Adam Blank Spring 2017

CSE 311

Foundations of Computing I

Some Reminders/Hints for HW3

Style matters in proofs!

Do NOT manipulate large statements by equivalences! This is horrible style and will lose points if there is a significantly cleaner proof.

Our inference rules can only prove things true. Do not "prove the negation false"! It doesn't make sense.

After you're done writing your proof, you should proof-read (heh) it.

Set Operations

$A \cup B = \{ x : (x \in A) \vee (x \in B) \}$ **Union**

$A \cap B = \{ x : (x \in A) \wedge (x \in B) \}$ **Intersection**

$A \setminus B = \{ x : (x \in A) \wedge (x \notin B) \}$ **Set Difference**

A = {1, 2, 3}
B = {4, 5, 6}
C = {3, 4}

QUESTIONS
Using A, B, C and set operations, make...
{6} = {1,2,3,4,5,6} = A∪B = A∪B∪C
{3} = A∩C
{1,2} =

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B = {4, 5, 6}
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QUESTIONS
Using A, B, C and set operations, make...
[6] = A ∪ B = A ∪ B ∪ C
{3} = C ∖ B = A ∖ B = A ∩ C
{1,2} = A ∖ C = (A ∪ B) ∖ C

More Set Operations

$A \oplus B = \{ x : (x \in A) \oplus (x \in B) \}$ **Symmetric Difference**

$\overline{A} = \{ x : x \notin A \}$ (with respect to universe U) **Complement**

A = {1, 2, 3}
B = {1, 4, 2, 6}
C = {1, 2, 3, 4}

QUESTIONS
Let S = {1, 2}.
If the universe is A, then $\overline{S}$ is... {3}
If the universe is B, then $\overline{S}$ is... {4,6}
If the universe is C, then $\overline{S}$ is...

More Set Operations

$A \oplus B = \{ x : (x \in A) \oplus (x \in B) \}$ **Symmetric Difference**

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B = {1, 4, 2, 6}
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QUESTIONS
Let S = {1, 2}.
If the universe is A, then $\overline{S}$ is... **A ∖ S = {3}**
If the universe is B, then $\overline{S}$ is... **B ∖ S = {4, 6}**
If the universe is C, then $\overline{S}$ is... **C ∖ S = {3, 4}**

Power Set

- Power Set of a set A = set of all subsets of A

$$\mathcal{P}(A) = \{ B : B \subseteq A \}$$

- Let Days = {M, W, F}. Suppose we wanted to know the possible ways that we could allocate class days to be cancelled. Let's call this set $\mathcal{P}(\text{Days})$.

e.g. $\mathcal{P}(\text{Days}) = \{$

$\emptyset$, $\{M, W, F\}$

$\{M\}, \{W\}, \{F\}, \{M, M\}, M\}$

$\}$

e.g. $\mathcal{P}(\emptyset) = ?$

Power Set

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$\emptyset,$
$\{M\}, \{W\}, \{F\},$
$\{M, W\}, \{W, F\}, \{M, F\},$
$\{M, W, F\}$
$\}$

e.g. $\mathcal{P}(\emptyset) = \{\emptyset\}$

Cartesian Product

$$A \times B = \{ (a, b) : a \in A, b \in B \}$$

$\mathbb{R} \times \mathbb{R}$ is the real plane. You've seen ordered pairs before.

These are just for arbitrary sets.

$\mathbb{Z} \times \mathbb{Z}$ is "the set of all pairs of integers"

If A = {1, 2}, B = {a, b, c}, then A × B = {(1,a), (1,b), (1,c), (2,a), (2,b), (2,c)}.

$A \times \emptyset = \{(a, b) : a \in A \wedge b \in \emptyset\} = \{(a, b) : a \in A \wedge F\} = \emptyset$

Russell's Paradox

$$S = \{ x : x \notin x \}$$

Suppose for contradiction that S ∈ S...

S contains itself.

$S \notin S$

$S \notin S \rightarrow S \in S$

Russell's Paradox

$$S = \{ x : x \notin x \}$$

Suppose for contradiction that S ∈ S. Then, by definition of S, S ∉ S, but that's a contradiction.

Suppose for contradiction that S ∉ S. Then, by definition of the set comprehension, S ∈ S, but that's a contradiction.

This is reminiscent of the truth value of the statement "This statement is false."

It's Boolean algebra again

- Definition for ∪ based on ∨

$$A \cup B = \{ x : (x \in A) \vee (x \in B) \}$$

- Definition for ∩ based on ∧

$$A \cap B = \{ x : (x \in A) \wedge (x \in B) \}$$

- Complement works like ¬

$$\overline{A} = \{ x : x \notin A \}$$

(with respect to universe U)

De Morgan's Laws

Prove $\overline{A \cup B} = \overline{A} \cap \overline{B}$

Let U be the universe.

$$\begin{aligned}\overline{A \cup B} &= \{x : x \notin A \cup B\} \\ &= \{x : \neg(x \in A \cup B)\} \\ &= \{x : \neg((x \in A) \vee (x \in B))\} \\ &= \{x : (x \notin A) \wedge (x \notin B)\} \\ &= \{x : (x \in \overline{A}) \wedge (x \in \overline{B})\} \\ &= \{x : (x \in \overline{A})\} \cap \{x : (x \in \overline{B})\} \\ &= \overline{A} \cap \overline{B}\end{aligned}$$

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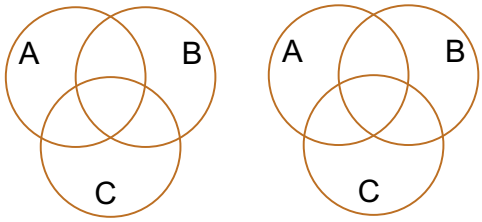
Distributive Laws

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$
$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$



One More

Prove $(A \cap B) \subseteq A$

Remember the definition of subset?

$X \subseteq Y \leftrightarrow \forall x\,(x \in X \rightarrow x \in Y)$

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Remember the definition of subset?

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Let x be an arbitrary element of $A \cap B$. Then, by definition of $A \cap B$, $x \in A$ and $x \in B$. It follows that $x \in A$, as required.

Representing Sets Using Bits

- Suppose universe U is $\{1,2,\ldots,n\}$
- Can represent set $B \subseteq U$ as a vector of bits:

  $b_1 b_2 \ldots b_n$ where $b_i = 1$ when $i \in B$

  $b_i = 0$ when $i \notin B$

  – Called the *characteristic vector* of set B

- Given characteristic vectors for A and B
  – What is characteristic vector for $A \cup B$? $A \cap B$?

UNIX/Linux File Permissions

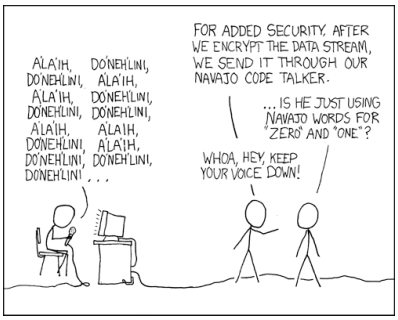
- `ls -l`

```
drwxr-xr-x ... Documents/
-rw-r--r-- ... file1
```

- Permissions maintained as bit vectors
  – Letter means bit is 1
  – "–" means bit is 0.

CSE 311: Foundations of Computing

Lecture 10: Modular Arithmetic



Number Theory (and applications to computing)

- Branch of Mathematics with direct relevance to computing
- Many significant applications
  – Cryptography
  – Hashing
  – Security
- Important tool set

Modular Arithmetic

- Arithmetic over a finite domain
- In computing, almost all computations are over a finite domain

I'm ALIVE!

```
public class Test {
   final static int SEC_IN_YEAR = 364*24*60*60*100;
   public static void main(String args[]) {
      System.out.println(
         "I will be alive for at least " +
         SEC_IN_YEAR * 101 + " seconds."
      );
   }
}
```

I'm ALIVE!

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public class Test {
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         "I will be alive for at least " +
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      );
   }
}
```

```
----jGRASP exec: java Test
I will be alive for at least -186619904 seconds.

----jGRASP: operation complete.
```

Divisibility

Definition: "a divides b"

For $a \in \mathbb{Z}, b \in \mathbb{Z}$ with $a \neq 0$:

$$a \mid b \leftrightarrow \exists(k \in \mathbb{Z})\ b = ka$$

Check Your Understanding. Which of the following are true?

5 | 1 **25 | 5** **5 | 5** **3 | 2**

1 | 5 **5 | 25** **0 | 1** **2 | 3**

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$$a \mid b \leftrightarrow \exists(k \in \mathbb{Z})\ b = ka$$

Check Your Understanding. Which of the following are true?

5 | 1	**25 | 5**	**5 | 5**	**3 | 2**
5 | 1 iff 1 = 5k	25 | 1 iff 1 = 25k	5 | 5 iff 5 = 5k	3 | 2 iff 2 = 3k
1 | 5	**5 | 25**	**0 | 1**	**2 | 3**
1 | 5 iff 5 = 1k	1 | 25 iff 25 = 1k	0 | 1 iff 1 = 0k	2 | 3 iff 3 = 2k

Division Theorem

Division Theorem

For $a \in \mathbb{Z}, d \in \mathbb{Z}^+$:
Then, there exists *unique* integers q, r with $0 \leq r < d$ such that $a = dq + r$.

To put it another way, if we take a/d, we get a dividend and a remainder: $q = a$ **div** d $\quad r = a$ **mod** d

Note: r ≥ 0 even if a < 0.
Not quite the same as `a % d`.

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```
public class Test2 {
    public static void main(String args[]) {
        int a = -5;
        int d = 2;
        System.out.println(a % d);
    }
}
```

```
----jGRASP exec: java Test2
-1

----jGRASP: operation complete.
```

Note: r ≥ 0 even if a < 0.
Not quite the same as `a % d`.

Arithmetic, mod 7

$a +_7 b = (a + b) \bmod 7$

$a \times_7 b = (a \times b) \bmod 7$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Modular Arithmetic

Definition: "a is congruent to b modulo m"

For $a \in \mathbb{Z}, b \in \mathbb{Z}, \mathrm{m} \in \mathbb{Z}$:

$$a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$$

Check Your Understanding. What do each of these mean? When are they true?

A ≡ 0 (mod 2)

1 ≡ 0 (mod 4)

A ≡ -1 (mod 17)

Modular Arithmetic

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Check Your Understanding. What do each of these mean? When are they true?

A ≡ 0 (mod 2)

This statement is the same as saying "A is even"; so, any A that is even (including negative even numbers) will work.

1 ≡ 0 (mod 4)

This statement is false. If we take it mod 1 instead, then the statement is true.

A ≡ -1 (mod 17)

If A = 17x – 1 = 17x + 16, then it works.
Note that (m – 1) mod m = ((m mod m) + (-1 mod m)) mod m
= (0 + -1) mod m = -1 mod m

Modular Arithmetic: A Property

Let a and b be integers, and let m be a positive integer. Then, a ≡ b (mod m) if and only if a mod m = b mod m.

Suppose that a ≡ b (mod m).

Suppose that a mod m = b mod m.

Modular Arithmetic: A Property

Let a and b be integers, and let m be a positive integer. Then, a ≡ b (mod m) if and only if a mod m = b mod m.

Suppose that a ≡ b (mod m).
Then, m | (a – b) by definition of congruence.
So, a – b = km for some integer k by definition of divides.
Therefore, a = b+km.
Taking both sides modulo m we get:
a mod m=(b+km) mod m = b mod m.

Suppose that a mod m = b mod m.
By the division theorem, a = mq + (a mod m) and
b = ms + (b mod m) for some integers q,s.
Then, a – b = (mq + (a mod m)) – (mr + (b mod m))
= m(q – r) + (a mod m – b mod m)
= m(q – r) since a mod m = b mod m
Therefore, m | (a-b) and so $a \equiv b \pmod{m}$.

Modular Arithmetic: Another Property

Let m be a positive integer. If a ≡ b (mod m) and c ≡ d (mod m), then **a + c ≡ b + d (mod m)**

Suppose a ≡ b (mod m) and c ≡ d (mod m). Unrolling definitions gives us some k such that
a – b = km, and some j such that c – d = jm.

Adding the equations together gives us
(a + c) – (b + d) = m(k + j). Now, re-applying the definition of mod gives us a + c ≡ b + d (mod m).

Modular Arithmetic: Another-nother Property

Let m be a positive integer. If a ≡ b (mod m) and c ≡ d (mod m), then **ac ≡ bd (mod m)**

Suppose a ≡ b (mod m) and c ≡ d (mod m). Unrolling definitions gives us some k such that
a – b = km, and some j such that c – d = jm.

Then, a = km + b and c = jm + d. Multiplying both together gives us ac = (km + b)(jm + d) = kjm^2 + kmd + jmb + bd.

Re-arranging gives us ac – bd = m(kjm + kd + jb). Using the definition of mod gives us ac ≡ bd (mod m).

Example

Let n be an integer.
Prove that $n^2 \equiv 0 \pmod 4$ or $n^2 \equiv 1 \pmod 4$

Case 1 (n is even):
Suppose $n \equiv 0 \pmod 2$.
Then, $n = 2k$ for some k.
So, $n^2 = (2k)^2 = 4k^2$. So, by definition of congruence, $n^2 \equiv 0 \pmod 4$.

Case 2 (n is odd):
Suppose $n \equiv 1 \pmod 2$.
Then, $n = 2k + 1$ for some k.
So, $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$. So, by definition of congruence, $n^2 \equiv 1 \pmod 4$.

Let's start by looking a a small example:

$0^2 = 0 \equiv 0 \pmod 4$
$1^2 = 1 \equiv 1 \pmod 4$
$2^2 = 4 \equiv 0 \pmod 4$
$3^2 = 9 \equiv 1 \pmod 4$
$4^2 = 16 \equiv 0 \pmod 4$

It looks like
$n \equiv 0 \pmod 2 \rightarrow n^2 \equiv 0 \pmod 4$, and
$n \equiv 1 \pmod 2 \rightarrow n^2 \equiv 1 \pmod 4$.

Bitwise Operations

```
  01101101        Java:   z=x|y
v 00110111
  01111111

  00101010        Java:   z=x&y
^ 00001111
  00001010

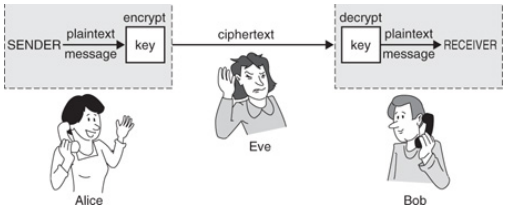
  01101101        Java:   z=x^y
⊕ 00110111
  01011010
```

A Useful Identity

- If x and y are bits: $(x \oplus y) \oplus y = ?$
  - $(x \oplus y) \oplus y = x \oplus (y \oplus y) = x \oplus 0 = x$

Private Key Cryptography

- **Alice** wants to communicate message secretly to **Bob** so that eavesdropper **Eve** who hears their conversation cannot tell what **Alice**'s message is.
- **Alice** and **Bob** can get together and privately share a secret key K ahead of time.



One-Time Pad

- Alice and Bob privately share random n-bit vector K
  - Eve does not know K
- Later, Alice has n-bit message m to send to Bob
  - Alice computes $C = m \oplus K$
  - Alice sends C to Bob
  - Bob computes $m = C \oplus K$ which is $(m \oplus K) \oplus K$
- Eve cannot figure out m from C unless she can guess K