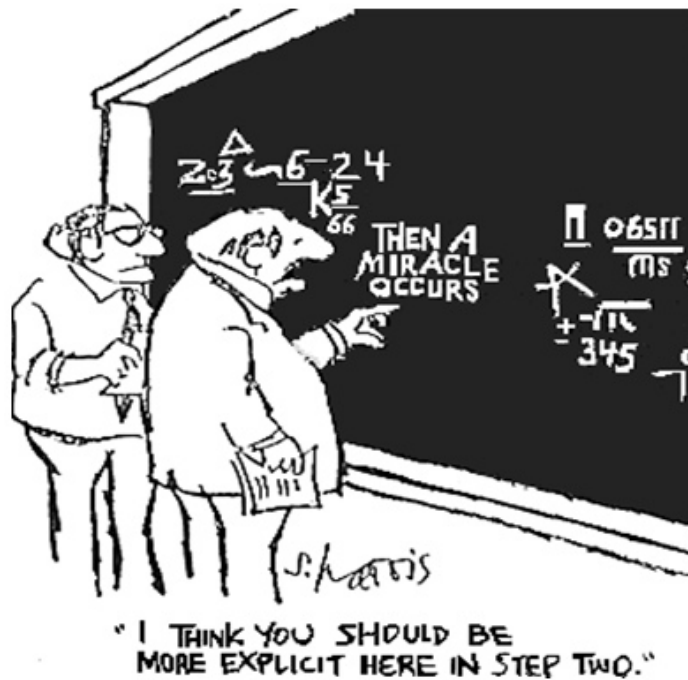


# CSE 311: Foundations of Computing

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## Lecture 23: Finite State Machine Minimization & NFAs



# State Minimization

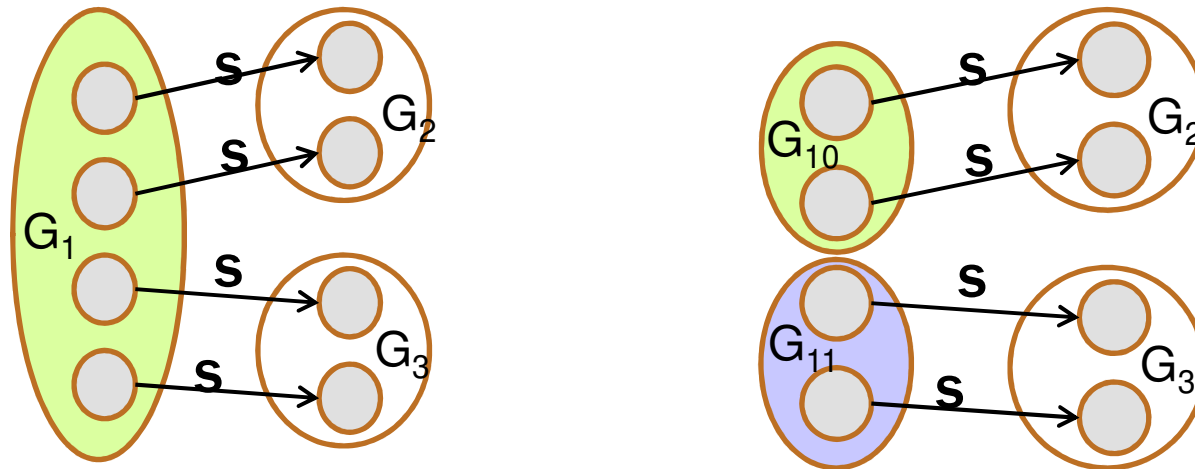
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- Many different FSMs (DFAs) for the same problem
- Take a given FSM and try to reduce its state set by combining states
  - Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this

# State Minimization Algorithm

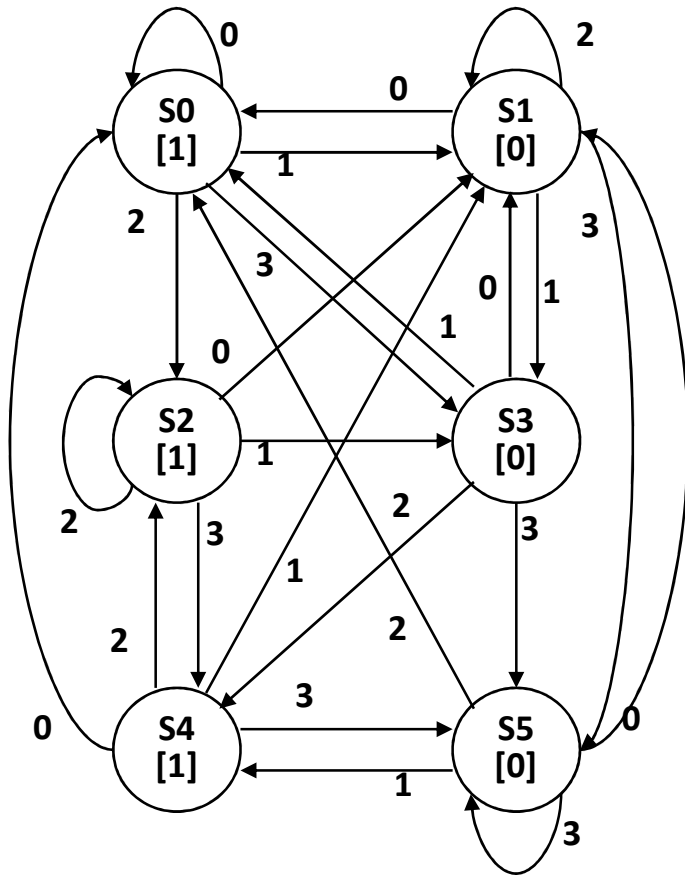
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1. Put states into groups based on their outputs (or whether they are final states or not)
2. Repeat the following until no change happens
  - a. If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** into smaller groups based on which group the states go to on **s**



3. Finally, convert groups to states

# State Minimization Example

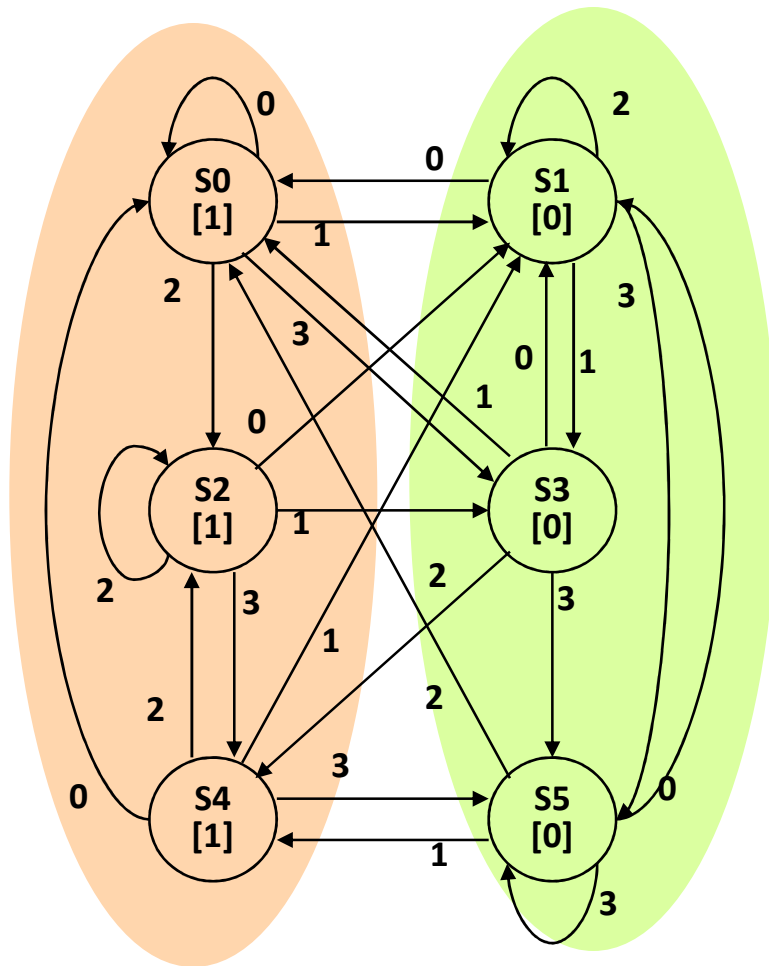


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

# State Minimization Example

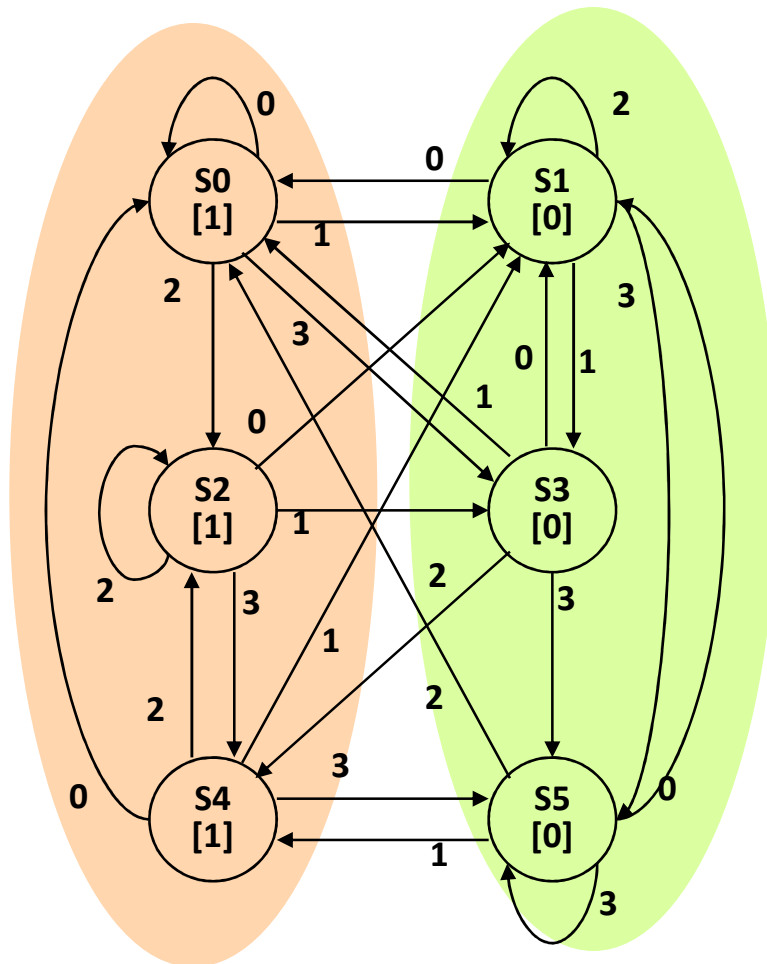


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

# State Minimization Example



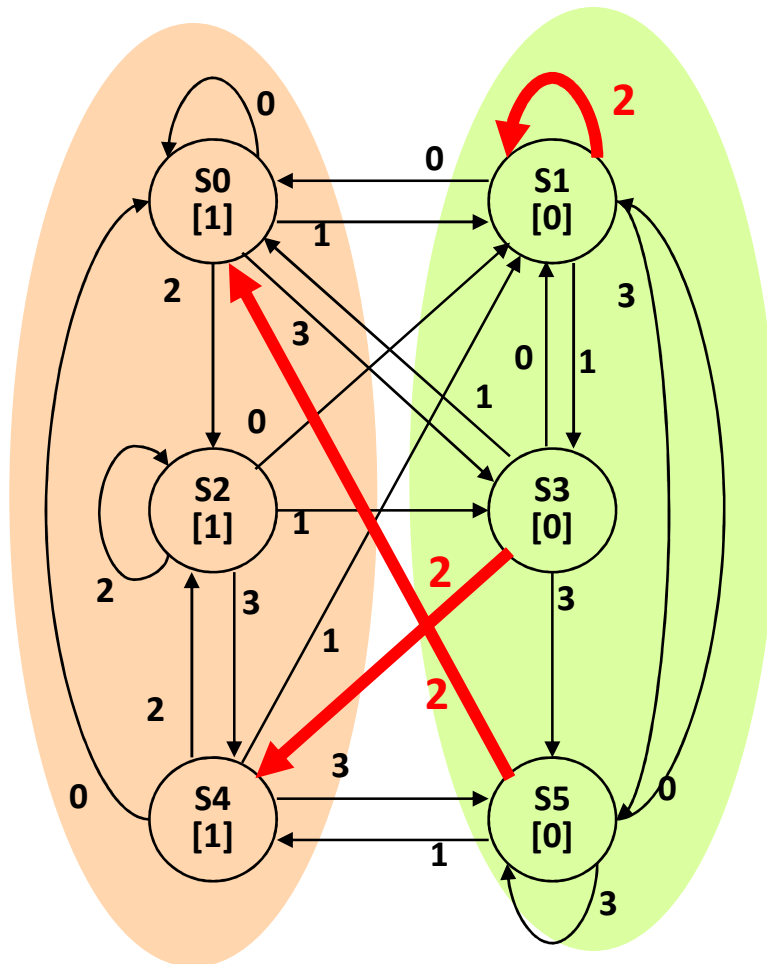
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** based on which group the states go to on **s**

# State Minimization Example



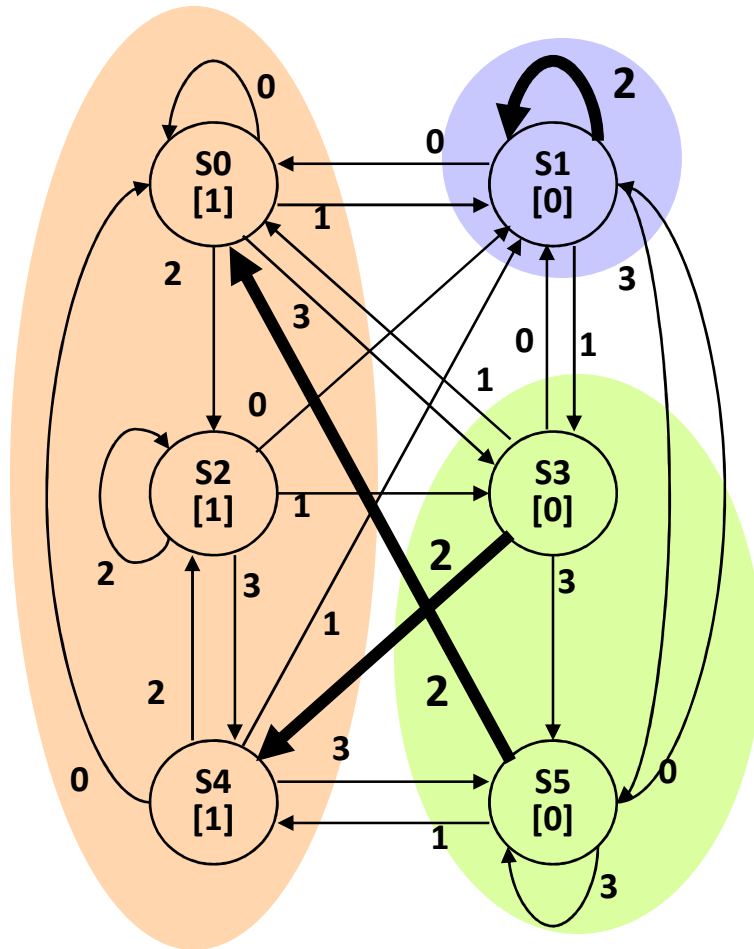
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** based on which group the states go to on **s**

# State Minimization Example



present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

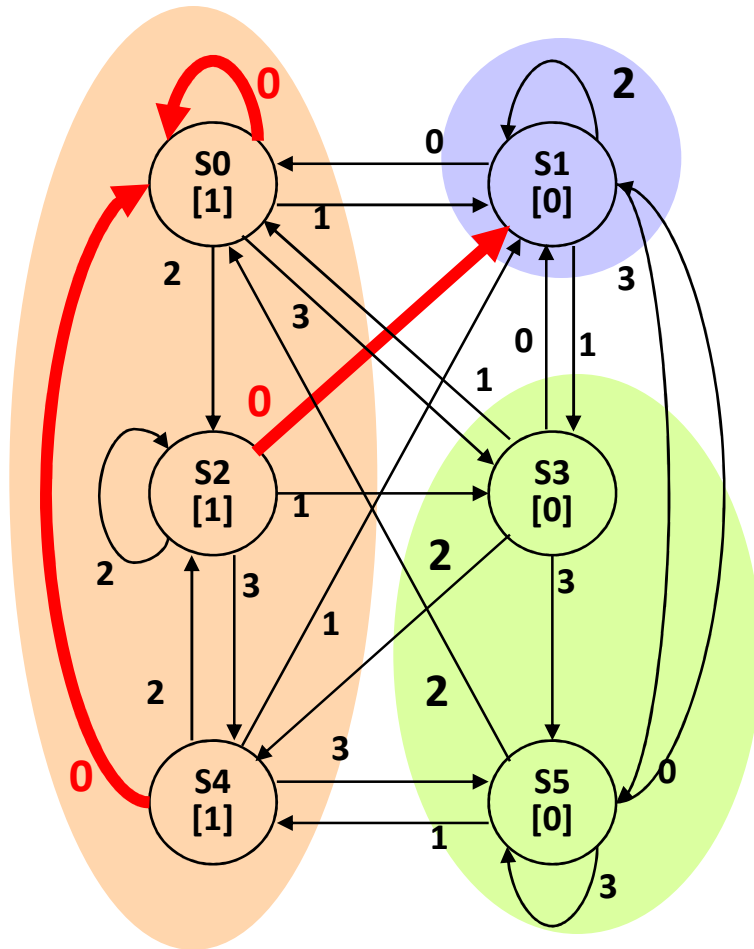
state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** based on which group the states go to on **s**



# State Minimization Example



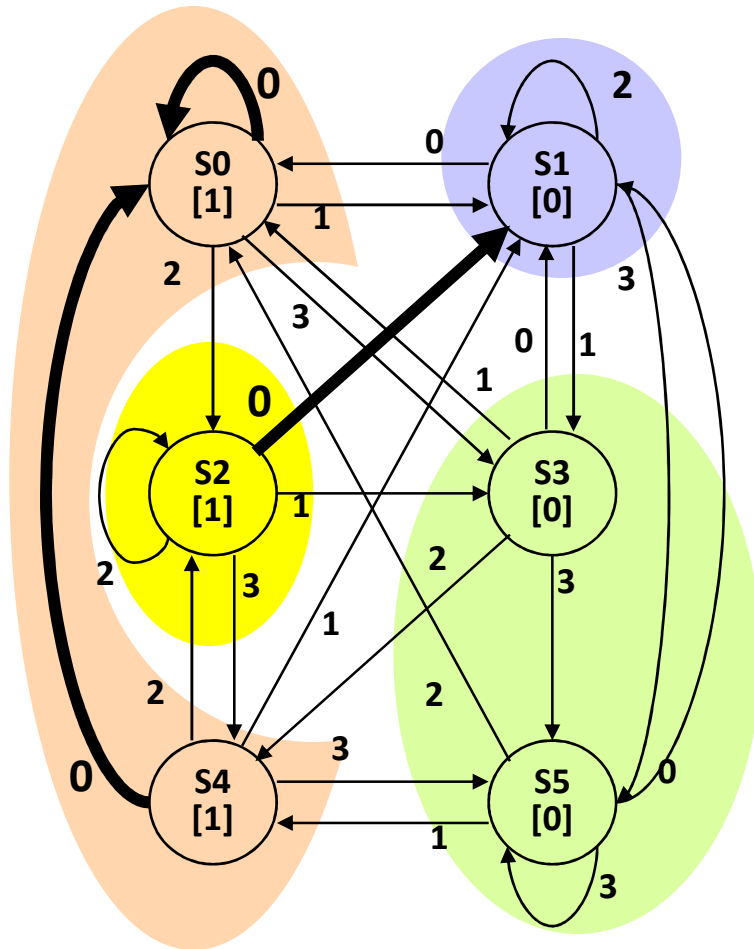
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** based on which group the states go to on **s**

# State Minimization Example



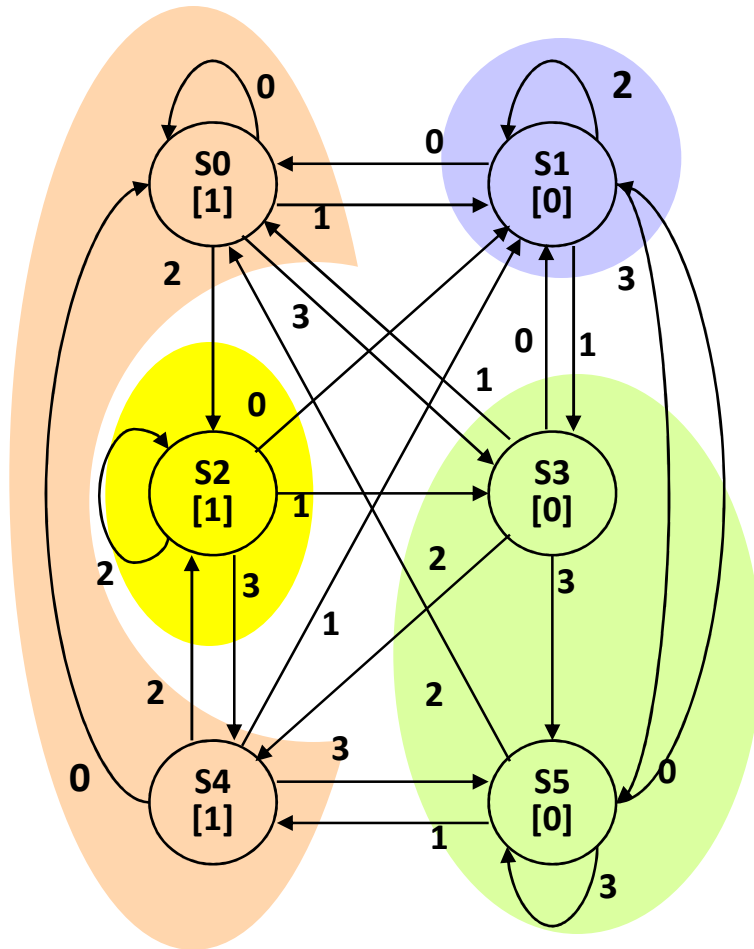
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group **G** agree on which group **s** leads to, split **G** based on which group the states go to on **s**

# State Minimization Example



present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

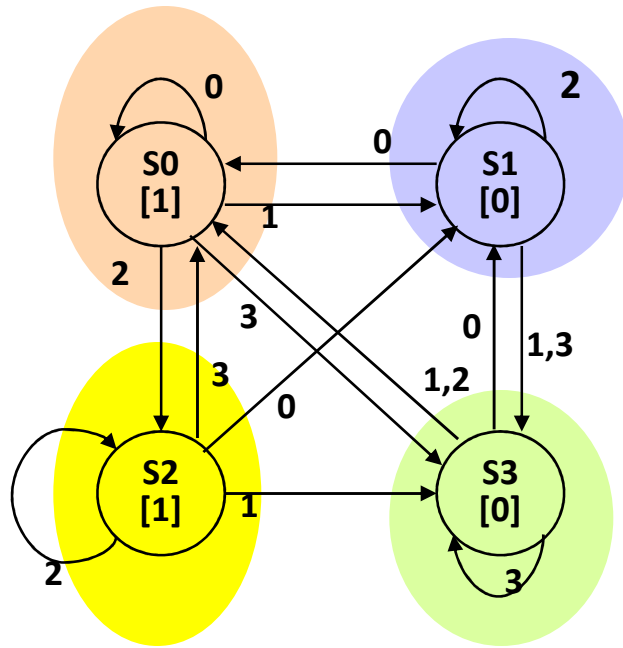
state transition table

Finally convert groups to states:

Can combine states S0-S4 and S3-S5.

In table replace all S4 with S0 and all S5 with S3

# Minimized Machine

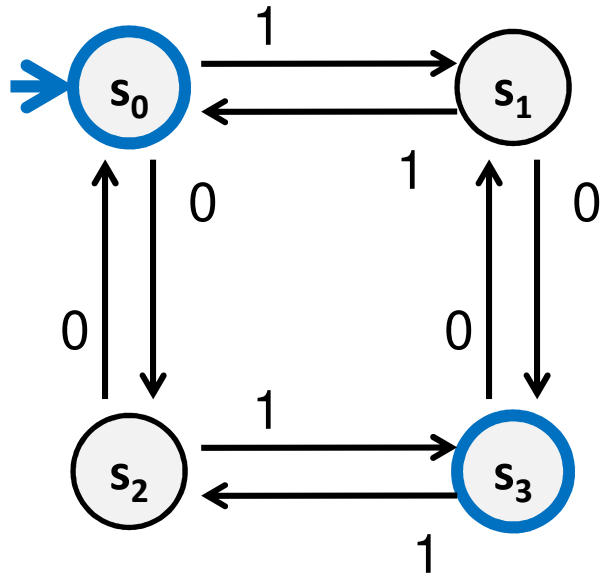


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S3	0
S2	S1	S3	S2	S0	1
S3	S1	S0	S0	S3	0

state transition table

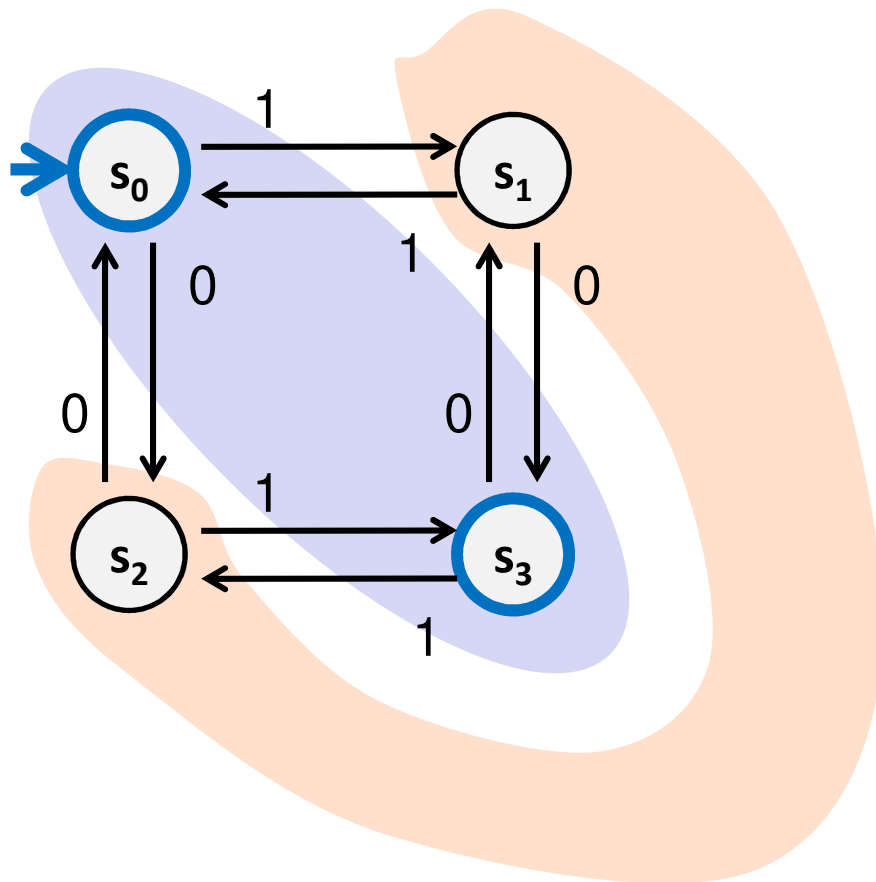
# A Simpler Minimization Example

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# A Simpler Minimization Example

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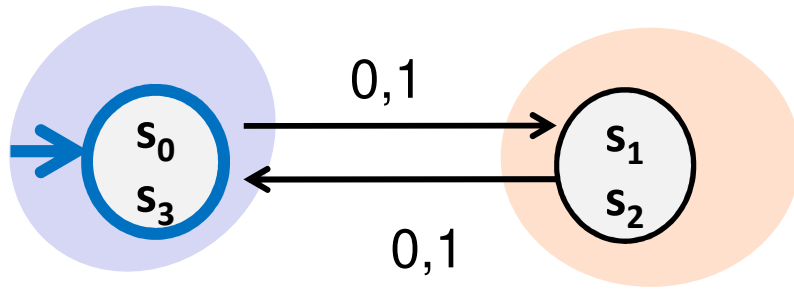


**Split states into  
final/non-final groups**

**Every symbol causes  
the DFA to go from one  
group to the other so  
neither group needs to  
be split**

# Minimized DFA

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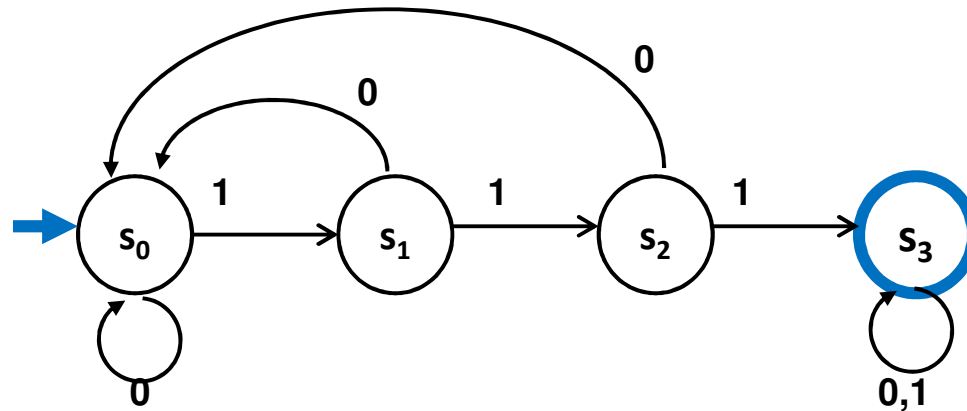


# Another way to look at DFAs

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Definition: The label of a path in a DFA is the concatenation of all the labels on its edges in order

Lemma:  $x$  is in the language recognized by a DFA iff  $x$  labels a path from the start state to some final state

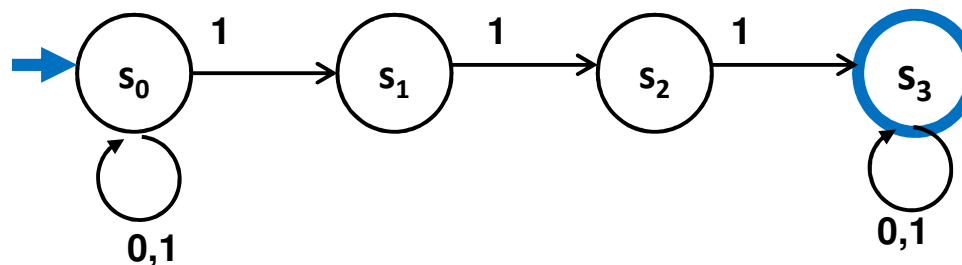




# Nondeterministic Finite Automata (NFA)

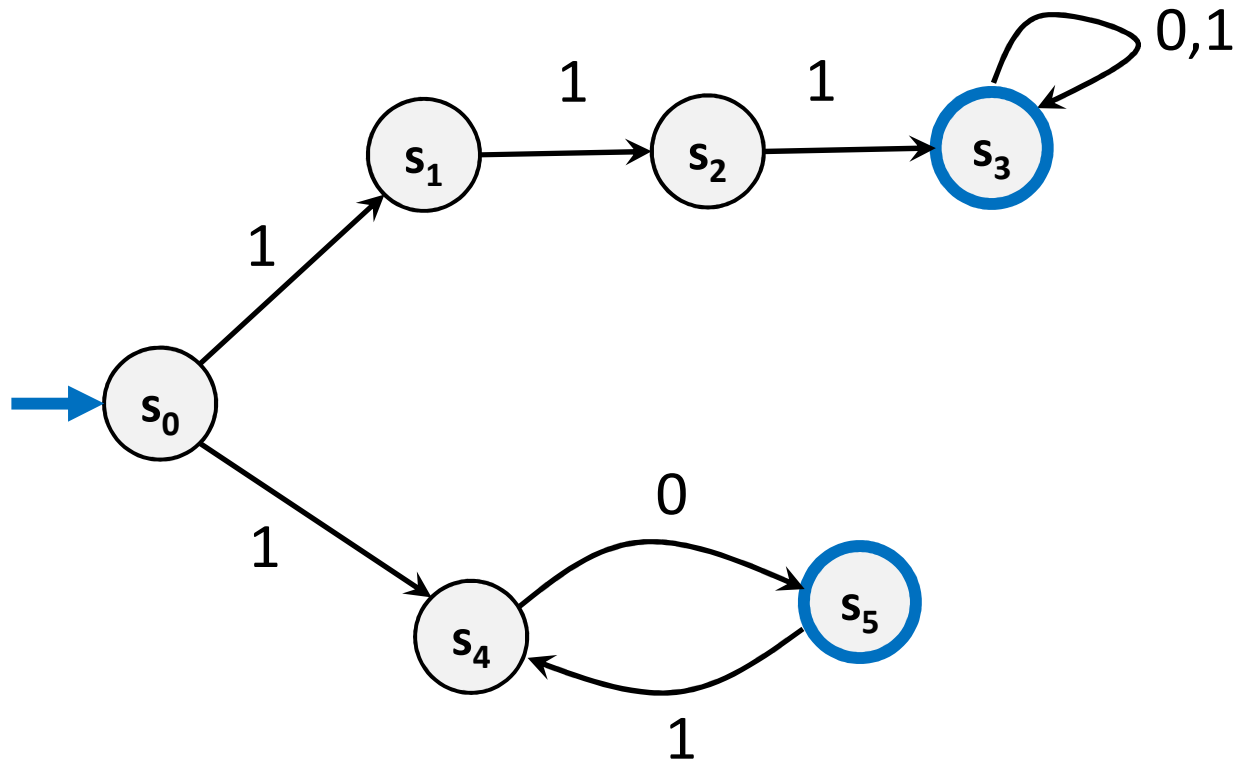
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- Graph with start state, final states, edges labeled by symbols (like DFA) but
  - Not required to have exactly 1 edge out of each state labeled by each symbol— can have 0 or  $>1$
  - Also can have edges labeled by empty string  $\epsilon$
- **Definition:**  $x$  is in the language recognized by an NFA if and only if  $x$  labels a path from the start state to some final state



## Consider This NFA

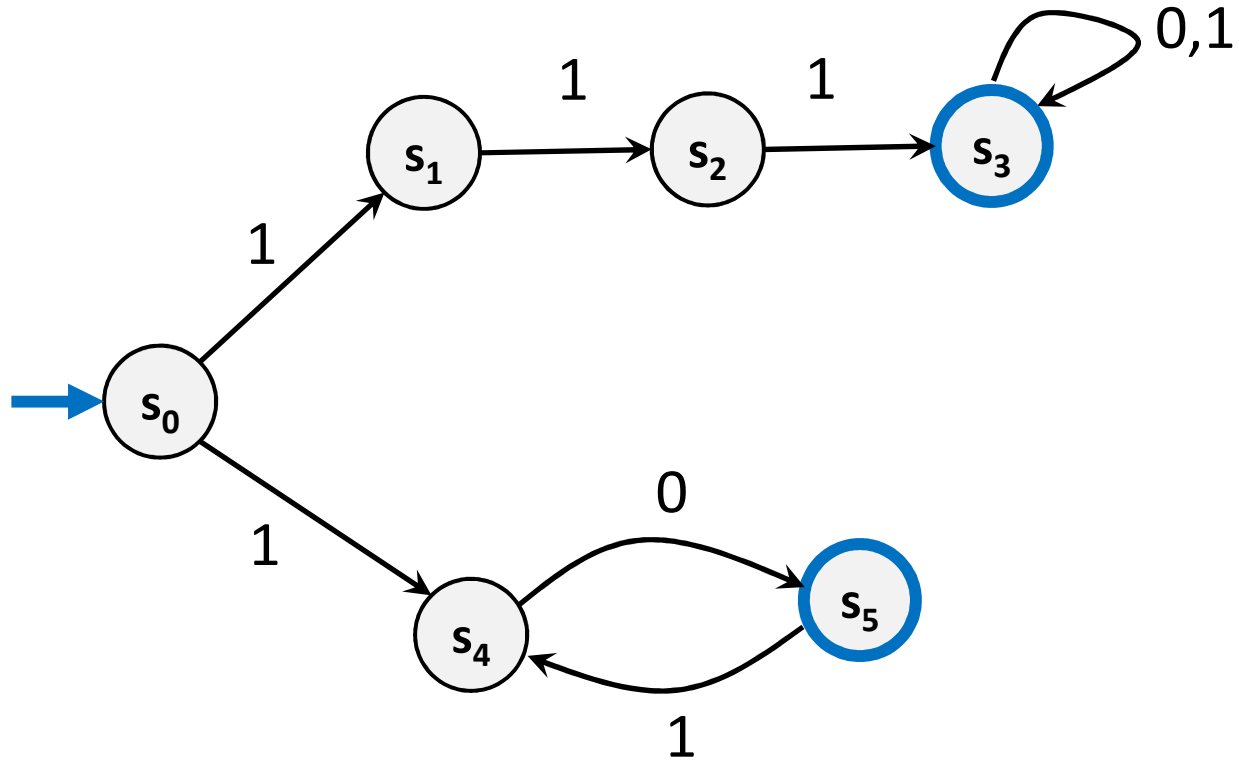
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What language does this NFA accept?

## Consider This NFA

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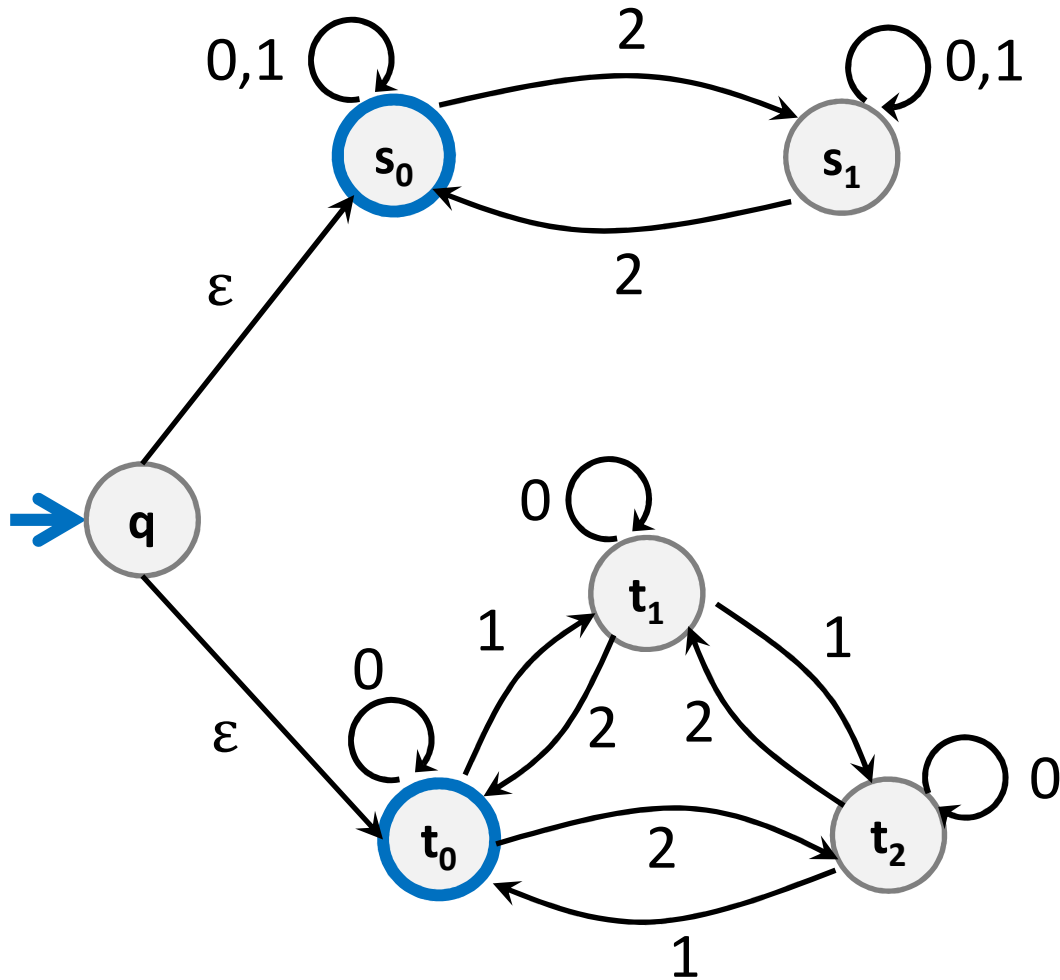


What language does this NFA accept?

$$10(10)^* \cup 111(0 \cup 1)^*$$

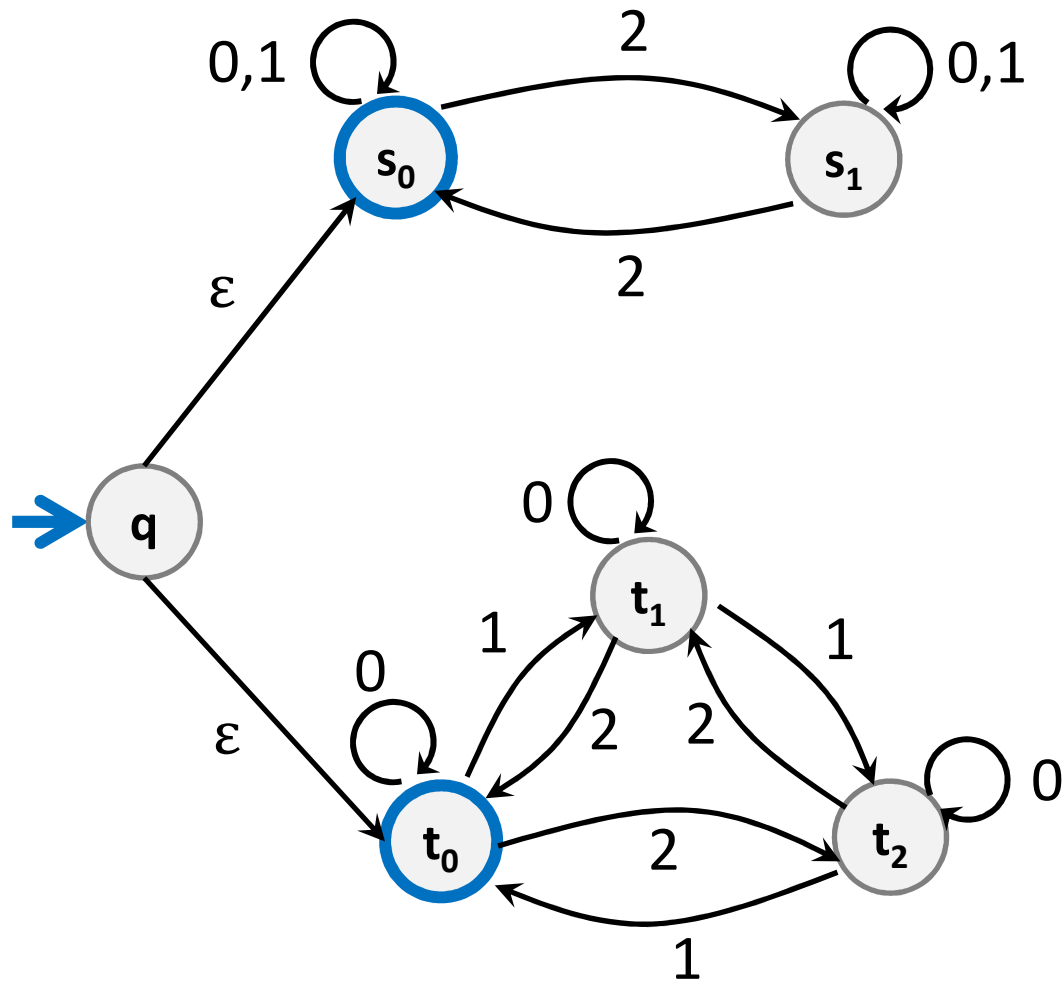
# NFA $\epsilon$ -moves

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# NFA $\epsilon$ -moves

Strings over  $\{0,1,2\}$  w/even # of 2's OR sum to 0 mod 3



# Three ways of thinking about NFAs

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- **Outside observer:** Is there a path labeled by  $x$  from the start state to some final state?
- **Perfect guesser:** The NFA has input  $x$  and whenever there is a choice of what to do it magically guesses a good one (if one exists)
- **Parallel exploration:** The NFA computation runs all possible computations on  $x$  step-by-step at the same time in parallel

**NFA for set of binary strings with a 1 in the 3<sup>rd</sup> position from the end**

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# NFA for set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

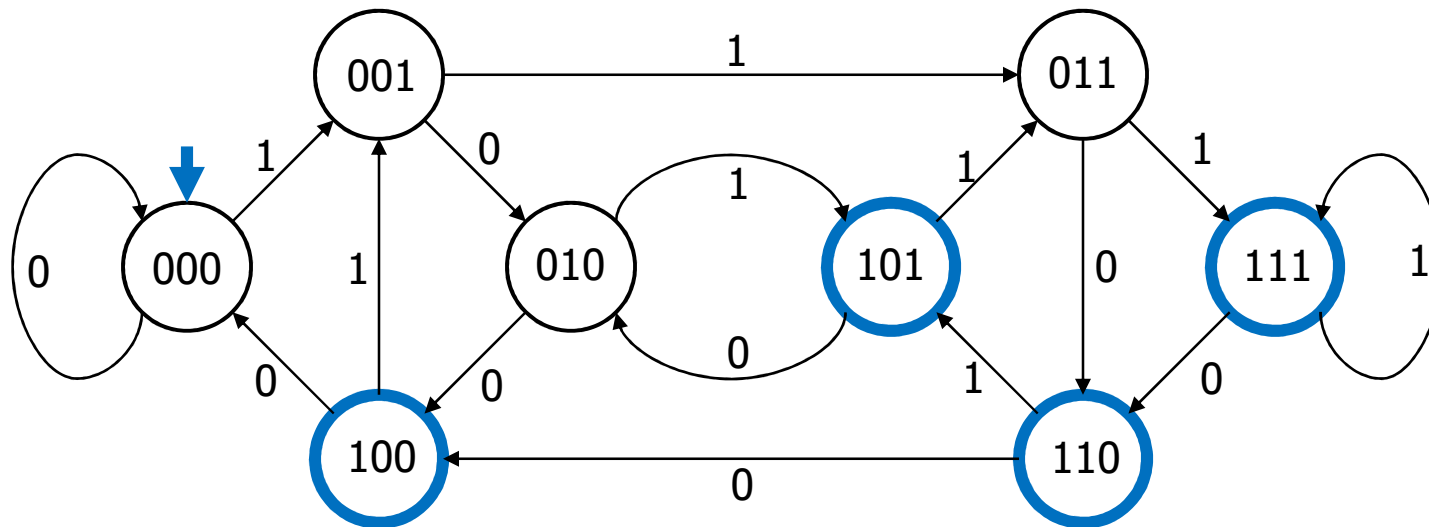
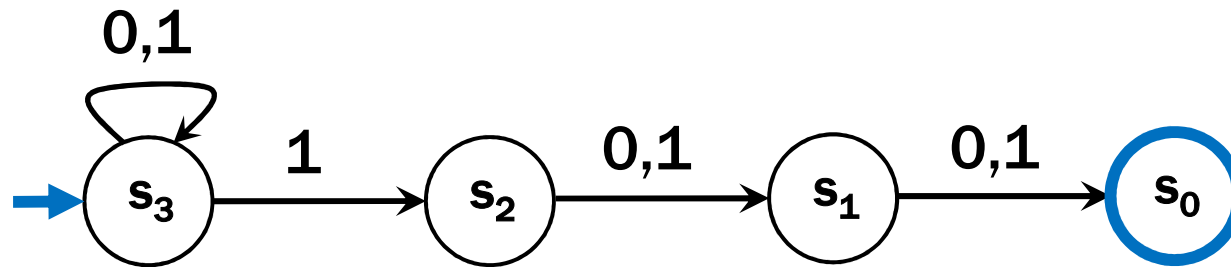
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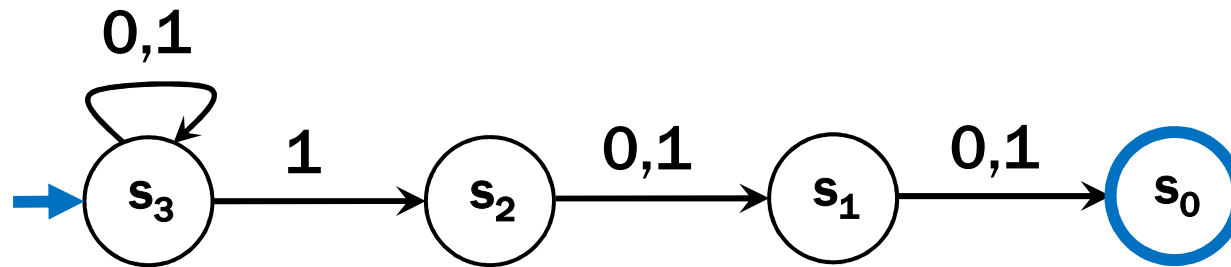
# Compare with the smallest DFA

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# Parallel Exploration view of an NFA

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Input string 0101100

