

# CSE 311: Foundations of Computing I

## Section 5: Midterm Practice Solutions

### 0. Propositional Logic

(a) Is the following expression a contingency, contradiction, or tautology?

$$(p \rightarrow q) \wedge (q \rightarrow r) \rightarrow (p \rightarrow r)$$

**Solution:**

Tautology.

(b) Show that  $\neg p \rightarrow (q \rightarrow r)$  and  $q \rightarrow (p \vee r)$  are logically equivalent.

**Solution:**

**Solution 1: Truth Table**

p	q	r	$\neg p$	$q \rightarrow r$	$\neg p \rightarrow (p \rightarrow r)$
T	T	T	F	T	T
T	T	F	F	F	T
T	F	T	F	T	T
T	F	F	F	T	T
F	T	T	T	T	T
F	T	F	T	F	F
F	F	T	T	T	T
F	F	F	T	T	T

p	q	r	$p \vee r$	$q \rightarrow (p \vee r)$
T	T	T	T	T
T	T	F	T	T
T	F	T	T	T
T	F	F	T	T
F	T	T	T	T
F	T	F	F	F
F	F	T	F	T
F	F	F	F	T

Since the last column is identical on both tables the two statements must be equivalent.

**Solution 2: Formal Proof**

$$\begin{aligned} \neg p \rightarrow (q \rightarrow r) &= \neg \neg p \vee (q \rightarrow r) && \text{By Law of Implication} \\ &= p \vee (q \rightarrow r) && \text{By Double Negation} \\ &= p \vee (\neg q \vee r) && \text{By Law of Implication} \\ &= (\neg q \vee r) \vee p && \text{By Commutativity} \\ &= \neg q \vee (r \vee p) && \text{By Associativity} \\ &= \neg q \vee (p \vee r) && \text{By Commutativity} \\ &= q \rightarrow (p \vee r) && \text{By Law of Implication} \end{aligned}$$

## 1. Predicate Logic

Let the domain of discourse be all plants and leaves. You may use the predicates  $\text{HasLeaf}(x, y) ::= "x \text{ has } y \text{ as a leaf}"$ ,  $\text{Equals}(x, y) ::= "x \text{ is the same object as } y"$ ,  $\text{Leaf}(x) ::= "x \text{ is a leaf}"$  and  $\text{Plant}(x) ::= "x \text{ is a plant}"$ ,  $\text{IsPurple}(x)$  be " $x$  is purple",  $\text{IsGolden}(x)$  be " $x$  is golden", and let the constant  $\text{LuckyLeaf}$  be the Lucky Leaf.

Translate the following sentences to predicate logic using quantifiers.

- (a) Every plant has at least 2 leaves.

**Solution:**

$$\forall x (\text{Plant}(x) \rightarrow \exists y \exists z (\neg \text{Equals}(y, z) \wedge \text{HasLeaf}(x, y) \wedge \text{HasLeaf}(x, z)))$$

- (b) Every plant has at most 2 leaves.

**Solution:**

$$\forall x (\text{Plant}(x) \rightarrow \forall a \forall b \forall c ((\text{HasLeaf}(x, a) \wedge \text{HasLeaf}(x, b) \wedge \text{HasLeaf}(x, c)) \rightarrow (\text{Equals}(a, b) \vee \text{Equals}(a, c) \vee \text{Equals}(b, c))))$$

- (c) There is exactly one plant that has no leaves.

**Solution:**

$$\exists x (\text{Plant}(x) \wedge \neg(\exists y(\text{HasLeaf}(x, y))) \wedge \forall z ((\text{Plant}(z) \wedge \neg(\exists k (\text{HasLeaf}(z, k)))) \rightarrow \text{Equals}(x, z)))$$

- (d) If a plant has the Lucky Leaf, all other leaves on that plant are golden, but the Lucky Leaf is purple, and then no other plants have golden or purple leaves.

**Solution:**

$$\begin{aligned} &\text{IsPurple}(\text{LuckyLeaf}) \wedge \forall x (\text{Plant}(x) \wedge \text{HasLeaf}(x, \text{LuckyLeaf}) \rightarrow \\ &\quad \forall y ((\text{HasLeaf}(x, y) \wedge \neg \text{Equal}(y, \text{LuckyLeaf})) \rightarrow \\ &\quad \text{IsGolden}(y)) \wedge \forall z ((\text{Plant}(z) \wedge \neg \text{Equal}(x, z)) \rightarrow \\ &\quad \forall k (\text{HasLeaf}(z, k) \rightarrow (\neg \text{IsPurple}(k) \wedge \neg \text{IsGolden}(k)))))) \end{aligned}$$

## 2. Proofs with Number Theory

- (a) Prove that if  $n^2 + 1$  is a perfect square, where  $n$  is an integer, then  $n$  is even.

### Solution:

Let  $n$  be an arbitrary integer, and suppose  $n^2 + 1$  is a perfect square. Then, by definition of perfect square,  $n^2 + 1 = k^2$  for some  $k \in \mathbb{N}$ . Since  $n$  and  $k$  are integers, we can define some integer  $z$  such that  $k = n + z$ . Now, substituting, we get:

$$\begin{aligned}n^2 + 1 &= (n + z)^2 \\n^2 + 1 &= n^2 + 2nz + z^2 \\1 &= 2nz + z^2 \\1 &= z(2n + z) \\\frac{1}{z} &= (2n + z)\end{aligned}$$

Since  $n$  and  $z$  are integers,  $2n + z$  is an integer, which means  $\frac{1}{z}$  is an integer. The only integers which satisfy this constraint are  $z = \pm 1$ , and in both these cases  $z = \frac{1}{z}$ , so we can subtract  $z$  from both sides to find  $n = 0$  as the only solution. Since  $n = 0$ , and 0 is even,  $n$  is even.

- (b) Prove that if  $n$  is a positive integer such that the sum of the divisors of  $n$  is  $n + 1$ , then  $n$  is prime.

### Solution:

Let  $n$  be an arbitrary positive integer, and suppose the sum of the divisors of  $n$  is  $n + 1$ . Note that  $n \mid n$ . Since the sum of divisors of  $n$  is  $n + 1$ ,  $n + 1 - n = 1$  must be the only other divisor. It follows, by definition of prime, that  $n$  is prime. Since we supposed  $n$  was a positive integer whose divisors sum to  $n + 1$ , we see that for any arbitrary positive integer  $n$ , if  $n$ 's divisors sum to  $n + 1$ ,  $n$  must be prime.

## 3. Induction

- (a) Prove for all  $n \in \mathbb{N}$  that if you have two groups of numbers,  $a_1, \dots, a_n$  and  $b_1, \dots, b_n$ , such that  $\forall(i \in [n]). a_i \leq b_i$ , then it must be that:

$$\sum_{i=1}^n a_i \leq \sum_{i=1}^n b_i$$

### Solution:

Let  $P(n)$  be the statement: "For any two groups of numbers,  $a_1, \dots, a_n$  and  $b_1, \dots, b_n$ , such that  $\forall(i \in [n]). a_i \leq b_i$ , it is true that:

$$\sum_{i=1}^n a_i \leq \sum_{i=1}^n b_i$$

defined for all  $n \in \mathbb{N}$ . We prove that  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction on  $n$ :

**Base Case** ( $n = 0$ ). We know that:

$$\begin{aligned} \sum_{i=1}^n a_i &= \sum_{i=1}^0 a_i \\ &= 0 \\ &\leq 0 \\ &= \sum_{i=1}^0 b_i \\ &= \sum_{i=1}^n b_i \end{aligned}$$

So the claim is true for  $n = 0$ .

**Induction Hypothesis.** Suppose that  $P(k)$  is true for some  $k \in \mathbb{N}$ .

**Induction Step.** Let the groups of numbers  $a_1, \dots, a_{k+1}$  and  $b_1, \dots, b_{k+1}$  be two groups such that  $a_i \leq b_i$  for all  $i \in [k+1]$ .

Note that

$$\begin{aligned} \sum_{i=1}^{k+1} a_i &= \sum_{i=1}^k a_i + a_{k+1} && \text{[Splitting the summation]} \\ &\leq \sum_{i=1}^k b_i + a_{k+1} && \text{[By IH]} \\ &\leq \sum_{i=1}^k b_i + b_{k+1} && \text{[By Assumption]} \\ &\leq \sum_{i=1}^{k+1} b_i && \text{[Algebra]} \end{aligned}$$

Thus we have shown that if the claim is true for  $k$ , it is true for  $k+1$ .

Therefore, we have shown  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction.

(b) For any  $n \in \mathbb{N}$ , define  $S_n$  to be the sum of the squares of the first  $n$  positive integers, or

$$S_n = \sum_{i=1}^n i^2.$$

For all  $n \in \mathbb{N}$ , prove that  $S_n = \frac{1}{6}n(n+1)(2n+1)$ .

**Solution:**

Let  $P(n)$  be the statement “ $S_n = \frac{1}{6}n(n+1)(2n+1)$ ” defined for all  $n \in \mathbb{N}$ . We prove that  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction on  $n$ .

**Base Case.** When  $n = 0$ , we know the sum of the squares of the first  $n$  positive integers is the sum of no terms, so we have a sum of 0. Thus,  $S_0 = 0$ . Since  $\frac{1}{6}(0)(0+1)((2)(0)+1) = 0$ , we know that  $P(0)$  is true.

**Induction Hypothesis.** Suppose that  $P(k)$  is true for some  $k \in \mathbb{N}$ .

**Induction Step.** Examining  $S_{k+1}$ , we see that

$$S_{k+1} = \sum_{i=1}^{k+1} i^2 = \sum_{i=1}^k i^2 + (k+1)^2 = S_k + (k+1)^2.$$

By the induction hypothesis, we know that  $S_k = \frac{1}{6}k(k+1)(2k+1)$ . Therefore, we can substitute and rewrite the expression as follows:

$$\begin{aligned} S_{k+1} &= S_k + (k+1)^2 \\ &= \frac{1}{6}k(k+1)(2k+1) + (k+1)^2 \\ &= (k+1) \left( \frac{1}{6}k(2k+1) + (k+1) \right) \\ &= \frac{1}{6}(k+1)(k(2k+1) + 6(k+1)) \\ &= \frac{1}{6}(k+1)(2k^2 + 7k + 6) \\ &= \frac{1}{6}(k+1)(k+2)(2k+3) \\ &= \frac{1}{6}(k+1)((k+1)+1)(2(k+1)+1) \end{aligned}$$

Thus, we can conclude that  $P(k+1)$  is true.

Therefore, because the base case and induction step hold,  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction.

(c) Define the triangle numbers as  $\Delta_n = 1 + 2 + \cdots + n$ , where  $n \in \mathbb{N}$ . Theorem:  $\Delta_n = \frac{n(n+1)}{2}$ .

Prove the following equality for all  $n \in \mathbb{N}$ :

$$\sum_{i=0}^n i^3 = \Delta_n^2$$

**Solution:**

First, note that  $\Delta_n = \sum_{i=0}^n i$ . So, we are trying to prove  $\sum_{i=0}^n i^3 = \left( \sum_{i=0}^n i \right)^2$ .

Let  $P(n)$  be the statement:

$$\sum_{i=0}^n i^3 = \left( \sum_{i=0}^n i \right)^2$$

We prove that  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction on  $n$ .

**Base Case.** When  $n = 0$ , we have  $0^3 = 0^2$ , so  $P(0)$  holds.

**Induction Hypothesis.** Suppose that  $P(k)$  is true for some  $k \in \mathbb{N}$ .

**Induction Step.** We show  $P(k + 1)$ :

$$\begin{aligned}
 \sum_{i=0}^{k+1} i^3 &= \sum_{i=1}^k i^3 + (k + 1)^3 && \text{[Take out a term]} \\
 &= \left( \sum_{i=0}^k i \right)^2 && \text{[Induction Hypothesis]} \\
 &= \left( \frac{k(k + 1)}{2} \right)^2 + (k + 1)^3 && \text{[Substitution from part (a)]} \\
 &= (k + 1)^2 \left( \frac{k^2}{2^2} + (k + 1) \right) && \text{[Factor } (k + 1)^2\text{]} \\
 &= (k + 1)^2 \left( \frac{k^2 + 4k + 4}{4} \right) && \text{[Add via comon denominator]} \\
 &= (k + 1)^2 \left( \frac{(k + 2)^2}{4} \right) && \text{[Factor numerator]} \\
 &= \left( \frac{(k + 1)(k + 2)}{2} \right)^2 && \text{[Take out the square]} \\
 &= \left( \sum_{i=0}^{k+1} i \right)^2 && \text{[Substitution from part (a)]}
 \end{aligned}$$

Therefore,  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction.

(d) Prove that  $9 \mid n^3 + (n + 1)^3 + (n + 2)^3$  for all  $n > 1$  by induction.

**Solution:**

Let  $P(n)$  be “ $9 \mid n^3 + (n + 1)^3 + (n + 2)^3$ ”. We will prove  $P(n)$  for all integers  $n > 1$  by induction.

**Base Case ( $n = 2$ ):**  $2^3 + (2 + 1)^3 + (2 + 2)^3 = 8 + 27 + 64 = 99 = 9 \cdot 11$ , so  $9 \mid 2^3 + (2 + 1)^3 + (2 + 2)^3$ , so  $P(2)$  holds.

**Induction Hypothesis:** Suppose that  $P(j)$  is true for some  $j > 1$ . Note that this is equivalent to assuming that  $j^3 + (j + 1)^3 + (j + 2)^3 = 9k$  for some integer  $k$ .

**Induction Step:** We want to show that  $9 \mid (j + 1)^3 + (j + 2)^3 + (j + 3)^3$ . First, note that our induction hypothesis  $P(j)$  is equivalent to assuming that  $j^3 + (j + 1)^3 + (j + 2)^3 = 9k$  for some integer  $k$ .  
Now

$$\begin{aligned}
 (j + 1)^3 + (j + 2)^3 + (j + 3)^3 &= (j + 3)^3 + 9k - j^3 \text{ for some integer } k \quad \text{[Induction Hypothesis]} \\
 &= j^3 + 9j^2 + 27j + 27 + 9k - j^3 \\
 &= 9j^2 + 27j + 27 + 9k \\
 &= 9(j^2 + 3j + 3 + k)
 \end{aligned}$$

So  $9 \mid (j + 1)^3 + (j + 2)^3 + (j + 3)^3$ , so  $P(j) \rightarrow P(j + 1)$  for some arbitrary integer  $j > 1$ . Thus,  $P(n)$  holds for all naturals  $n > 1$  by induction.

(e) Prove that  $6n + 6 < 2^n$  for all  $n \geq 6$ .

### Solution:

Let  $P(n)$  be " $6n + 6 < 2^n$ ". We will prove  $P(n)$  for all integers  $n \geq 6$  by induction.

**Base Case** ( $n = 6$ ):  $6n + 6 = 6 \cdot 6 + 6 = 42 < 64 = 2^6 = 2^n$ , so  $P(6)$  holds.

**Induction Hypothesis:** Suppose that  $P(j)$  is true for some  $j \geq 6$ .

**Induction Step:** We want to show that  $6(j + 1) + 6 < 2^{j+1}$

We have that

$$\begin{aligned} 6(j + 1) + 6 &= 6j + 6 + 6 \\ &< 2^j + 6 && \text{[Induction Hypothesis]} \\ &< 2^j + 2^j && \text{[Since } 2^j > 6, \text{ since } j \geq 6\text{]} \\ &< 2 \cdot 2^j \\ &< 2^{j+1} \end{aligned}$$

So  $P(j) \rightarrow P(j + 1)$  for an arbitrary integer  $j \geq 6$ . Therefore,  $P(n)$  holds for all naturals  $n \geq 6$  by induction.

(f) Define

$$H_i = 1 + \frac{1}{2} + \cdots + \frac{1}{i}$$

Prove that  $H_{2^n} \geq 1 + \frac{n}{2}$  for  $n \in \mathbb{N}$ .

### Solution:

We define  $H_i$  more formally as  $\sum_{k=1}^i \frac{1}{k}$ . Let  $P(n)$  be " $H_{2^n} \geq 1 + \frac{n}{2}$ ". We will prove  $P(n)$  for all  $n \in \mathbb{N}$  by induction.

**Base Case** ( $n = 0$ ):  $H_{2^0} = H_2 = H_1 = \sum_{k=1}^1 \frac{1}{k} = 1 \geq 1 + \frac{0}{2} = 1 + \frac{n}{2}$ , so  $P(0)$  holds.

**Induction Hypothesis:** Suppose  $P(j)$  is true for some  $j \in \mathbb{N}$ .

**Induction Step:** We want to show that  $H_{2^{j+1}} \geq 1 + \frac{j+1}{2}$

Now

$$\begin{aligned} H_{2^{j+1}} &= \sum_{k=1}^{2^{j+1}} \frac{1}{k} \\ &= \sum_{k=1}^{2^j} \frac{1}{k} + \sum_{k=2^j+1}^{2^{j+1}} \frac{1}{k} \\ &\geq 1 + \frac{j}{2} + \sum_{k=2^j+1}^{2^{j+1}} \frac{1}{k} && \text{[Induction Hypothesis]} \\ &\geq 1 + \frac{j}{2} + 2^j \cdot \frac{1}{2^{j+1}} && \text{[There are } 2^j \text{ terms in } [2^j + 1, 2^{j+1}] \text{ and each is at least } \frac{1}{2^{j+1}}\text{]} \\ &\geq 1 + \frac{j}{2} + \frac{2^j}{2^{j+1}} \\ &\geq 1 + \frac{j}{2} + \frac{1}{2} \\ &\geq 1 + \frac{j+1}{2} \end{aligned}$$

So  $P(j) \rightarrow P(j + 1)$  for an arbitrary integer  $j \in \mathbb{N}$ . Therefore,  $P(n)$  holds for all  $n \in \mathbb{N}$  by induction.

## 4. Set Proofs

Prove for any sets  $A$  and  $B$ ,  $P((A \cup B) \setminus B) \subseteq P(A)$ .

### Solution:

Let  $A, B$  be arbitrary sets. Let  $X$  be an arbitrary element of  $P((A \cup B) \setminus B)$ . Note that  $X \subseteq (A \cup B) \setminus B$  by the definition of powersets. Let  $x \in X$  be arbitrary. Then, by the definition of subsets  $x \in ((A \cup B) \setminus B)$ . Note that,  $x \in (A \cup B) \wedge x \notin B$  by the definition of set difference. Then,  $((x \in A \vee x \in B) \wedge x \notin B)$  by the definition of set union. Then, it follows that  $x \in A \wedge x \notin B$ . Then, it follows that  $x \in A$ . Since  $x$  was arbitrary, by the definition of subsets  $X \subseteq A$ . Then by the definition of powersets  $X \in P(A)$ . Since  $X$  was arbitrary, by the definition of subsets  $P((A \cup B) \setminus B) \subseteq P(A)$ .