# **CSE 311: Foundations of Computing**

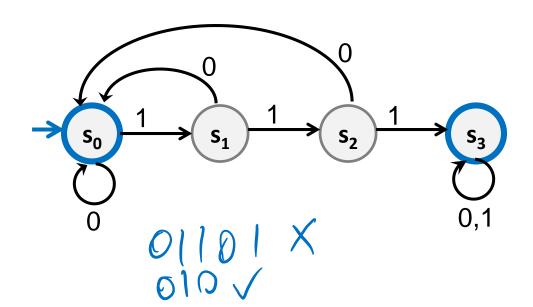
**Lecture 22: DFAs and Finite State Machines with Output** 



#### **Finite State Machines**

- States
- Transitions on input symbols
- Start state and final states
- The "language recognized" by the machine is the set of strings that reach a final state from the start

Old State	0	1
s <sub>0</sub>	s <sub>0</sub>	S <sub>1</sub>
S <sub>1</sub>	s <sub>0</sub>	S <sub>2</sub>
S <sub>2</sub>	s <sub>0</sub>	S <sub>3</sub>
S <sub>3</sub>	S <sub>3</sub>	S <sub>3</sub>

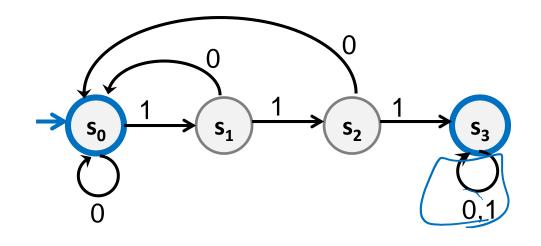


#### **Finite State Machines**

• Each machine designed for strings over some fixed alphabet  $\Sigma$ .

 Must have a transition defined from each state for every symbol in Σ.

Old State	0	1
s <sub>0</sub>	s <sub>0</sub>	S <sub>1</sub>
S <sub>1</sub>	$s_0$	S <sub>2</sub>
S <sub>2</sub>	$s_0$	S <sub>3</sub>
S <sub>3</sub>	<b>S</b> <sub>3</sub>	S <sub>3</sub>



# Applications of FSMs (a.k.a. Finite Automata)

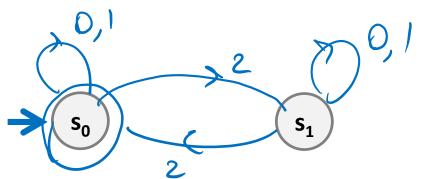
- Implementation of regular expression matching in programs like grep
- Control structures for sequential logic in digital circuits
- Algorithms for communication and cachecoherence protocols
  - Each agent runs its own FSM
- Design specifications for reactive systems
  - Components are communicating FSMs

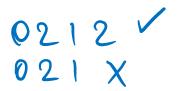
# Applications of FSMs (a.k.a. Finite Automata)

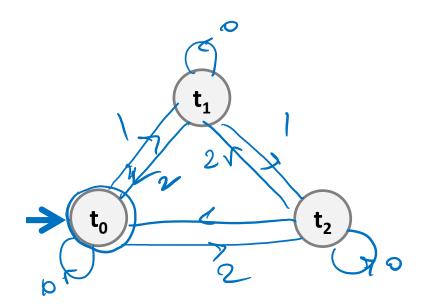
- Formal verification of systems
  - Is an unsafe state reachable?
- Computer games
  - FSMs provide worlds to explore
- Minimization algorithms for FSMs can be extended to more general models used in
  - Text prediction
  - Speech recognition

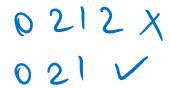
# **Strings over** {0, 1, 2}

#### M<sub>1</sub>: Strings with an even number of 2's



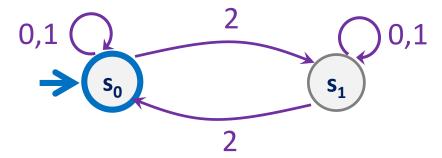


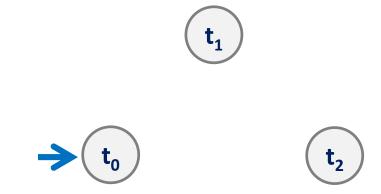




# **Strings over** {0, 1, 2}

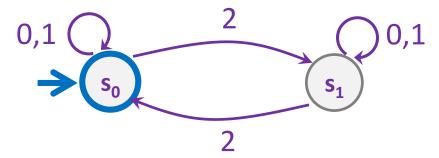
M<sub>1</sub>: Strings with an even number of 2's

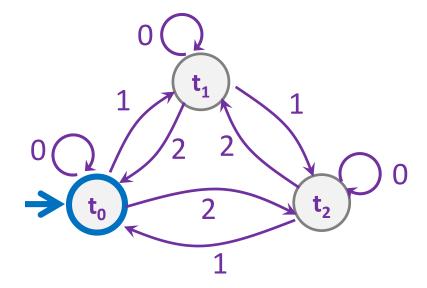




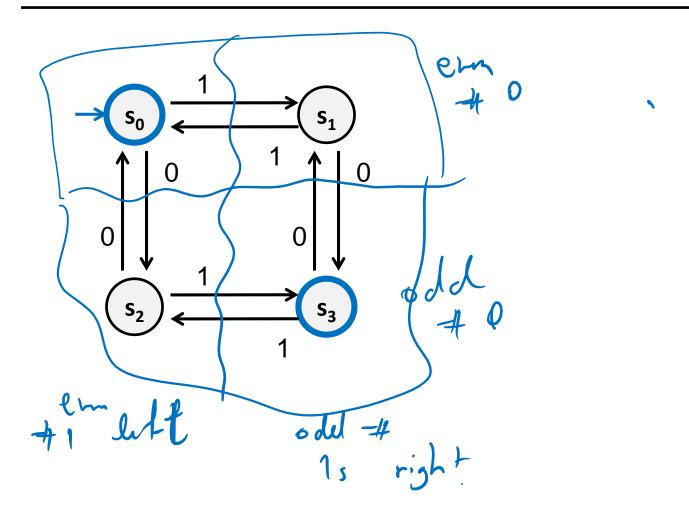


M<sub>1</sub>: Strings with an even number of 2's

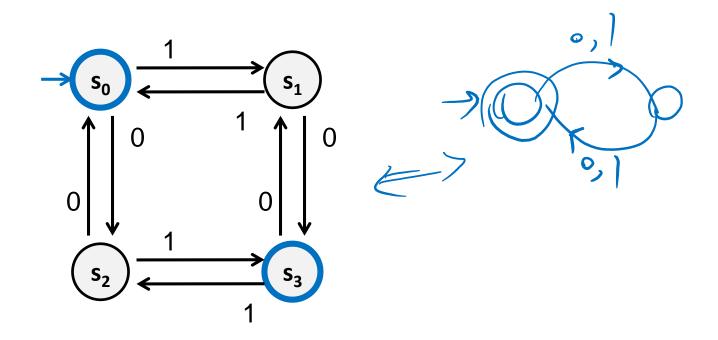




# What language does this machine recognize?



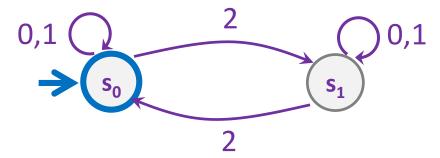
# What language does this machine recognize?

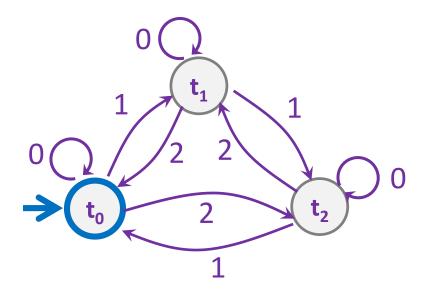


The set of all binary strings with # of 1's  $\equiv$  # of 0's (mod 2) (both are even or both are odd).

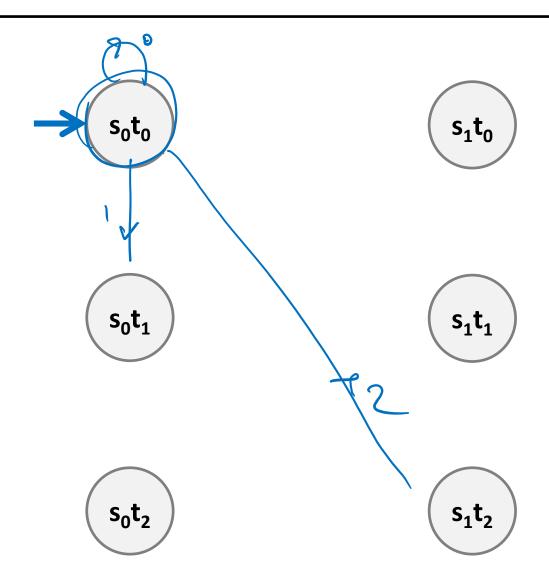
# **Strings over** {0, 1, 2}

# M<sub>1</sub>: Strings with an even number of 2's

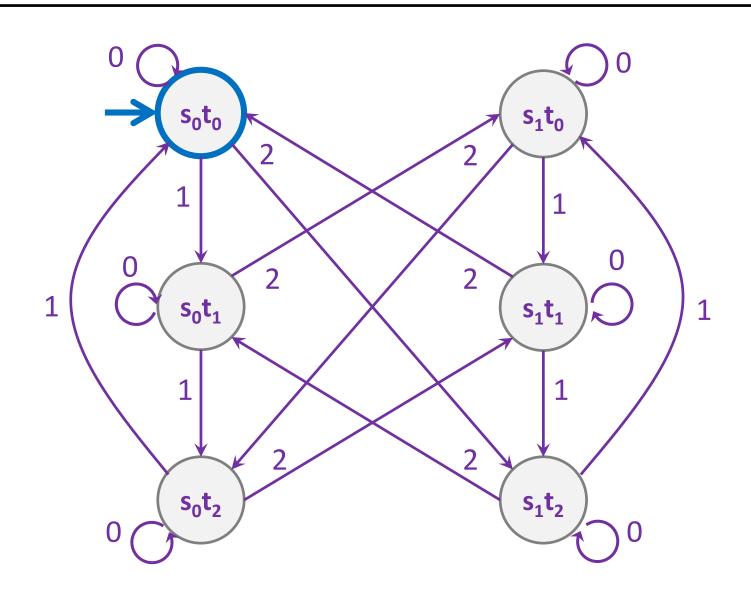




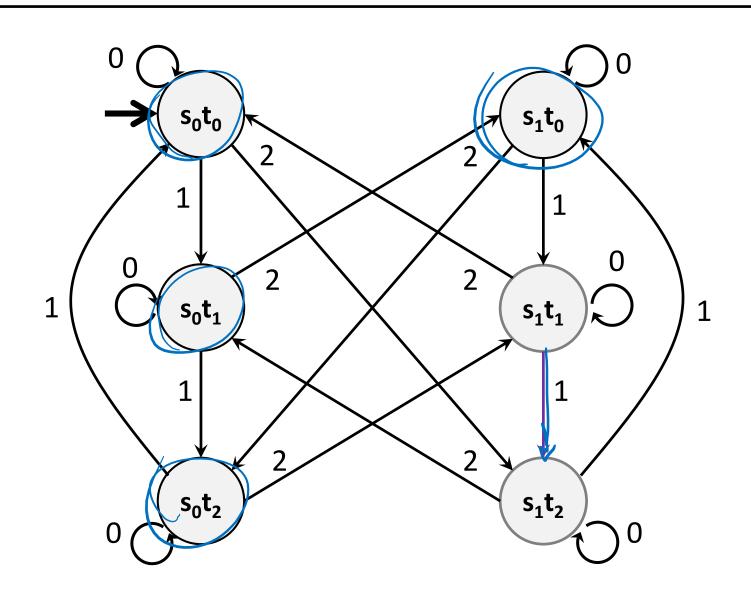
## Strings over {0,1,2} w/ even number of 2's and mod 3 sum 0



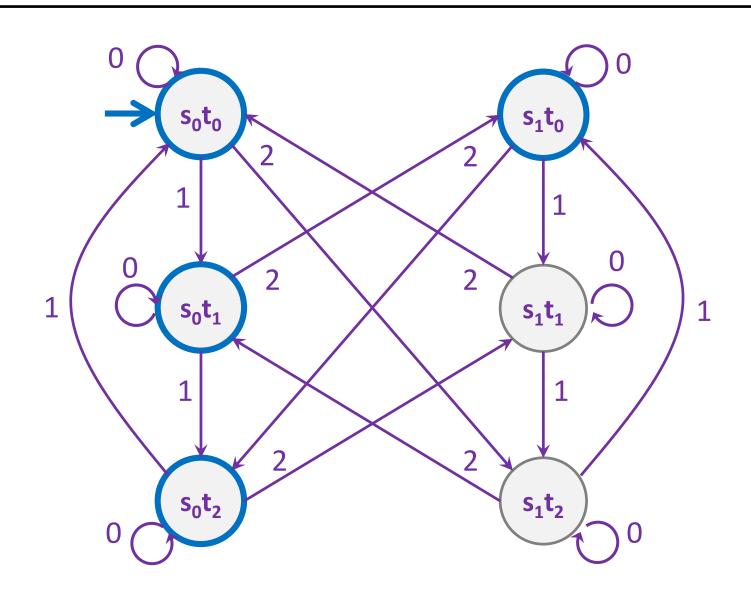
Strings over {0,1,2} w/ even number of 2's and mod 3 sum 0



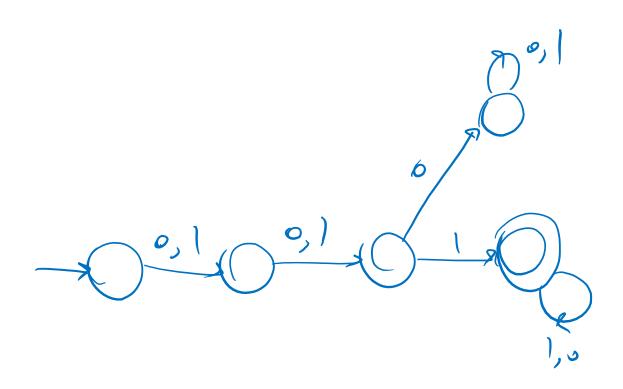
Strings over {0,1,2} w/ even number of 2's OR mod 3 sum 0?



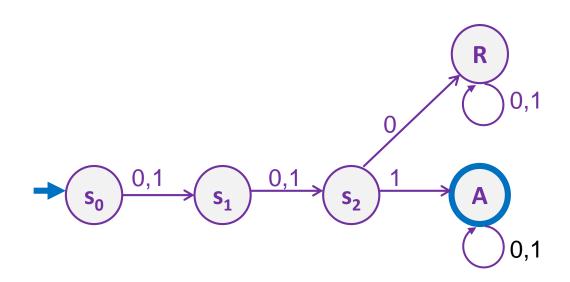
Strings over {0,1,2} w/ even number of 2's OR mod 3 sum 0



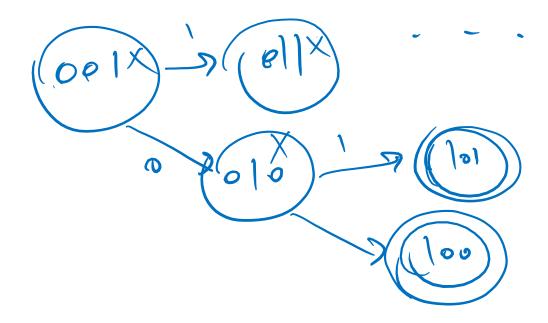
#### The set of binary strings with a 1 in the 3<sup>rd</sup> position from the start



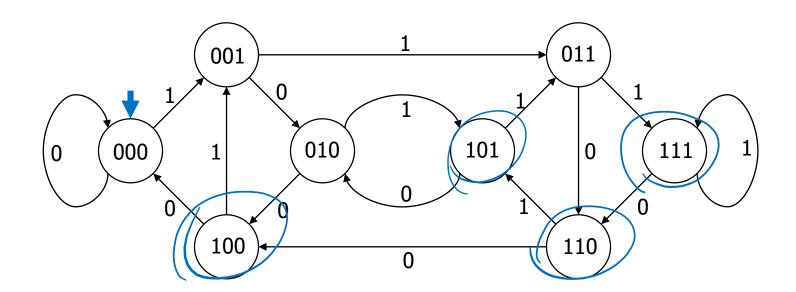
#### The set of binary strings with a 1 in the 3<sup>rd</sup> position from the start



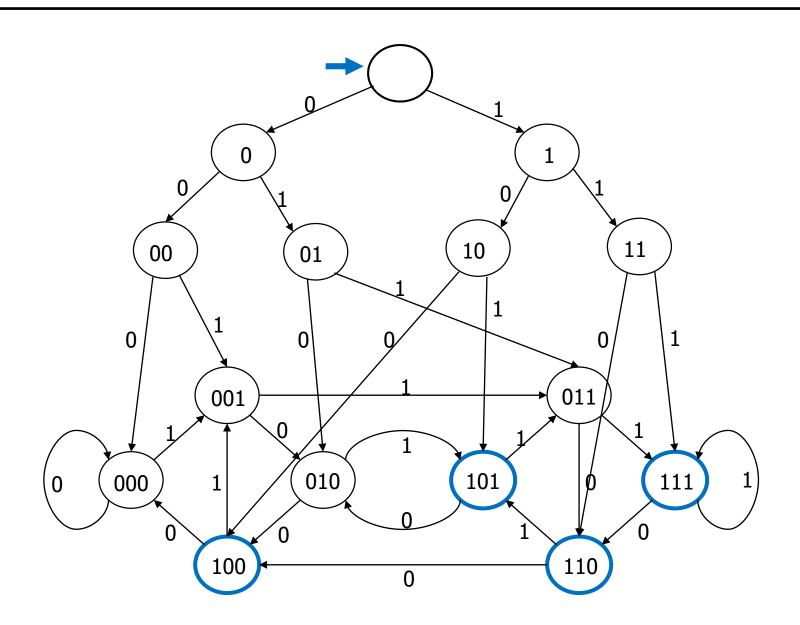
#### The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end



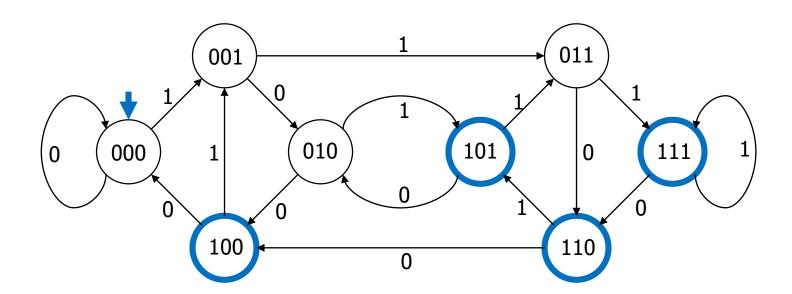
# 3 bit shift register "Remember the last three bits"



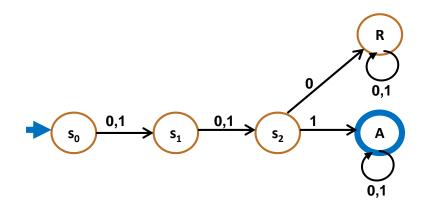
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

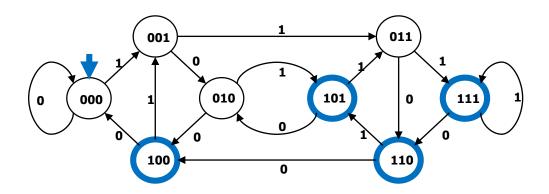


#### The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end



# The beginning versus the end





# **Adding Output to Finite State Machines**

- So far we have considered finite state machines that just accept/reject strings
  - called "Deterministic Finite Automata" or DFAs

- Now we consider finite state machines that with output
  - These are the kinds used as controllers



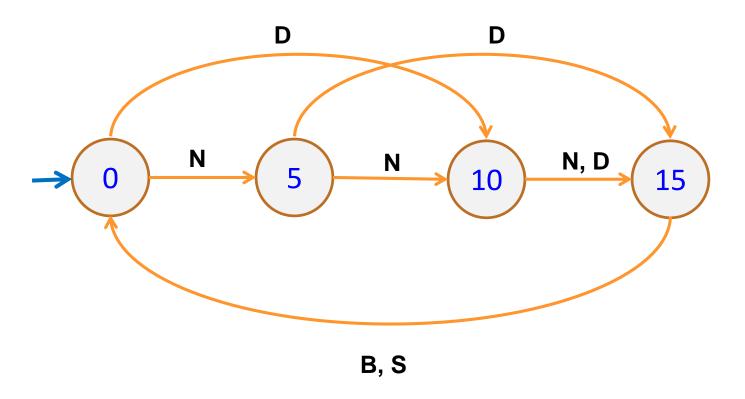
# **Vending Machine**



# Enter 15 cents in dimes or nickels Press S or B for a candy bar

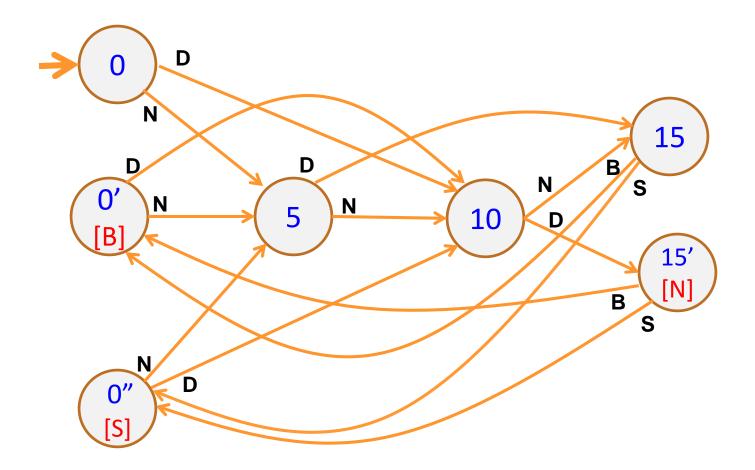


# **Vending Machine, v0.1**



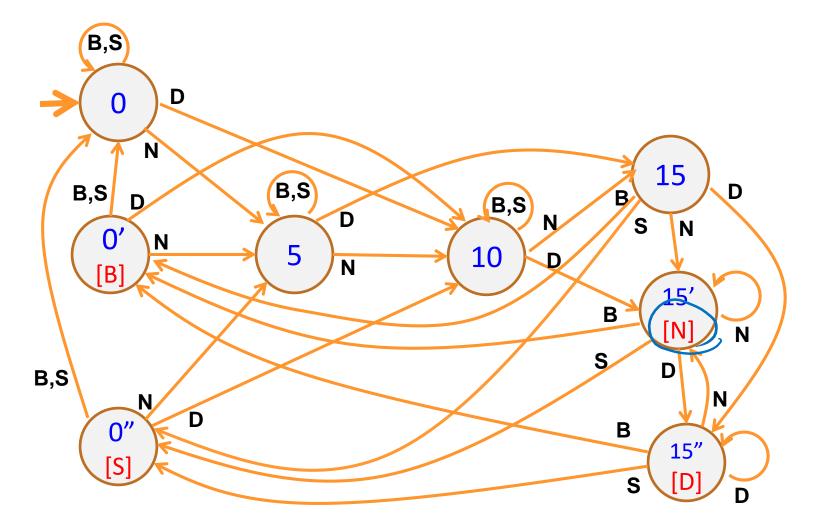
Basic transitions on **N** (nickel), **D** (dime), **B** (butterfinger), **S** (snickers)

# **Vending Machine, v0.2**



Adding output to states: N – Nickel, S – Snickers, B – Butterfinger

# Vending Machine, v1.0



Adding additional "unexpected" transitions to cover all symbols for each state