

**CSE  
31F**

# Foundations of Computing I

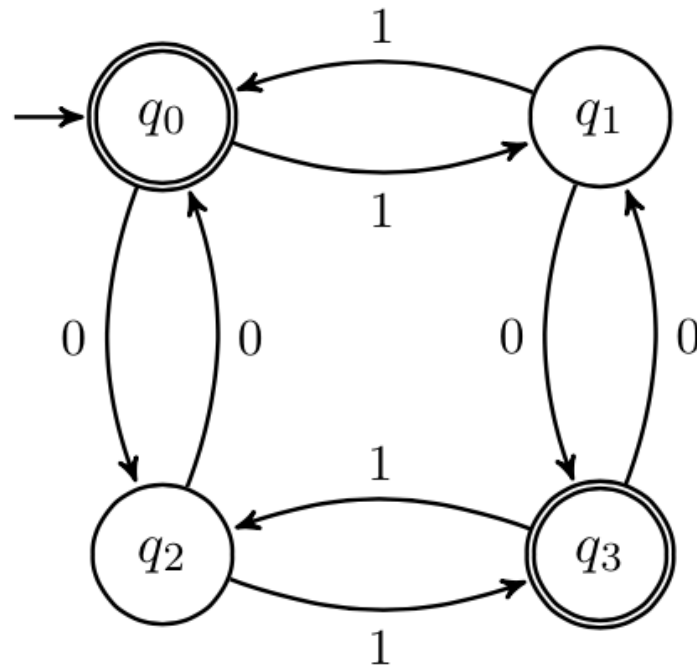
\* All slides are a combined effort between  
previous instructors of the course

# CSE 311: Foundations of Computing

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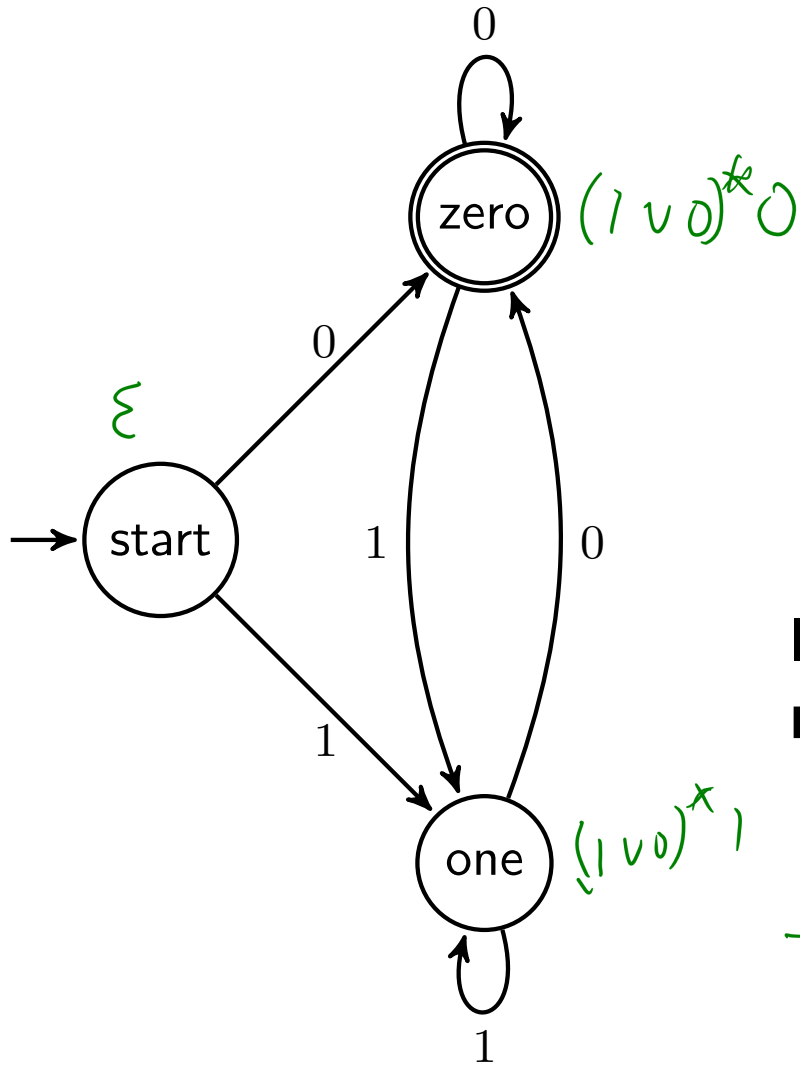
## Lecture 20: Finite State Machines (DFAs)

Hi!!  
😊



# A Weird Sort of Programming!

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0101  $\rightarrow$  one

101  $\rightarrow$  one

10

0

What does this “thing” do?

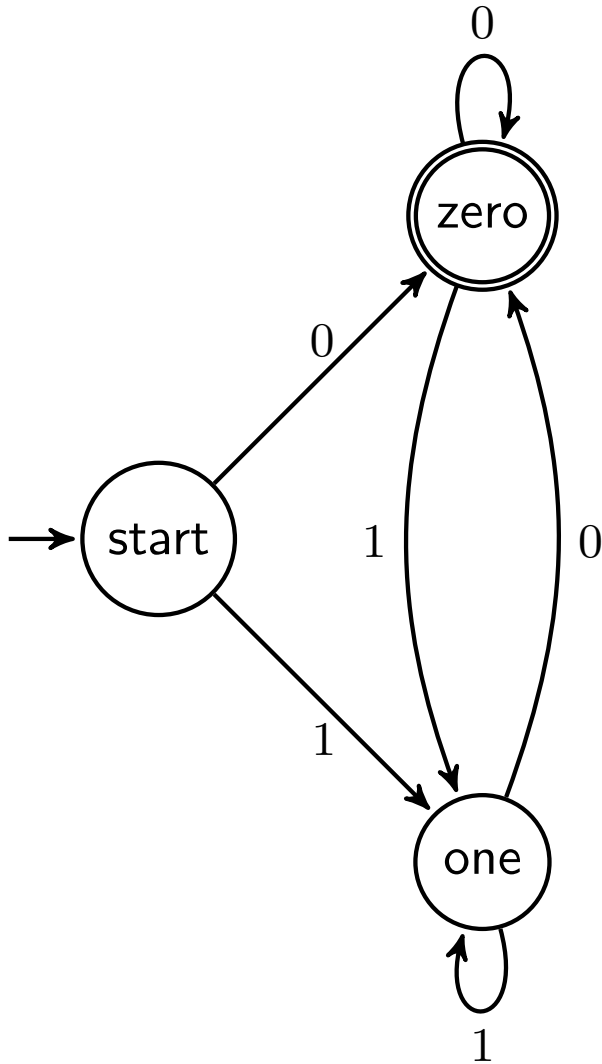
Take a guess!

If you had to give this “method” a name, what would it be?

boolean isEven (String s)

# A Weird Sort of Programming!

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What does this “thing” do?

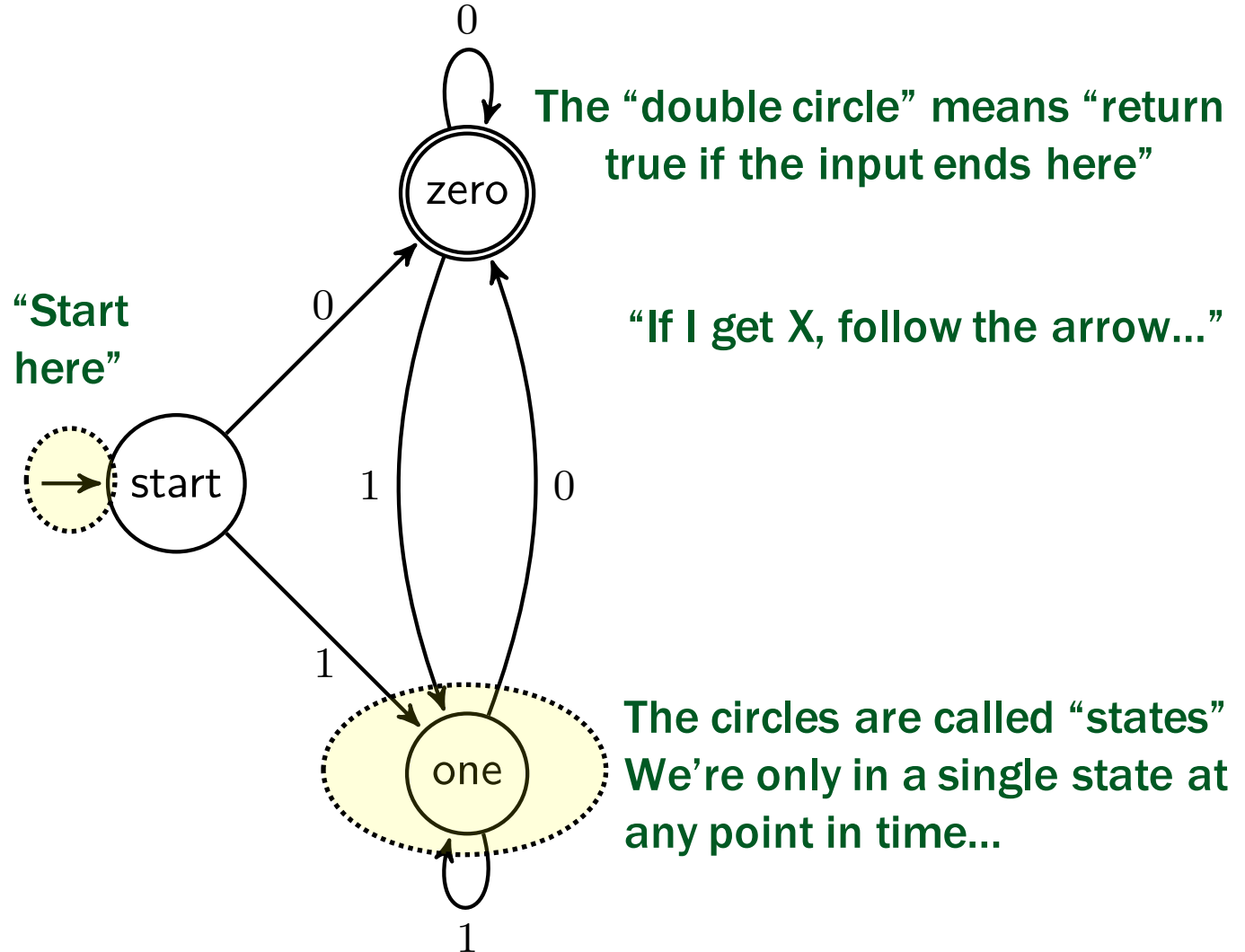
Take a guess!

If you had to give this “method” a name, what would it be?

`boolean isEven(binary s)`

# Finite State Machines (“DFAs”)

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# Applications of FSMs (a.k.a. finite automata)

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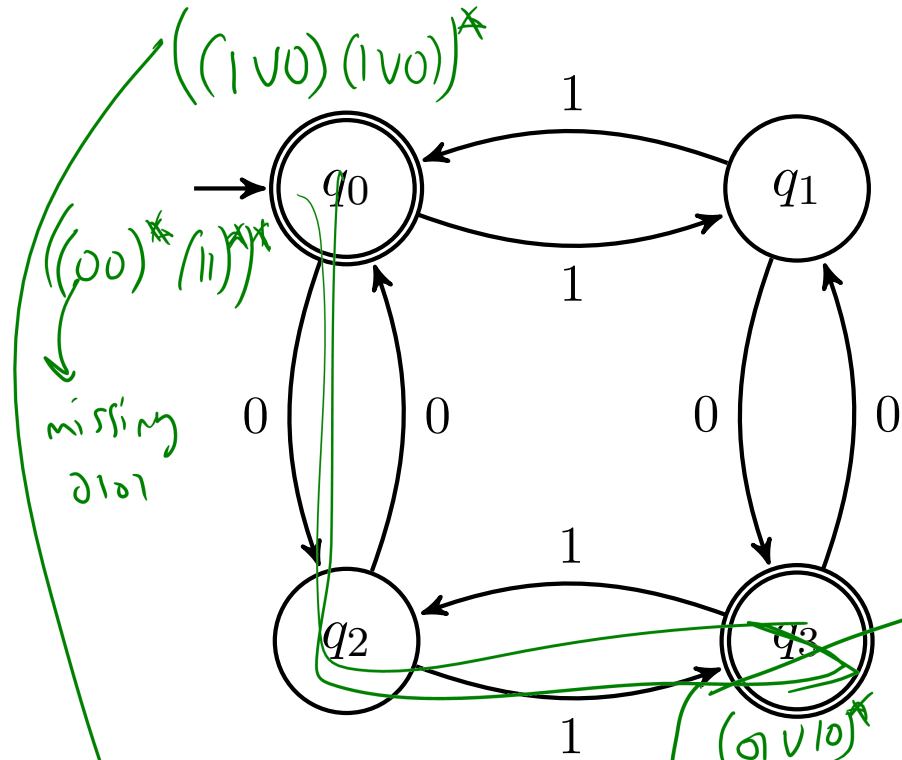
- Implementation of regular expression matching in programs like `grep`
- Control structures for sequential logic in digital circuits
- Algorithms for communication and cache-coherence protocols
  - Each agent runs its own FSM
- Design specifications for reactive systems
  - Components are communicating FSMs

# Applications of FSMs (a.k.a. finite automata)

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- **Formal verification of systems**
  - Is an unsafe state reachable?
- **Computer games**
  - FSMs provide worlds to explore
- **Minimization algorithms for FSMs can be extended to more general models used in**
  - Text prediction
  - Speech recognition

# What language does this machine recognize?



$((1100)(1100))^*$

$((00)^*(11))^*$

missing 0101

$q_0:$

$((00)^*(11)^*(1010)^*(0101)^*)^*$

$0110)^*$

$(1001)^*$

$(0110)^*$

$((00)^*(11)^*(0110))^*$

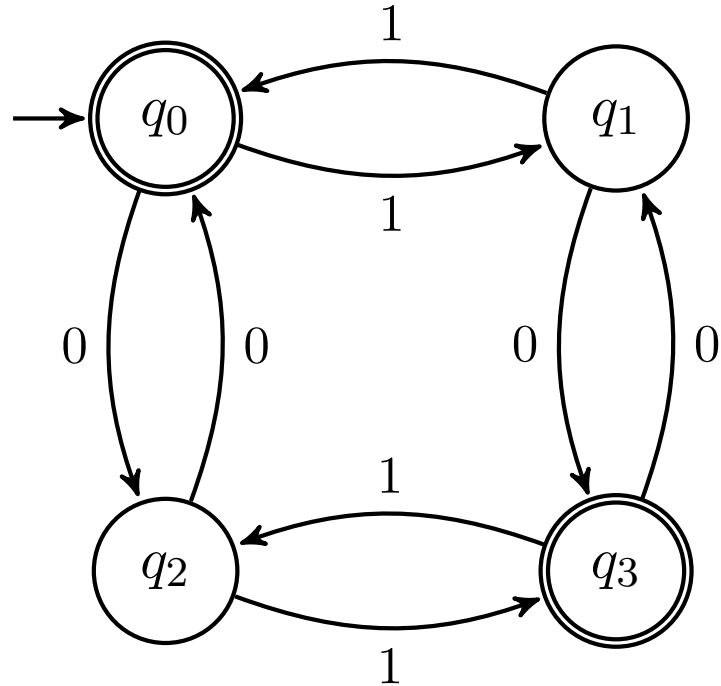
binary strings w/ even length

$(00)^* \cup (0101)^*$   
 $\cup (0011)^*$   
 $\cup (11)^*$   
 $\cup (01)^*$



# What language does this machine recognize?

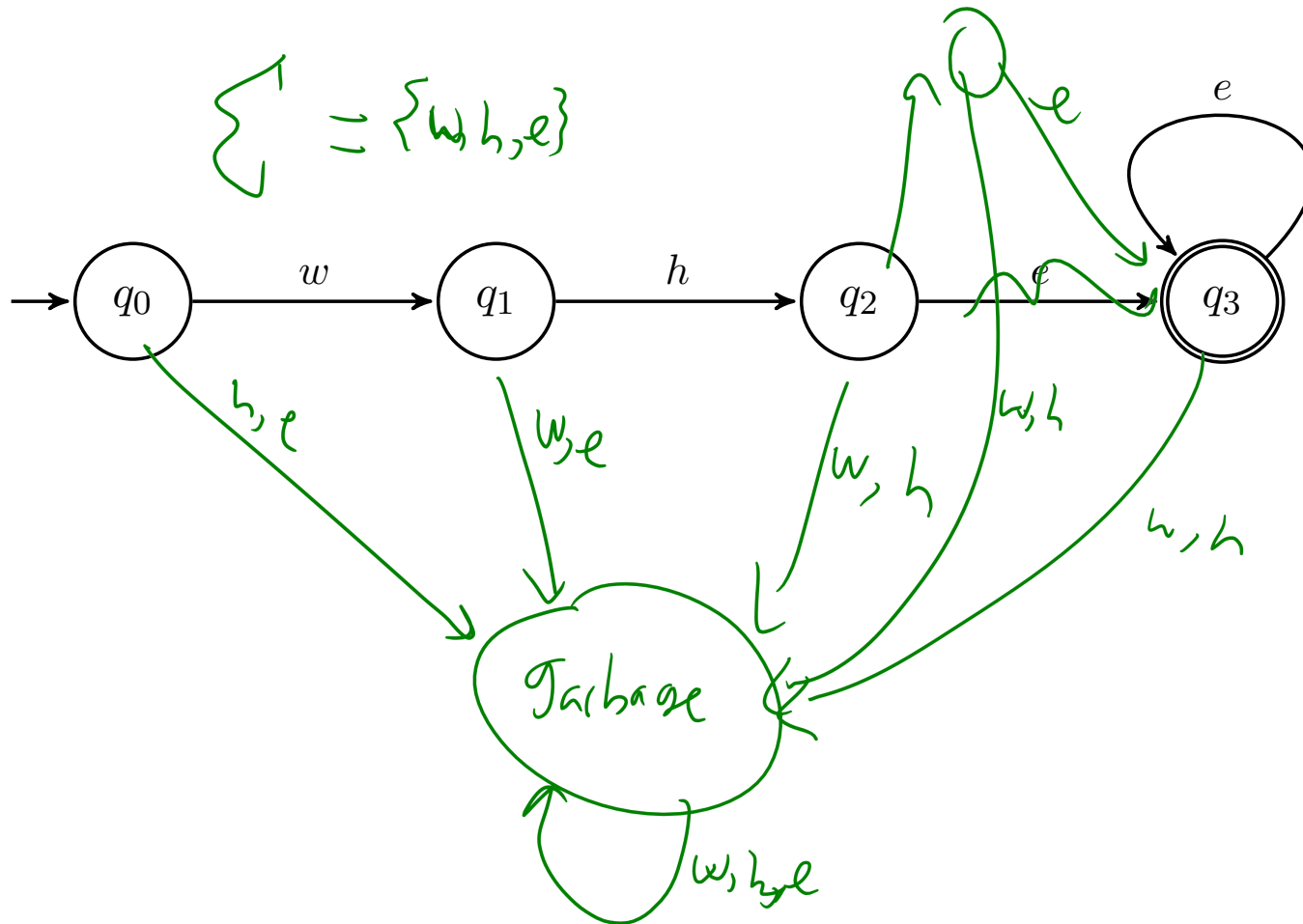
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All binary strings with even length

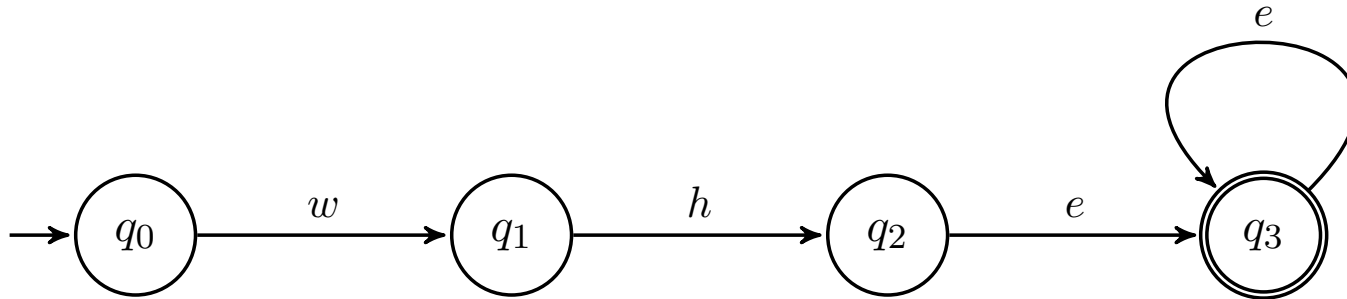
# Why is this not a DFA? Fix it!

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# Why is this not a DFA? Fix it!

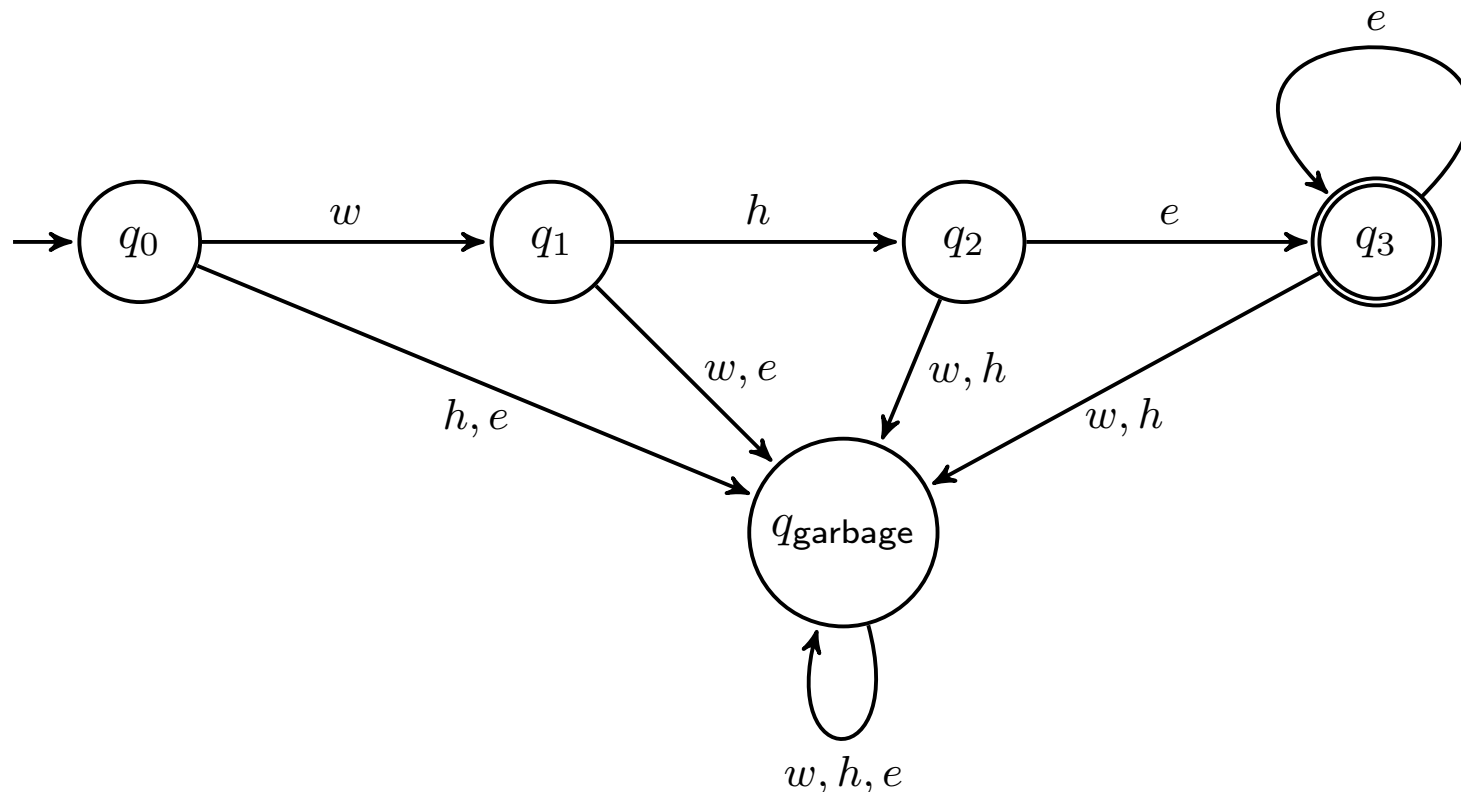
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**DFAs must have a transition for every character at every state!**

# Why is this not a DFA? Fix it!

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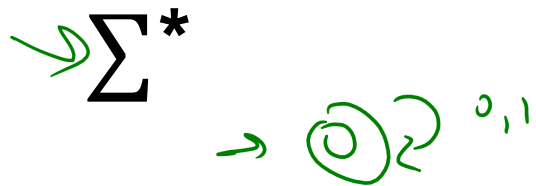
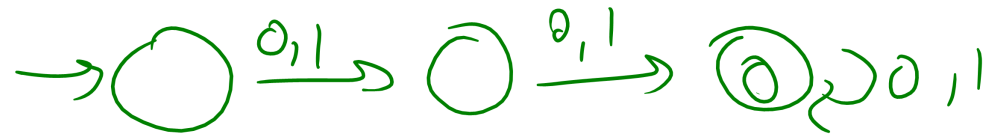
“Garbage states” are a useful concept. Whenever we KNOW we can’t accept the string, just send it to a state that will always go back to itself. This is the way of saying “return false” in DFA-land.

For each of the following languages, create a DFA

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$\emptyset$  

$\{x \in \{0,1\}^* : \text{len}(x) > 1\}$

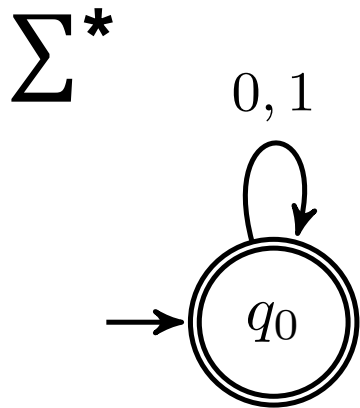
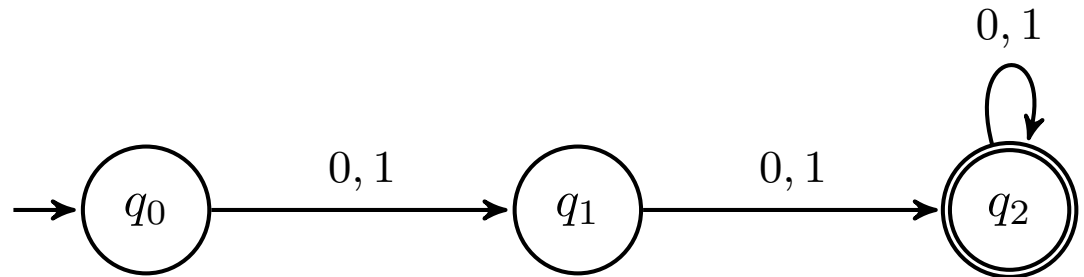
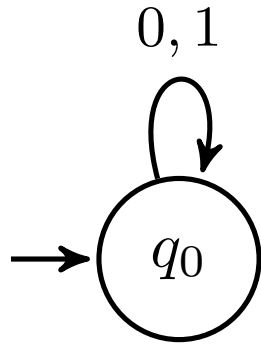


$\varepsilon \neq \emptyset$   
 $\{\varepsilon\}$

# For each of the following languages, create a DFA

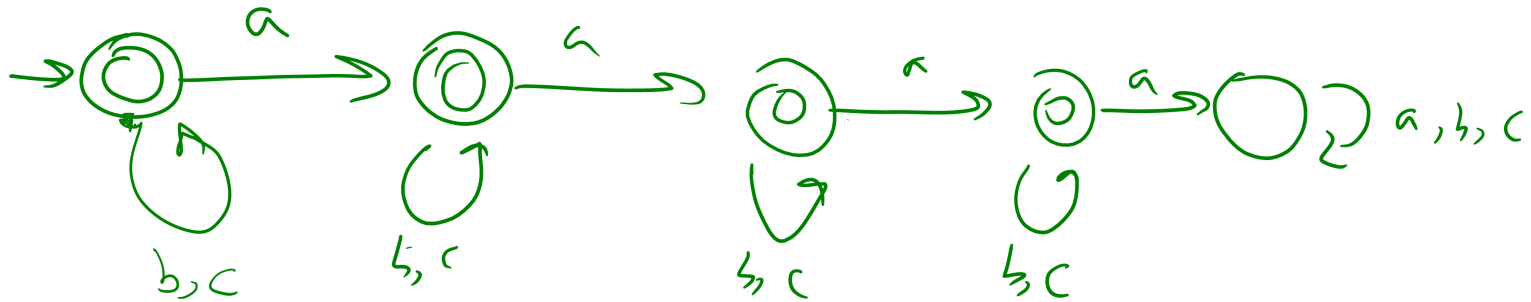
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$\emptyset$   $\{x \in \{0,1\}^* : \text{len}(x) > 1\}$



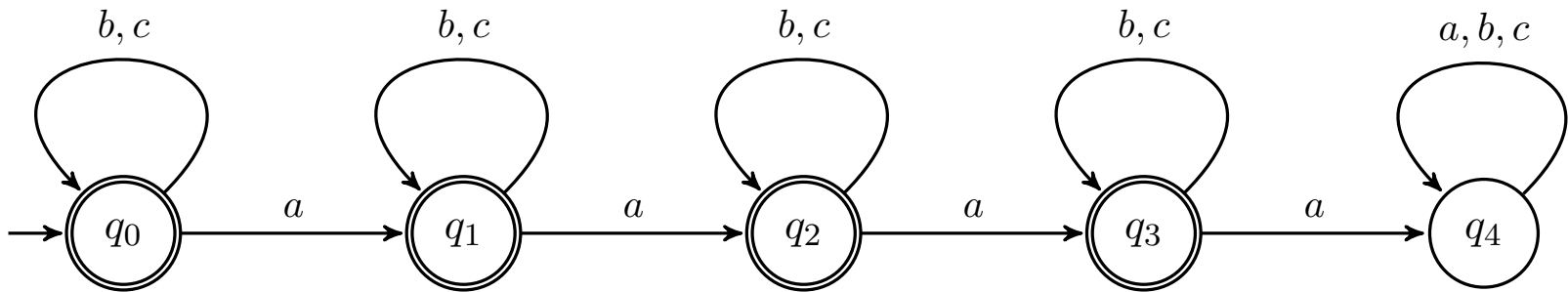
# FSM that accepts strings of a's, b's, c's with no more than 3 a's

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# FSM that accepts strings of a's, b's, c's with no more than 3 a's

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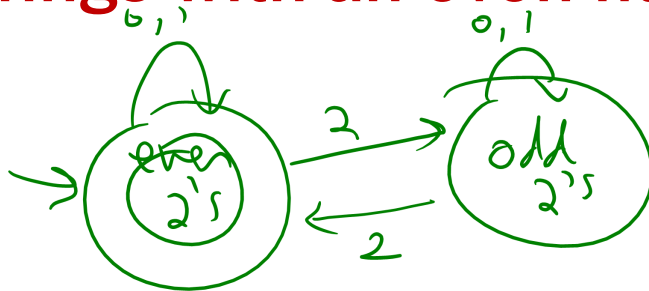




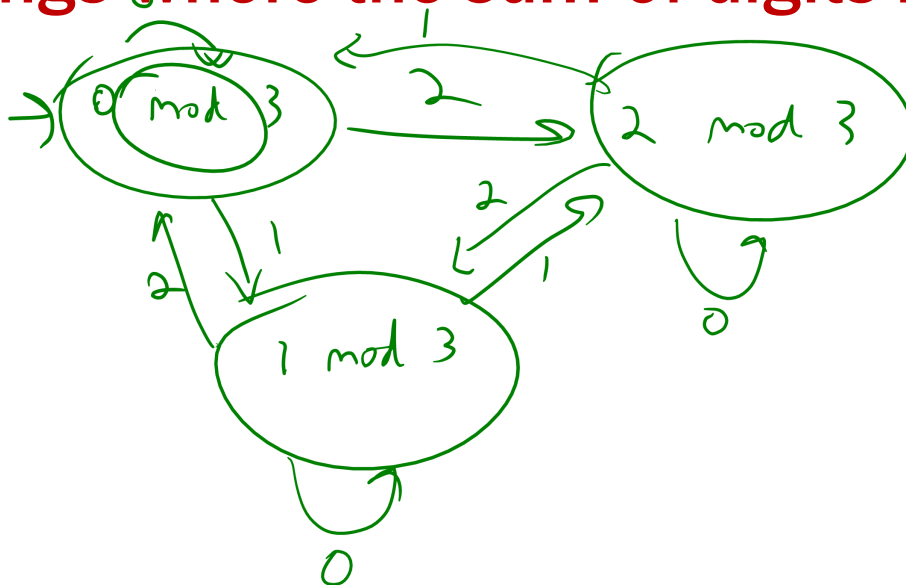
# Strings over $\{0, 1, 2\}^*$

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**$M_1$ : Strings with an even number of 2's**



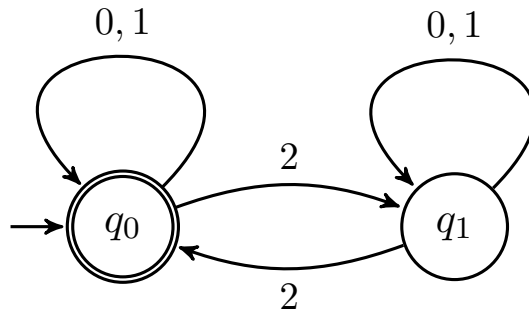
**$M_2$ : Strings where the sum of digits mod 3 is 0**



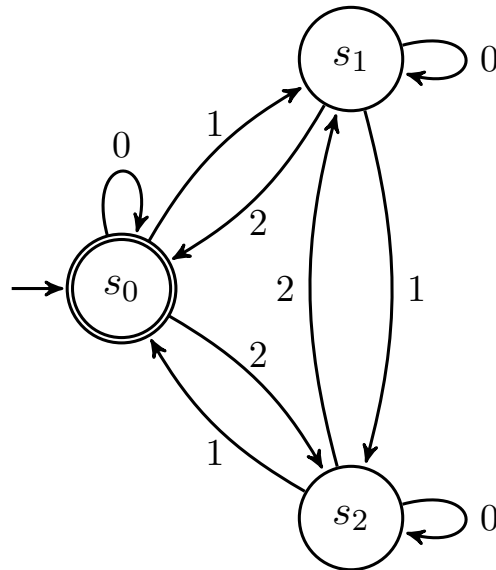
# Strings over $\{0, 1, 2\}^*$

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**$M_1$ : Strings with an even number of 2's**

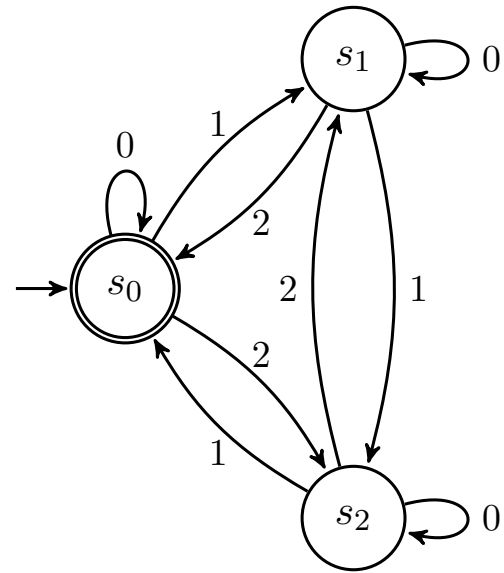
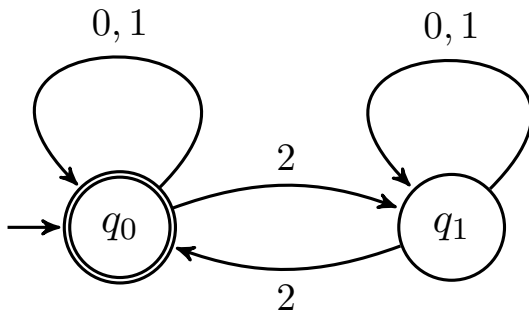


**$M_2$ : Strings where the sum of digits mod 3 is 0**



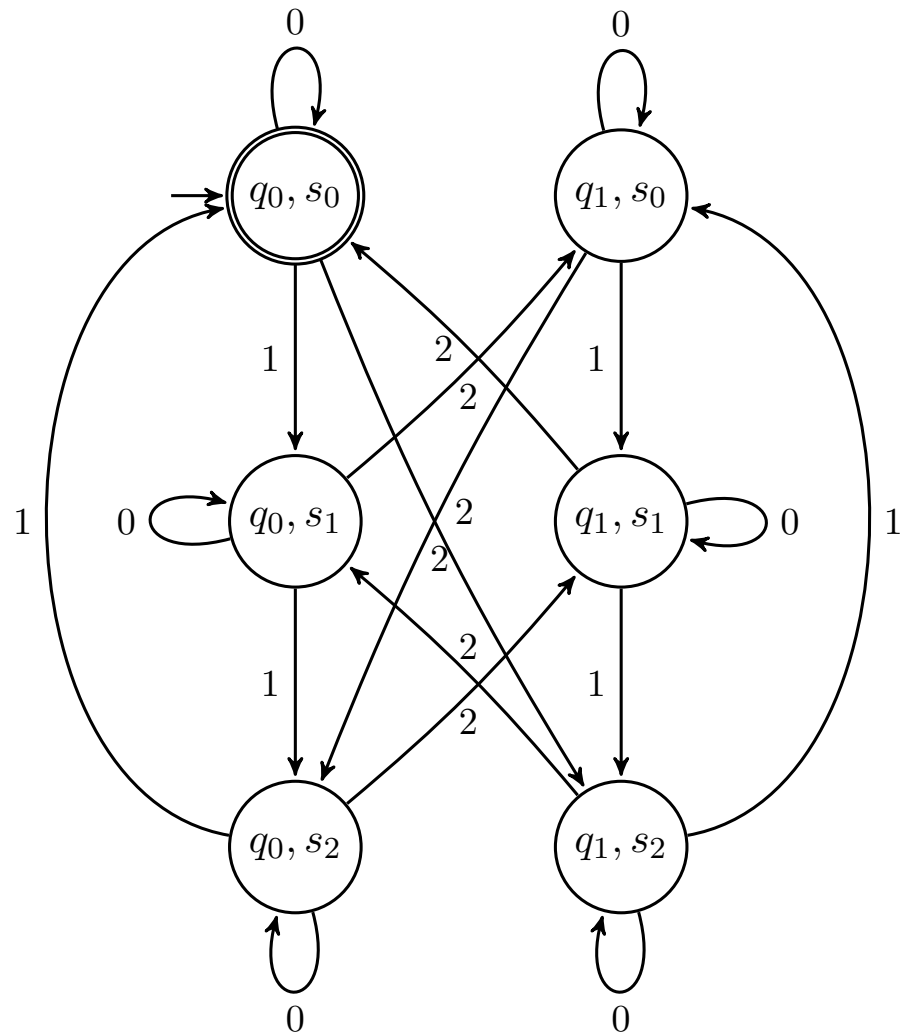
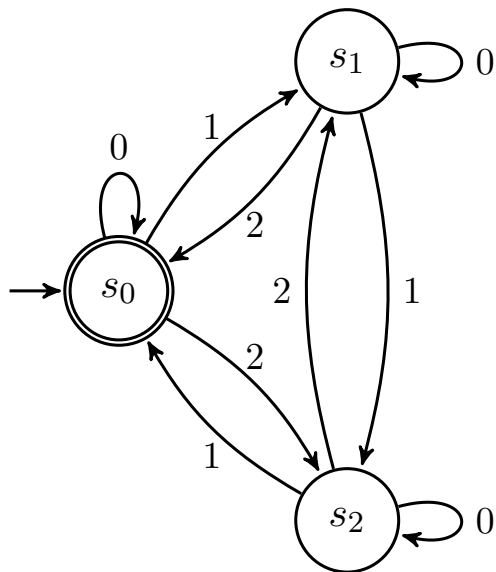
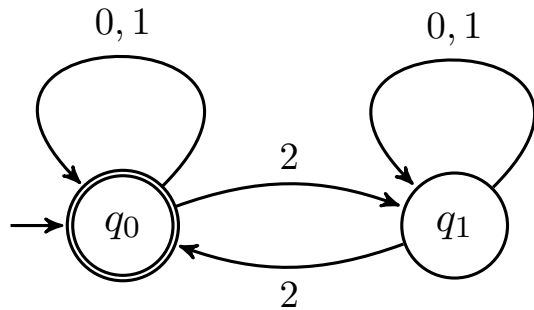
# Strings with an even number of 2's AND a mod 3 sum of 0

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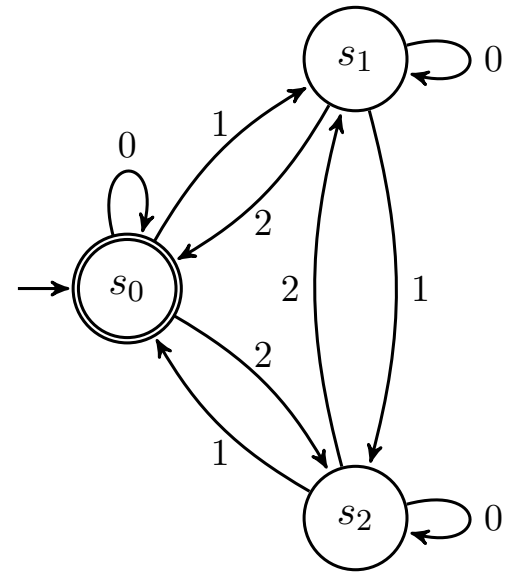
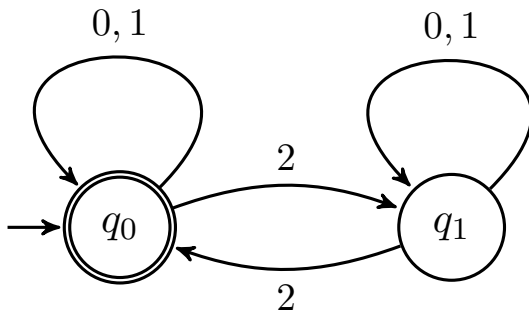
# Strings with an even number of 2's AND a mod 3 sum of 0

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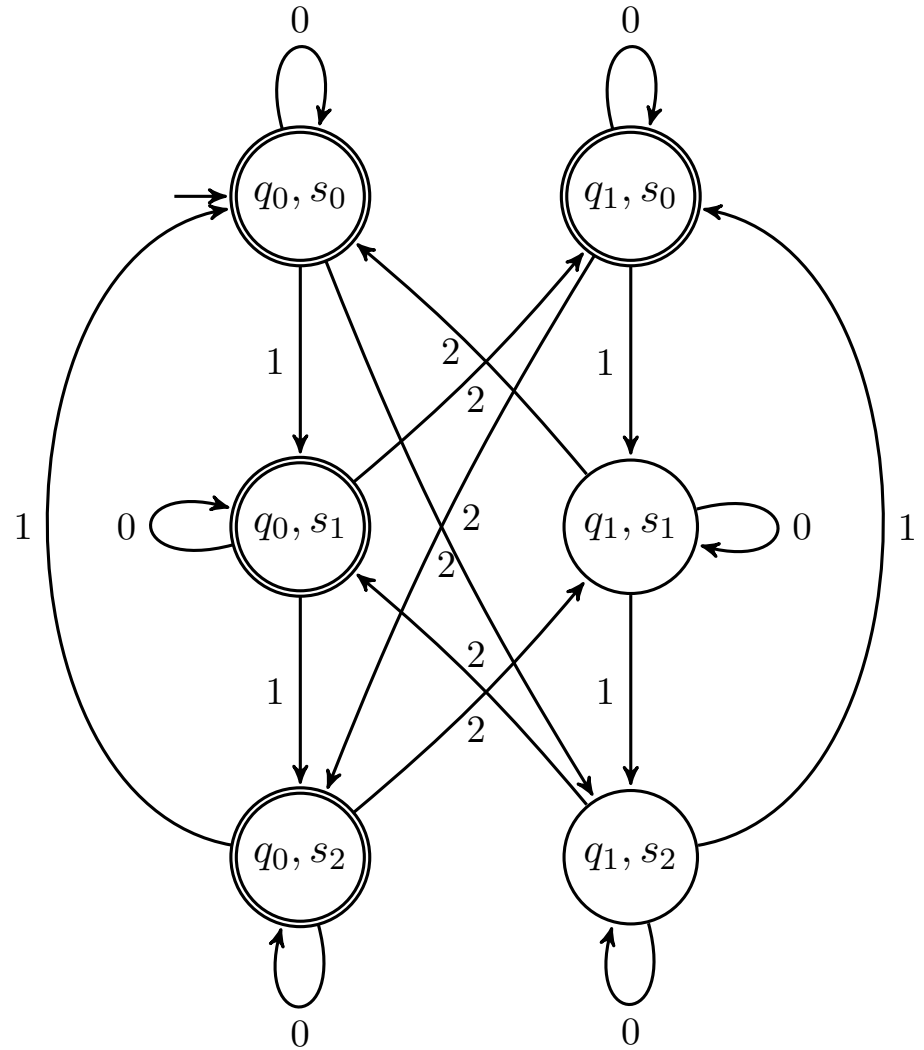
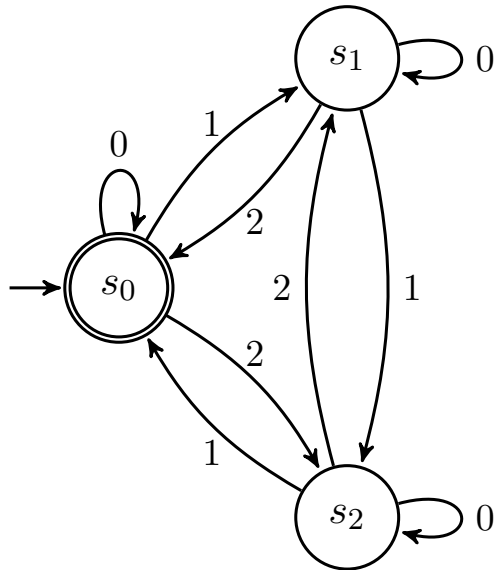
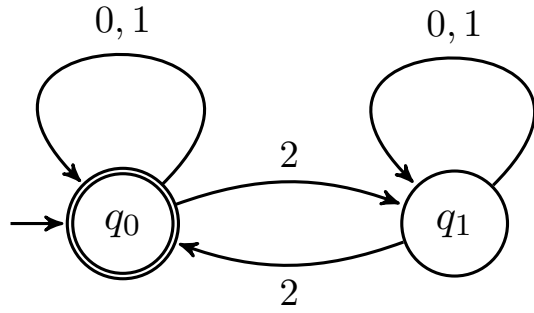
# Strings with an even number of 2's OR a mod 3 sum of 0

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# Strings with an even number of 2's OR a mod 3 sum of 0

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**FSM that accepts binary strings with a 1 three positions from the start**

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# FSM that accepts binary strings with a 1 three positions from the start

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