Adam Blank Spring 2016



Foundations of Computing I

* All slides are a combined effort between previous instructors of the course

Building Precedence in Arithmetic Expressions

- E expression (start symbol)
- T term F factor I identifier N number

 $E \rightarrow T \mid E+T$

 $T \rightarrow F \mid F*T$

 $F \rightarrow (E) \mid I \mid N$

 $I \rightarrow x \mid y \mid z$

 $N \rightarrow 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9$

Backus-Naur Form (The same thing...)

BNF (Backus-Naur Form) grammars

- Originally used to define programming languages
- Variables denoted by long names in angle brackets, e.g.

<identifier>, <if-then-else-statement>, <assignment-statement>, <condition>

::= used instead of \rightarrow

BNF for C

Parse Trees

Back to middle school:

<sentence>::=<noun phrase><verb phrase>

<noun phrase>::==<article><adjective><noun>

<verb phrase>::=<verb><adverb>|<verb><object>

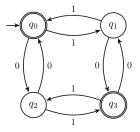
<object>::=<noun phrase>

Parse:

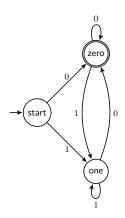
The yellow duck squeaked loudly The red truck hit a parked car

CSE 311: Foundations of Computing

Lecture 20: Finite State Machines (DFAs)



A Weird Sort of Programming!

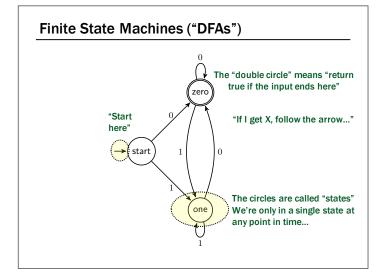


What does this "thing" do?

Take a guess!

If you had to give this "method" a name, what would it be?

boolean isEven(binary s)



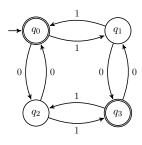
Applications of FSMs (a.k.a. finite automata)

- Implementation of regular expression matching in programs like grep
- Control structures for sequential logic in digital circuits
- Algorithms for communication and cachecoherence protocols
 - Each agent runs its own FSM
- · Design specifications for reactive systems
 - Components are communicating FSMs

Applications of FSMs (a.k.a. finite automata)

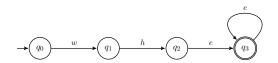
- · Formal verification of systems
 - Is an unsafe state reachable?
- · Computer games
 - FSMs provide worlds to explore
- Minimization algorithms for FSMs can be extended to more general models used in
 - Text prediction
 - Speech recognition

What language does this machine recognize?



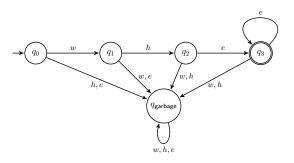
All binary strings with even length

Why is this not a DFA? Fix it!



DFAs must have a transition for every character at every state!

Why is this not a DFA? Fix it!



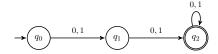
"Garbage states" are a useful concept. Whenever we KNOW we can't accept the string, just send it to a state that will always go back to itself. This is the way of saying "return false" in DFA-land.

For each of the following languages, create a DFA

Ø

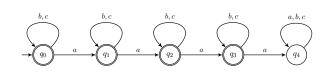
 ${x \in {0,1}}^* : len(x) > 1}$





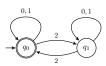


FSM that accepts strings of a's, b's, c's with no more than 3 a's



Strings over {0, 1, 2}*

M₁: Strings with an even number of 2's



 M_2 : Strings where the sum of digits mod 3 is 0

