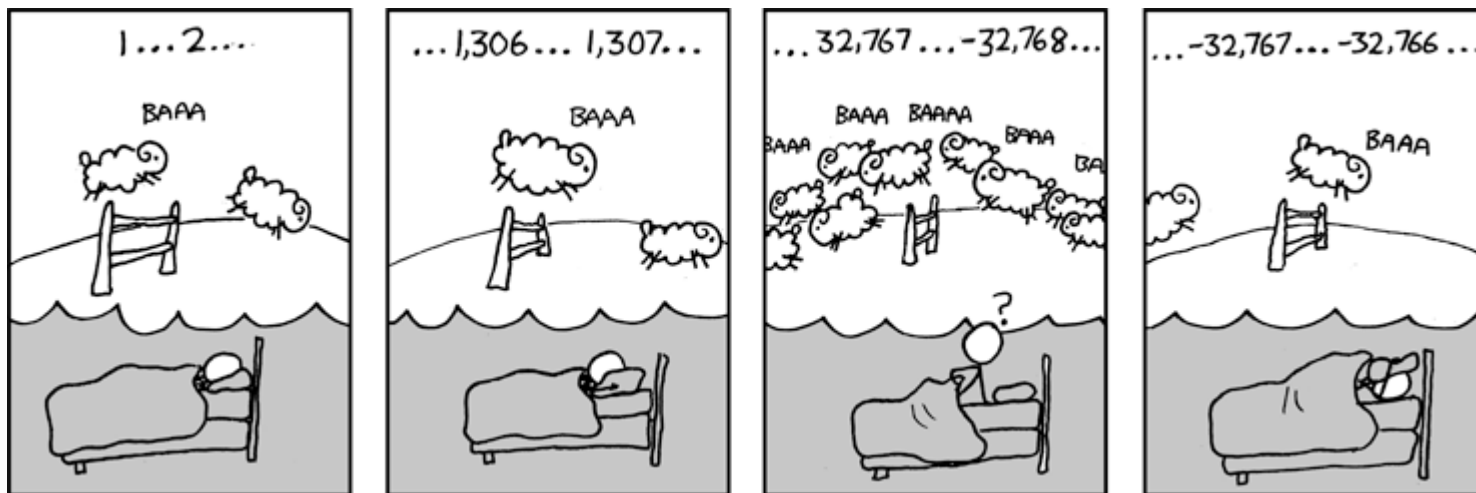


# CSE 311: Foundations of Computing

## Lecture 11: Modular Arithmetic and Applications

2.6(d) Sol<sup>n</sup> (arithmetic) see <sup>black</sup> board



Backward  
reasoning

~~Prof:  $2 = 1$   
 $\therefore 0.2 = 0.1$   
 $0 = 0$~~

# Last Class: Divisibility

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## Definition: “a divides b”

For  $a \in \mathbb{Z}, b \in \mathbb{Z}$  with  $a \neq 0$ :

$$a \mid b \leftrightarrow \exists k \in \mathbb{Z} (b = ka)$$

Check Your Understanding. Which of the following are true?

$$5 \mid 1$$

$$5 \mid 1 \text{ iff } 1 = 5k$$

$$25 \mid 5$$

$$25 \mid 5 \text{ iff } 5 = 25k$$

$$5 \mid 0$$

$$5 \mid 0 \text{ iff } 0 = 5k$$

$$3 \mid 2$$

$$3 \mid 2 \text{ iff } 2 = 3k$$

$$1 \mid 5$$

$$1 \mid 5 \text{ iff } 5 = 1k$$

$$5 \mid 25$$

$$5 \mid 25 \text{ iff } 25 = 5k$$

$$0 \mid 5$$

$$0 \mid 5 \text{ iff } 5 = 0k$$

$$2 \mid 3$$

$$2 \mid 3 \text{ iff } 3 = 2k$$

# Last Class: Division Theorem

## Division Theorem

For  $a \in \mathbb{Z}, d \in \mathbb{Z}$  with  $d > 0$   
there exist *unique* integers  $q, r$  with  $0 \leq r < d$   
such that  $a = dq + r$ .

To put it another way, if we divide  $d$  into  $a$ , we get a  
unique quotient  $q = a \text{ div } d$   
and non-negative remainder  $r = a \text{ mod } d$

$$a = (a \text{ div } d) \cdot d + (a \text{ mod } d)$$

```
public class Test2 {  
    public static void main(String args[]) {  
        int a = -5;  
        int d = 2;  
        System.out.println(a % d);  
    }  
}
```

```
-----jGRASP exec: java Test2  
-1  
-----jGRASP: operation complete.
```

Note:  $r \geq 0$  even if  $a < 0$ .  
Not quite the same as  $a \% d$ .

# Last Class: Arithmetic, mod 7

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$$a +_7 b = (a + b) \bmod 7$$

$$a \times_7 b = (a \times b) \bmod 7$$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

# Modular Arithmetic

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**Definition: “a is congruent to b modulo m”**

For  $a, b, m \in \mathbb{Z}$  with  $m > 0$

$$a \equiv b \pmod{m} \leftrightarrow \underline{m \mid (a - b)}$$

**Check Your Understanding. What do each of these mean?  
When are they true?**

$$x \equiv 0 \pmod{2}$$

This statement is the same as saying “x is even”; so, any x that is even (including negative even numbers) will work.

$$-1 \equiv 19 \pmod{5}$$

This statement is true.  $19 - (-1) = 20$  which is divisible by 5

$$y \equiv 2 \pmod{7}$$

This statement is true for y in  $\{ \dots, -12, -5, 2, 9, 16, \dots \}$ . In other words, all y of the form  $2+7k$  for k an integer.

# Modular Arithmetic: A Property

$$a \text{ div } m = \lfloor \frac{a}{m} \rfloor$$

Let  $a, b, m$  be integers with  $m > 0$ .

Then,  $a \equiv b \pmod{m}$  if and only if  $a \text{ mod } m = b \text{ mod } m$ .

Suppose that  $a \equiv b \pmod{m}$ .

$$\begin{aligned} \therefore m \mid (a-b) & \quad \text{by defn of } \equiv \pmod{m} \\ \therefore a-b = km & \quad \text{for some integer } k \\ \therefore a = b + km & \quad \text{By defn } b = qm + (b \text{ mod } m) \\ & \quad \text{(for some integer } k) \\ \therefore a = qm + (b \text{ mod } m) + km & \\ & = (k+q)m + (b \text{ mod } m) \\ & \quad \text{By defn } 0 \leq (b \text{ mod } m) < m \\ \therefore a \text{ mod } m & = b \text{ mod } m \end{aligned}$$

Suppose that  $a \text{ mod } m = b \text{ mod } m$ .

(call this  $r$ )

$$\begin{aligned} \therefore a &= qm + r \quad 0 \leq r < m \\ b &= q'm + r \quad \text{for integers } q, q' \\ \therefore a-b &= qm - q'm = (q-q')m \\ \therefore m \mid (a-b) & \\ \therefore a &\equiv b \pmod{m} \end{aligned}$$

# Modular Arithmetic: A Property

---

Let  $a, b, m$  be integers with  $m > 0$ .

Then,  $a \equiv b \pmod{m}$  if and only if  $a \bmod m = b \bmod m$ .

Suppose that  $a \equiv b \pmod{m}$ .

Then,  $m \mid (a - b)$  by definition of congruence.

So,  $a - b = km$  for some integer  $k$  by definition of divides.

Therefore,  $a = b + km$ .

Taking both sides modulo  $m$  we get:

$$a \bmod m = (b + km) \bmod m = b \bmod m.$$

Suppose that  $a \bmod m = b \bmod m$ .

By the division theorem,  $a = mq + (a \bmod m)$  and

$$b = ms + (b \bmod m) \text{ for some integers } q, s.$$

$$\begin{aligned} \text{Then, } a - b &= (mq + (a \bmod m)) - (ms + (b \bmod m)) \\ &= m(q - s) + (a \bmod m - b \bmod m) \\ &= m(q - s) \text{ since } a \bmod m = b \bmod m \end{aligned}$$

Therefore,  $m \mid (a - b)$  and so  $a \equiv b \pmod{m}$ .

# The mod $m$ function vs the $\equiv \pmod{m}$ predicate

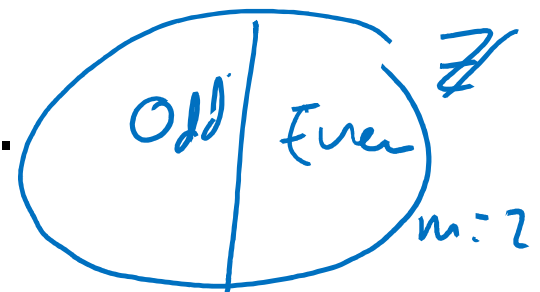
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- **What we have just shown**

- The mod  $m$  function takes any  $a \in \mathbb{Z}$  and maps it to a remainder  $a \bmod m \in \{0, 1, \dots, m - 1\}$ .
- Imagine grouping together all integers that have the same value of the mod  $m$  function  
That is, the same remainder in  $\{0, 1, \dots, m - 1\}$ .
- The  $\equiv \pmod{m}$  predicate compares  $a, b \in \mathbb{Z}$ . It is true if and only if the mod  $m$  function has the same value on  $a$  and on  $b$ .



That is,  $a$  and  $b$  are in the same group.





# Modular Arithmetic: Addition Property

$$a \equiv b \pmod{m}$$

$$a \equiv b$$

$$a + c \equiv b + d$$

Intuition

Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $a + c \equiv b + d \pmod{m}$

Proof

Since  $a \equiv b \pmod{m}$  we get  
 $m \mid (a - b)$

$\therefore a - b = km$  for some integer  $k$

Since  $c \equiv d \pmod{m}$  we get  
 $m \mid (c - d)$

$\therefore c - d = lm$  for some integer  $l$

$\therefore a = b + km$  and  $c = d + lm$

$\therefore a + c = b + d + km + lm = b + d + (k + l)m$

$$(a + c) - (b + d) = (k + l)m$$

$$m \mid (a + c) - (b + d)$$

$$\therefore a + c \equiv b + d \pmod{m} \quad \square$$

# Modular Arithmetic: Addition Property

---

Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $a + c \equiv b + d \pmod{m}$

Suppose that  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ . Unrolling definitions gives us some  $k$  such that  $a - b = km$ , and some  $j$  such that  $c - d = jm$ .

Adding the equations together gives us

$(a + c) - (b + d) = m(k + j)$ . Now, re-applying the definition of congruence gives us  $a + c \equiv b + d \pmod{m}$ .

# Modular Arithmetic: Multiplication Property

---

Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $ac \equiv bd \pmod{m}$

Proof Since  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$   
we have  $m \mid (a-b)$  and  $m \mid (c-d)$   
 $\therefore a-b = km$  and  $c-d = lm$  for some  
integers  $k, l$ .  
 $\therefore a = b + km$  and  $c = d + lm$   
 $a \cdot c = (b + km)(d + lm) = b \cdot d + b \cdot lm + km \cdot d + klm^2$   
 $= b \cdot d + m(bl + kd + klm)$   
 $\therefore a \cdot c - b \cdot d = m(bl + kd + klm)$   
 $\therefore m \mid (a \cdot c - b \cdot d)$   
 $\therefore ac \equiv bd \pmod{m} \quad \square$

# Modular Arithmetic: Multiplication Property

---

Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $ac \equiv bd \pmod{m}$

Suppose that  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ . Unrolling definitions gives us some  $k$  such that  $a - b = km$ , and some  $j$  such that  $c - d = jm$ .

Then,  $a = km + b$  and  $c = jm + d$ . Multiplying both together gives us  $ac = (km + b)(jm + d) = kjm^2 + kmd + bjm + bd$ .

Re-arranging gives us  $ac - bd = m(kjm + kd + bj)$ .

Using the definition of congruence gives us  $ac \equiv bd \pmod{m}$ .

# Example

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Let  $n$  be an integer.

Prove that  $n^2 \equiv \mathbf{0 \pmod{4}}$  or  $n^2 \equiv \mathbf{1 \pmod{4}}$

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

# Example

---

Let  $n$  be an integer.

Prove that  $n^2 \equiv 0 \pmod{4}$  or  $n^2 \equiv 1 \pmod{4}$

Case 1 ( $n$  is even):

$$\begin{aligned} \therefore n &= 2k \text{ for some integer } k \\ \therefore n^2 &= 4k^2 \\ \therefore n^2 \pmod{4} &= 0 \\ \therefore n^2 &\equiv 0 \pmod{4} \end{aligned}$$

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

It looks like

Case 2 ( $n$  is odd):

$$\begin{aligned} \therefore n &= 2k+1 \text{ for some integer } k \\ \therefore n^2 &= (2k+1)^2 = 4k^2 + 4k + 1 \\ &= 4(k^2 + k) + 1 \\ \therefore n^2 \pmod{4} &= 1 \\ \therefore n^2 &\equiv 1 \pmod{4} \end{aligned}$$

$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$ , and

$n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$ .

# Example

---

Let  $n$  be an integer.

Prove that  $n^2 \equiv 0 \pmod{4}$  or  $n^2 \equiv 1 \pmod{4}$

Case 1 ( $n$  is even):

Suppose  $n \equiv 0 \pmod{2}$ .

Then,  $n = 2k$  for some integer  $k$ .

So,  $n^2 = (2k)^2 = 4k^2$ . So, by

definition of congruence,

$n^2 \equiv 0 \pmod{4}$ .

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

It looks like

Case 2 ( $n$  is odd):

Suppose  $n \equiv 1 \pmod{2}$ .

Then,  $n = 2k + 1$  for some integer  $k$ .

So,  $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$ .

So, by definition of congruence,  $n^2 \equiv 1 \pmod{4}$ .

$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$ , and

$n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$ .

# n-bit Unsigned Integer Representation

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- Represent integer  $x$  as sum of powers of 2:

If  $\sum_{i=0}^{n-1} b_i 2^i$  where each  $b_i \in \{0,1\}$

then representation is  $b_{n-1} \dots b_2 b_1 b_0$

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

- For  $n = 8$ :  
99: 0110 0011  
18: 0001 0010

---

$$117: 0111 0101$$



# Sign-Magnitude Integer Representation

---

## *n*-bit signed integers

Suppose that  $-2^{n-1} < x < 2^{n-1}$

First bit as the sign,  $n - 1$  bits for the value

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

For  $n = 8$ :

99: 0110 0011

-18: 1001 0010

1111 0101

Any problems with this representation?

# Two's Complement Representation

---

$n$  bit signed integers, first bit will still be the sign bit

Suppose that  $0 \leq x < 2^{n-1}$ ,  
 $x$  is represented by the binary representation of  $x$

Suppose that  $0 \leq x \leq 2^{n-1}$ ,  
 $-x$  is represented by the binary representation of  $2^n - x$

**Key property:** Two's complement representation of any number  $y$  is equivalent to  $y \bmod 2^n$  so arithmetic works **mod**  $2^n$

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

For  $n = 8$ :

$$99: 0110\ 0011$$

$$-18: 1110\ 1110$$

$$2^n - x$$

$$128 + 64 + 32 - 16 + 8 + 4 + 2 + 1 = 255 \\ = 256 - 1 \\ 2^8 - 1$$

# Sign-Magnitude vs. Two's Complement

---

-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
1111	1110	1101	1100	1011	1010	1001	0000	0001	0010	0011	0100	0101	0110	0111

Sign-bit 1000

<u>-8</u>	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
1000	1001	1010	1011	1100	1101	1110	1111	0000	0001	0010	0011	0100	0101	0110	0111

Two's complement



# Two's Complement Representation

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- For  $0 < x \leq 2^{n-1}$ ,  $-x$  is represented by the binary representation of  $2^n - x$ 
  - That is, the two's complement representation of any number  $y$  has the same value as  $y$  modulo  $2^n$ .
- To compute this: Flip the bits of  $x$  then add 1:
  - All 1's string is  $2^n - 1$ , so

Flip the bits of  $x \equiv$  replace  $x$  by  $2^n - 1 - x$

Then add 1 to get  $2^n - x$

$18 = 16 + 2$

$$\begin{array}{r}
 \begin{array}{cccccccc}
 & \overset{1}{\cdot} & \overset{1}{\cdot} & \overset{1}{\cdot} & \overset{1}{\cdot} & \overset{1}{\cdot} & \overset{1}{\cdot} & \overset{1}{\cdot} \\
 y & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\
 \hline
 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\
 & & & & & & & +1 \\
 \hline
 & & & & & & & 1 \\
 & & & & & & & 1 \\
 & & & & & & & 1 \\
 & & & & & & & 1 \\
 & & & & & & & 0 \\
 & & & & & & & 1 \\
 & & & & & & & 1 \\
 & & & & & & & 1 \\
 & & & & & & & 0
 \end{array}
 & \leftarrow 2^n - 1 & & 11101110
 \end{array}$$

# Basic Applications of mod

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- Hashing
- Pseudo random number generation
- Simple cipher

# Hashing

---

## Scenario:

Map a small number of data values from a large domain  $\{0, 1, \dots, M - 1\}$  ...

...into a small set of locations  $\{0, 1, \dots, n - 1\}$  so one can quickly check if some value is present

- $\text{hash}(x) = x \bmod p$  for  $p$  a prime close to  $n$ 
  - or  $\text{hash}(x) = (ax + b) \bmod p$
- Depends on all of the bits of the data
  - helps avoid collisions due to similar values
  - need to manage them if they occur

# Pseudo-Random Number Generation

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## Linear Congruential method

$$x_{n+1} = (a x_n + c) \bmod m$$

**Choose random  $x_0, a, c, m$  and produce a long sequence of  $x_n$ 's**

# Simple Ciphers

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- **Caesar cipher**,  $A = 1$ ,  $B = 2, \dots$ 
  - HELLO WORLD
- **Shift cipher**
  - $f(p) = (p + k) \bmod 26$
  - $f^{-1}(p) = (p - k) \bmod 26$
- **More general**
  - $f(p) = (ap + b) \bmod 26$