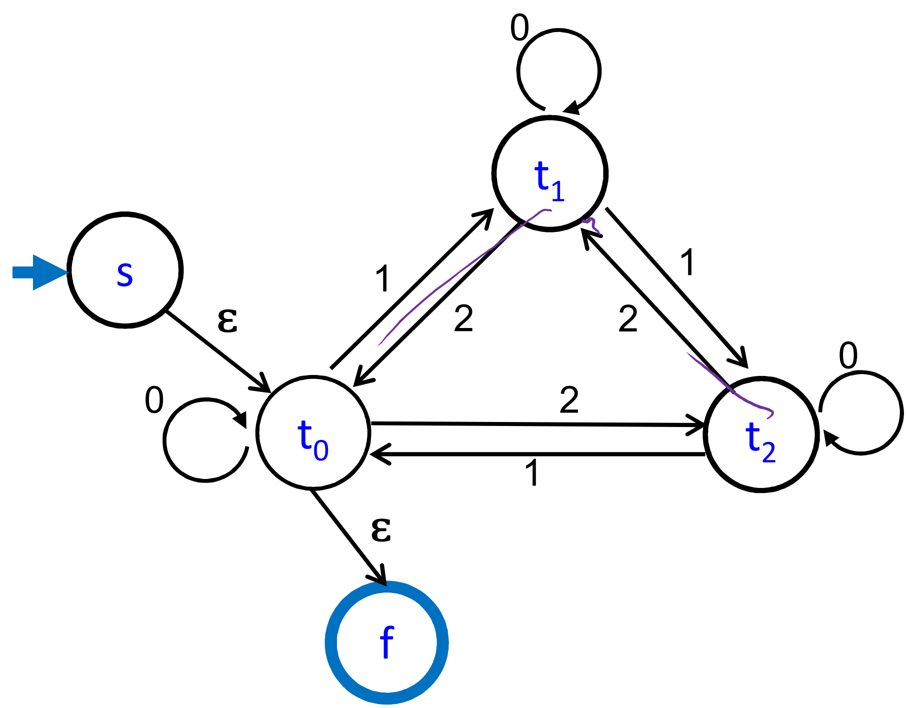
Regular expressions to add to edges

$t_0 \rightarrow t_1 \rightarrow t_0$: 10*2

$t_0 \rightarrow t_1 \rightarrow t_2$: 10*1

$t_2 \rightarrow t_1 \rightarrow t_0$: 20*2

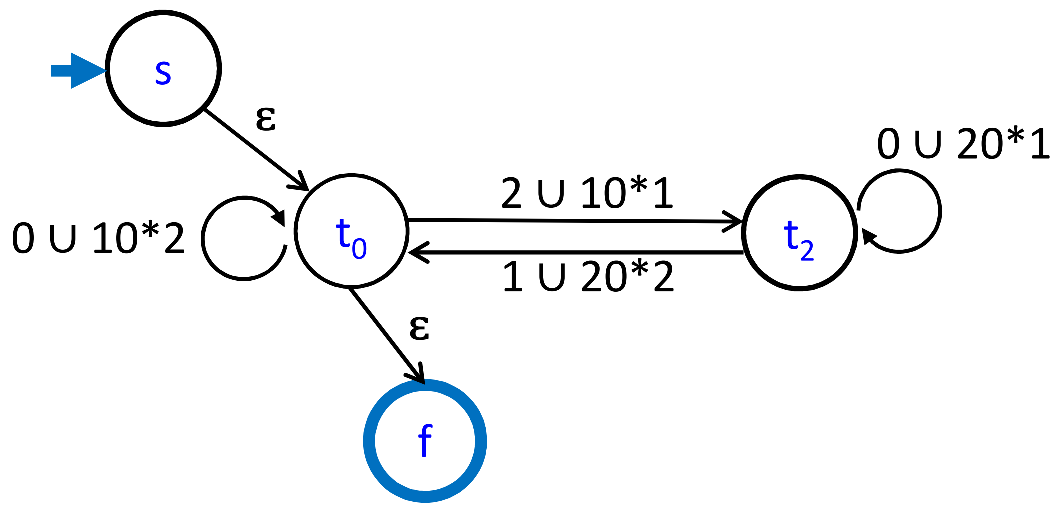
$t_2 \rightarrow t_1 \rightarrow t_2$: 20*1



Splicing out a state t_1

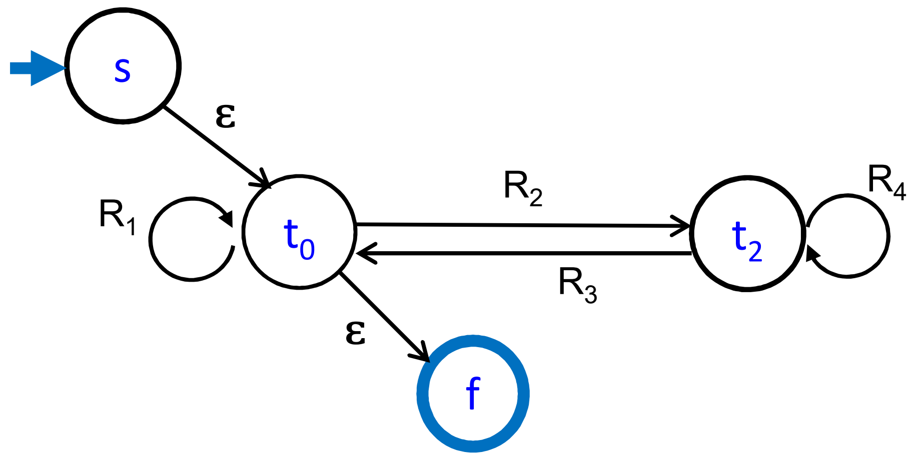
Regular expressions to add to edges

$t_0 \rightarrow t_1 \rightarrow t_0$: 10*2
$t_0 \rightarrow t_1 \rightarrow t_2$: 10*1
$t_2 \rightarrow t_1 \rightarrow t_0$: 20*2
$t_2 \rightarrow t_1 \rightarrow t_2$: 20*1

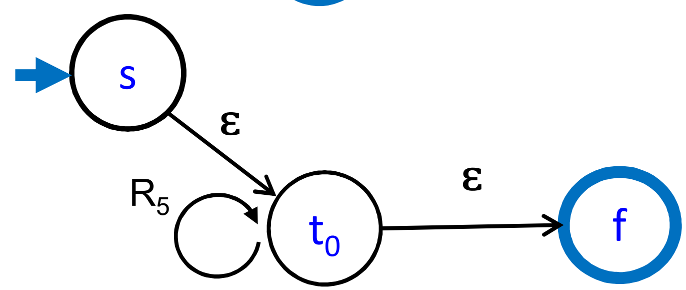


Splicing out state t_2 (and then t_0)

R_1: $0 \cup 10^*2$
R_2: $2 \cup 10^*1$
R_3: $1 \cup 20^*2$
R_4: $0 \cup 20^*1$



R_5: $R_1 \cup R_2R_4^*R_3$



Final regular expression: R_5^*=
$(0 \cup 10^*2 \cup (2 \cup 10^*1)(0 \cup 20^*1)^*(1 \cup 20^*2))^*$