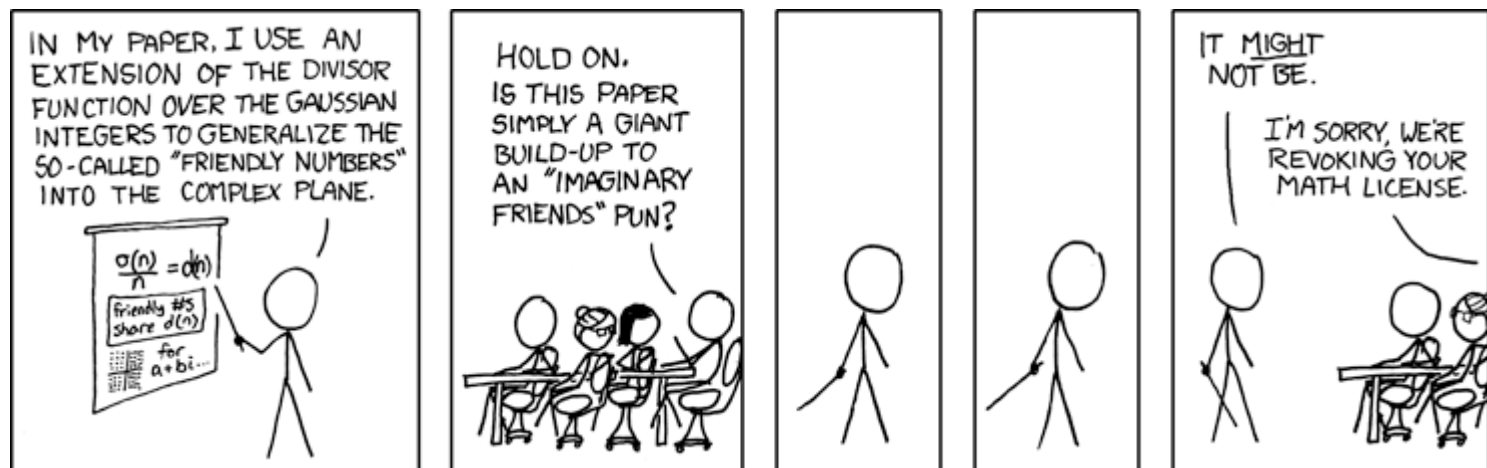


CSE 311: Foundations of Computing

Lecture 11: Modular Arithmetic and Applications



Divisibility

$$a \mid b \iff \exists k \quad b = ka$$

Definition: “a divides b”

For $a \in \mathbb{Z}, b \in \mathbb{Z}$ with $a \neq 0$:

$$a \mid b \iff \exists k \in \mathbb{Z} \quad (b = ka)$$

Check Your Understanding. Which of the following are true?

$$5 \mid 1$$

$$5 \mid 1 \text{ iff } 1 = 5k$$

$$25 \mid 5$$

$$25 \mid 5 \text{ iff } 5 = 25k$$

$$5 \mid 0$$

$$5 \mid 0 \text{ iff } 0 = 5k$$

$$3 \mid 2$$

$$3 \mid 2 \text{ iff } 2 = 3k$$

$$1 \mid 5$$

$$1 \mid 5 \text{ iff } 5 = 1k$$

$$5 \mid 25$$

$$5 \mid 25 \text{ iff } 25 = 5k$$

$$0 \mid 5$$

$$0 \mid 5 \text{ iff } 5 = 0k$$

$$2 \mid 3$$

$$2 \mid 3 \text{ iff } 3 = 2k$$

Division Theorem

$$\begin{array}{l} a = -17 \quad d = 4 \\ q = -5 \quad r = 3 \quad -17 = d \cdot (-5) + 3 \end{array}$$

Division Theorem

For $a \in \mathbb{Z}, d \in \mathbb{Z}$ with $d > 0$

there exist *unique* integers q, r with $0 \leq r < d$
such that $a = dq + r$.

To put it another way, if we divide d into a , we get a
unique quotient $q = a \text{ div } d$
and non-negative remainder $r = a \text{ mod } d$

$$\begin{array}{l} a = dq + r \\ \swarrow \quad \searrow \\ 0 \leq r \leq d - 1 \end{array}$$

$$\begin{array}{l} a = +17 \quad d = 4 \\ +17 = d \cdot \underbrace{4}_q + \underbrace{1}_r \end{array}$$

Note: $r \geq 0$ even if $a < 0$.
Not quite the same as $a \% d$.

Division Theorem

Division Theorem

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there exist *unique* integers q, r with $0 \leq r < d$
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To put it another way, if we divide d into a , we get a
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```
public class Test2 {  
    public static void main(String args[]) {  
        int a = -5;  
        int d = 2;  
        System.out.println(a % d);  
    }  
}
```

```
----jGRASP exec: java Test2  
-1  
----jGRASP: operation complete.
```

Note: $r \geq 0$ even if $a < 0$.
Not quite the same as $a \% d$.

Arithmetic, mod 7

$$a +_7 b = (a + b) \bmod 7$$

$$a \times_7 b = (a \times b) \bmod 7$$

$$6 \times 3 = 4$$

$$5 + 4 = 2$$

$$2 \times 6 = 5$$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Modular Arithmetic

Definition: “a is congruent to b modulo m”

For $a, b, m \in \mathbb{Z}$ with $m > 0$

$$a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$$

Check Your Understanding. What do each of these mean?
When are they true?

$$x \equiv 0 \pmod{2} \quad 2 \mid x - 0 \quad 2 \mid x$$

x is even.

$$-1 \equiv 19 \pmod{5} \quad 5 \mid -1 - 19 = -20 \quad \text{true}$$

$$y \equiv 2 \pmod{7} \quad 7 \mid y - 2 \quad y = 7k + 2$$

Modular Arithmetic

Definition: “a is congruent to b modulo m”

For $a, b, m \in \mathbb{Z}$ with $m > 0$

$$a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$$

**Check Your Understanding. What do each of these mean?
When are they true?**

$$x \equiv 0 \pmod{2}$$

This statement is the same as saying “x is even”; so, any x that is even (including negative even numbers) will work.

$$-1 \equiv 19 \pmod{5}$$

This statement is true. $19 - (-1) = 20$ which is divisible by 5

$$y \equiv 2 \pmod{7}$$

This statement is true for y in $\{ \dots, -12, -5, 2, 9, 16, \dots \}$. In other words, all y of the form $2+7k$ for k an integer.

Modular Arithmetic: A Property

Let a, b, m be integers with $m > 0$.

Then, $a \equiv b \pmod{m}$ if and only if $a \bmod m = b \bmod m$.

Suppose that $a \equiv b \pmod{m}$.

$m \mid a - b$ by definition of congruence \Rightarrow
So $a - b = km$ for some integer k .
 $a = b + km$. Take both sides mod m .
 $a \bmod m = (b + km) \bmod m = b \bmod m$.

Conclusion $a \bmod m = b \bmod m$.

Suppose that $a \bmod m = b \bmod m$.

$a = m(a \text{ div } m) + a \bmod m$ by division Thm.
 $b = m(b \text{ div } m) + b \bmod m$ by division Thm.
 $a - b = m(a \text{ div } m - b \text{ div } m) + a \bmod m - b \bmod m = m(a \text{ div } m - b \text{ div } m)$.
 $a - b = km$ for some k .
 $m \mid a - b$

Conclusion $a \equiv b \pmod{m}$.

Modular Arithmetic: A Property

Let a, b, m be integers with $m > 0$.

Then, $a \equiv b \pmod{m}$ if and only if $a \bmod m = b \bmod m$.

Suppose that $a \equiv b \pmod{m}$.

Then, $m \mid (a - b)$ by definition of congruence.

So, $a - b = km$ for some integer k by definition of divides.

Therefore, $a = b + km$.

Taking both sides modulo m we get:

$$a \bmod m = (b + km) \bmod m = b \bmod m.$$

Suppose that $a \bmod m = b \bmod m$.

By the division theorem, $a = mq + (a \bmod m)$ and

$$b = ms + (b \bmod m) \text{ for some integers } q, s.$$

$$\text{Then, } a - b = (mq + (a \bmod m)) - (ms + (b \bmod m))$$

$$= m(q - s) + (a \bmod m - b \bmod m)$$

$$= m(q - s) \text{ since } a \bmod m = b \bmod m$$

Therefore, $m \mid (a - b)$ and so $a \equiv b \pmod{m}$.

The mod m function vs the $\equiv (\text{mod } m)$ predicate

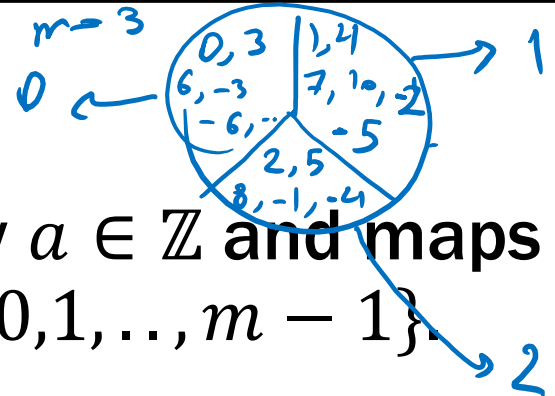
- What we have just shown

- The mod m function takes any $a \in \mathbb{Z}$ and maps it to a remainder $a \bmod m \in \{0, 1, \dots, m - 1\}$.
- Imagine grouping together all integers that have the same value of the mod m function

That is, the same remainder in $\{0, 1, \dots, m - 1\}$.

- The $\equiv (\text{mod } m)$ predicate compares $a, b \in \mathbb{Z}$. It is true if and only if the mod m function has the same value on a and on b .

That is, a and b are in the same group.



Modular Arithmetic: Addition Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $a + c \equiv b + d \pmod{m}$

pf. $m \mid a - b$ by def of cong
 $a - b = km$ for some int k
+ $m \mid c - d$ by def of cong
 $c - d = lm$ by def of div for some l .
 $a + c - (b + d) = \underbrace{(k + l)}_n \cdot m$ by algebra.

$$m \cdot n = a + c - (b + d) \text{ for some } n.$$

$$m \mid a + c - (b + d)$$
$$\vdash a + c \equiv b + d \pmod{m}$$

Modular Arithmetic: Addition Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $a + c \equiv b + d \pmod{m}$

Suppose that $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$. Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Adding the equations together gives us $(a + c) - (b + d) = m(k + j)$. Now, re-applying the definition of congruence gives us $a + c \equiv b + d \pmod{m}$.

Modular Arithmetic: Multiplication Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $ac \equiv bd \pmod{m}$

pf. $m \mid a - b$ by def of cong
 $\underline{a - b = k \cdot m}$ for some int k .
 $m \mid c - d$ def of cong. $\underline{c - d = l \cdot m}$ for some int l .
 $a = b + km$
 $c = d + lm$ multiply $ac = bd + b \cdot lm + d \cdot km + k \cdot l \cdot m^2$
 $ac - bd = m(\underbrace{bl + dk + klm}_n)$

$$ac - bd = n \cdot m \text{ for some int } n.$$

$$m \mid ac - bd$$

$$ac \equiv bd \pmod{m}$$

Modular Arithmetic: Multiplication Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $ac \equiv bd \pmod{m}$

Suppose that $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$. Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Then, $a = km + b$ and $c = jm + d$. Multiplying both together gives us $ac = (km + b)(jm + d) = kjm^2 + kmd + bjm + bd$.

Re-arranging gives us $ac - bd = m(kjm + kd + bj)$.

Using the definition of congruence gives us $ac \equiv bd \pmod{m}$.

Example

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

Example

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$

Case 1 (n is even):

$$n = 2k \text{ for some int } k.$$

$$n^2 = 4k^2$$

$$4 \mid n^2$$

$$n^2 \equiv 0 \pmod{4}.$$

Case 2 (n is odd):

$$n = 2k+1 \text{ for some int } k$$

$$n^2 = (2k+1)^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$$

$$n^2 - 1 = 4(k^2 + k)$$

$$4 \mid n^2 - 1$$

$$n^2 \equiv 1 \pmod{4}.$$

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

It looks like

$$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}, \text{ and}$$

$$n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}.$$

Example

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$

Case 1 (n is even):

Suppose $n \equiv 0 \pmod{2}$.

Then, $n = 2k$ for some integer k .

So, $n^2 = (2k)^2 = 4k^2$. So, by

definition of congruence,

$n^2 \equiv 0 \pmod{4}$.

Let's start by looking at a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

It looks like

Case 2 (n is odd):

Suppose $n \equiv 1 \pmod{2}$.

Then, $n = 2k + 1$ for some integer k .

So, $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$.

So, by definition of congruence, $n^2 \equiv 1 \pmod{4}$.

$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$, and

$n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$.

n-bit Unsigned Integer Representation

- Represent integer x as sum of powers of 2:

If $\sum_{i=0}^{n-1} b_i 2^i$ where each $b_i \in \{0,1\}$

then representation is $b_{n-1} \dots b_2 b_1 b_0$

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

- For $n = 8$:

99: 0110 0011

18: 0001 0010

Sign-Magnitude Integer Representation

n-bit signed integers

Suppose that $-2^{n-1} < x < 2^{n-1}$

First bit as the sign, $n - 1$ bits for the value

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

For $n = 8$:

99: 0110 0011

-18: 1001 0010

Any problems with this representation?

Two's Complement Representation

n bit signed integers, first bit will still be the sign bit

Suppose that $0 \leq x < 2^{n-1}$,

x is represented by the binary representation of x

Suppose that $0 \leq x \leq 2^{n-1}$,

$-x$ is represented by the binary representation of $2^n - x$

Key property: Two's complement representation of any number y is equivalent to $y \bmod 2^n$ so arithmetic works **mod 2^n**

$$99 = 64 + 32 + 2 + 1$$

$$18 = 16 + 2$$

For $n = 8$:

99: 0110 0011

-18: 1110 1110

Sign-Magnitude vs. Two's Complement

-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
1111	1110	1101	1100	1011	1010	1001	0000	0001	0010	0011	0100	0101	0110	0111

Sign-bit

-8	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
1000	1001	1010	1011	1100	1101	1110	1111	0000	0001	0010	0011	0100	0101	0110	0111

Two's complement

Two's Complement Representation

- For $0 < x \leq 2^{n-1}$, $-x$ is represented by the binary representation of $2^n - x$
 - That is, the two's complement representation of any number y has the same value as y modulo 2^n .
- To compute this: Flip the bits of x then add 1:
 - All 1's string is $2^n - 1$, so
Flip the bits of $x \equiv$ replace x by $2^n - 1 - x$
Then add 1 to get $2^n - x$

Basic Applications of mod

- Hashing
- Pseudo random number generation
- Simple cipher

Hashing

Scenario:

Map a small number of data values from a large domain $\{0, 1, \dots, M - 1\}$...

...into a small set of locations $\{0, 1, \dots, n - 1\}$ so one can quickly check if some value is present

- $\text{hash}(x) = x \bmod p$ for p a prime close to n
 - or $\text{hash}(x) = (ax + b) \bmod p$
- Depends on all of the bits of the data
 - helps avoid collisions due to similar values
 - need to manage them if they occur

Pseudo-Random Number Generation

Linear Congruential method

$$x_{n+1} = (a x_n + c) \bmod m$$

Choose random x_0, a, c, m and produce a long sequence of x_n 's

Simple Ciphers

- **Caesar cipher**, $A = 1$, $B = 2$, . . .
 - HELLO WORLD
- **Shift cipher**
 - $f(p) = (p + k) \bmod 26$
 - $f^{-1}(p) = (p - k) \bmod 26$
- **More general**
 - $f(p) = (ap + b) \bmod 26$