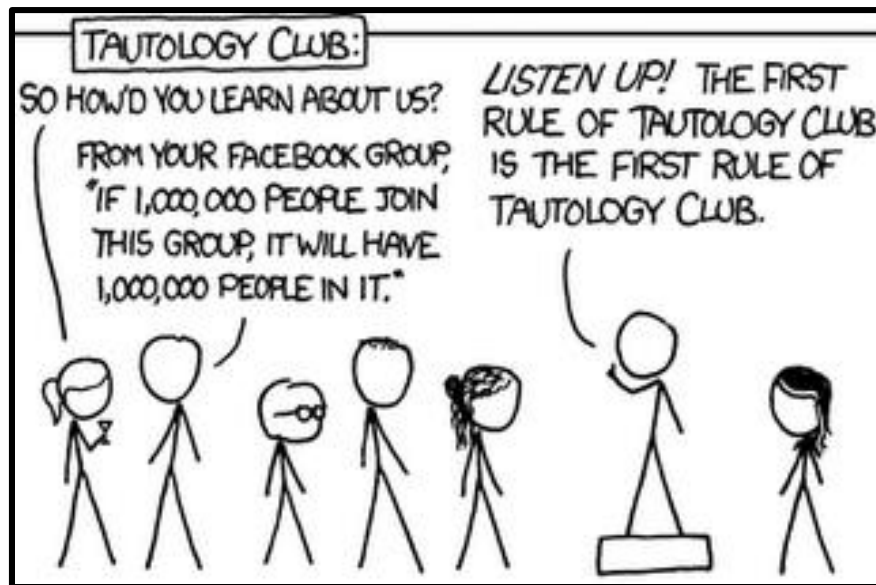


CSE 311: Foundations of Computing

Lecture 6: DNF, CNF and Predicate Logic



1-bit Binary Adder

A	$0 + 0 = 0$ (with $C_{OUT} = 0$)
<u>+ B</u>	$0 + 1 = 1$ (with $C_{OUT} = 0$)
S	$1 + 0 = 1$ (with $C_{OUT} = 0$)
(C_{OUT})	$1 + 1 = 0$ (with $C_{OUT} = 1$)



.

1-bit Binary Adder

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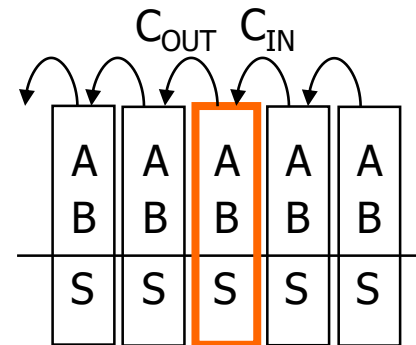
Idea: To chain these together, let's add a carry-in

1-bit Binary Adder

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Idea: To chain these together, let's add a carry-in

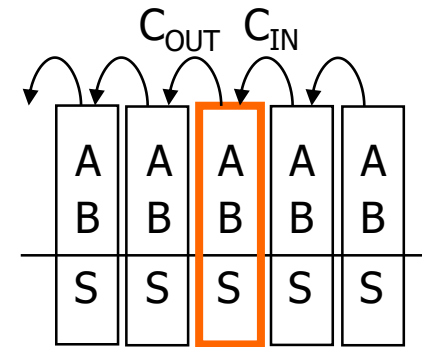
$$\begin{array}{r} (C_{IN}) \\ A \\ + B \\ \hline S \\ (C_{OUT}) \end{array}$$



$$\begin{array}{r} 1101 \\ 0101 \\ 1101 \\ \hline 10016 \end{array}$$

1-bit Binary Adder

- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



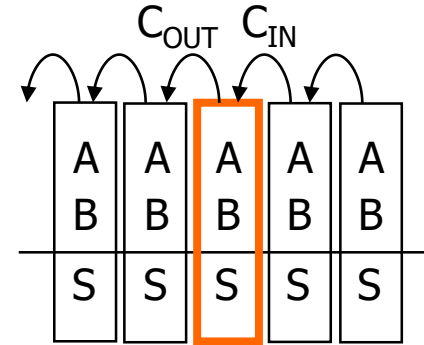
A	B	C _{IN}	C _{OUT}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$\begin{array}{r} A' \cdot B' \cdot C_{in} \\ + \\ \\ + \\ \\ + \end{array}$$



1-bit Binary Adder

- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



A	B	C _{IN}	C _{OUT}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Derive an expression for S

$$A' \cdot B' \cdot C_{IN}$$

$$A' \cdot B \cdot C_{IN}'$$

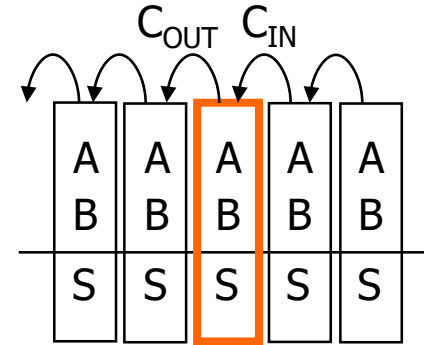
$$A \cdot B' \cdot C_{IN}'$$

$$A \cdot B \cdot C_{IN}$$

$$S = A' \cdot B' \cdot C_{IN} + A' \cdot B \cdot C_{IN}' + A \cdot B' \cdot C_{IN}' + A \cdot B \cdot C_{IN}$$

1-bit Binary Adder

- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



A	B	C _{IN}	C _{OUT}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Derive an expression for C_{OUT}

$$A' \cdot B \cdot C_{IN}$$

$$A \cdot B' \cdot C_{IN}$$

$$A \cdot B \cdot C_{IN}'$$

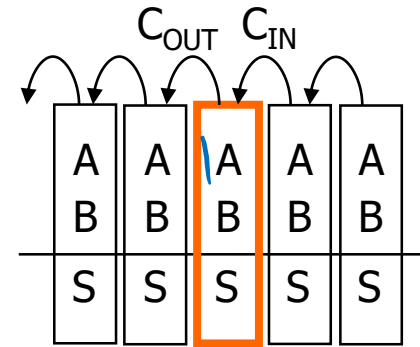
$$A \cdot B \cdot C_{IN}$$

$$C_{OUT} = A' \cdot B \cdot C_{IN} + A \cdot B' \cdot C_{IN} + A \cdot B \cdot C_{IN}' + A \cdot B \cdot C_{IN}$$

$$S = A' \cdot B' \cdot C_{IN} + A' \cdot B \cdot C_{IN}' + A \cdot B' \cdot C_{IN}' + A \cdot B \cdot C_{IN}$$

1-bit Binary Adder

- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



A	B	C _{IN}	C _{OUT}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$S = A' \cdot B' \cdot C_{IN} + A' \cdot B \cdot C_{IN}' + A \cdot B' \cdot C_{IN}' + A \cdot B \cdot C_{IN}$$

$$C_{OUT} = A' \cdot B \cdot C_{IN} + A \cdot B' \cdot C_{IN} + A \cdot B \cdot C_{IN}' + A \cdot B \cdot C_{IN}$$

Apply Theorems to Simplify Expressions

The theorems of Boolean algebra can simplify expressions

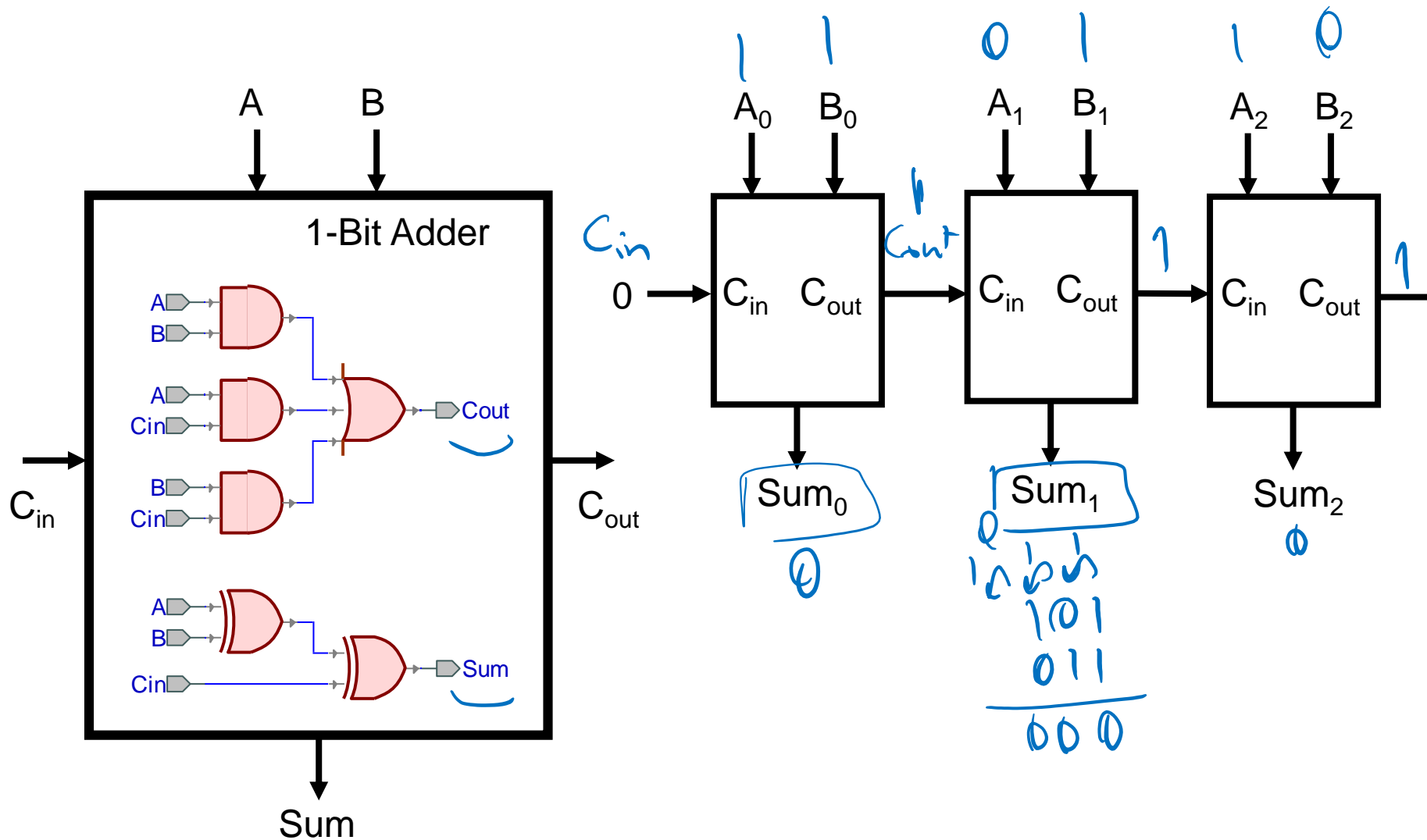
– e.g., full adder's carry-out function

$$p \vee q \equiv p \vee q \vee q$$

$$\begin{aligned}\text{Cout} &= A' B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= A' B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + \boxed{A B \text{Cin}} + \boxed{A B \text{Cin}} \\ &= A' B \text{Cin} + A B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= (A' + A) B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= (1) B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + \boxed{A B \text{Cin}} + \boxed{A B \text{Cin}} \\ &= B \text{Cin} + A B' \text{Cin} + A B \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= B \text{Cin} + A (B' + B) \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= B \text{Cin} + A (1) \text{Cin} + A B \text{Cin}' + A B \text{Cin} \\ &= B \text{Cin} + A \text{Cin} + A B (\text{Cin}' + \text{Cin}) \\ &= B \text{Cin} + A \text{Cin} + A B (1) \\ &= \boxed{B \text{Cin} + A \text{Cin} + A B}\end{aligned}$$

adding extra terms
creates new factoring
opportunities

A 2-bit Ripple-Carry Adder



Mapping Truth Tables to Logic Gates

Given a truth table:

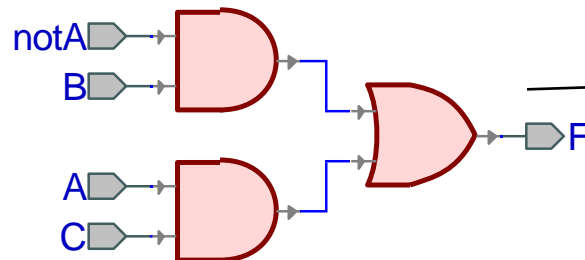
1. Write the Boolean expression
2. Minimize the Boolean expression
3. Draw as gates
4. Map to available gates

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

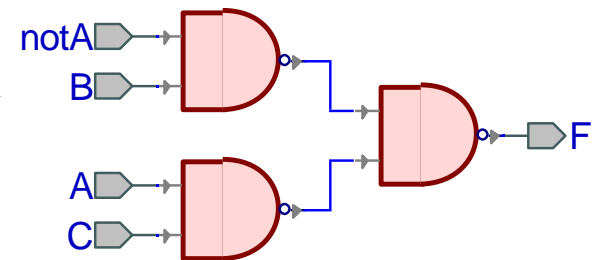
(2) ↓

$$\begin{aligned} F &= A'BC' + A'BC + AB'C + ABC \\ &= A'B(C' + C) + AC(B' + B) \\ &= A'B + AC \end{aligned}$$

(3) ↘



(4) →



Canonical Forms

- **Truth table is the unique signature of a Boolean Function**
- **The same truth table can have many gate realizations**
 - We've seen this already
 - Depends on how good we are at Boolean simplification
- **Canonical forms**
 - Standard forms for a Boolean expression
 - We all come up with the same expression

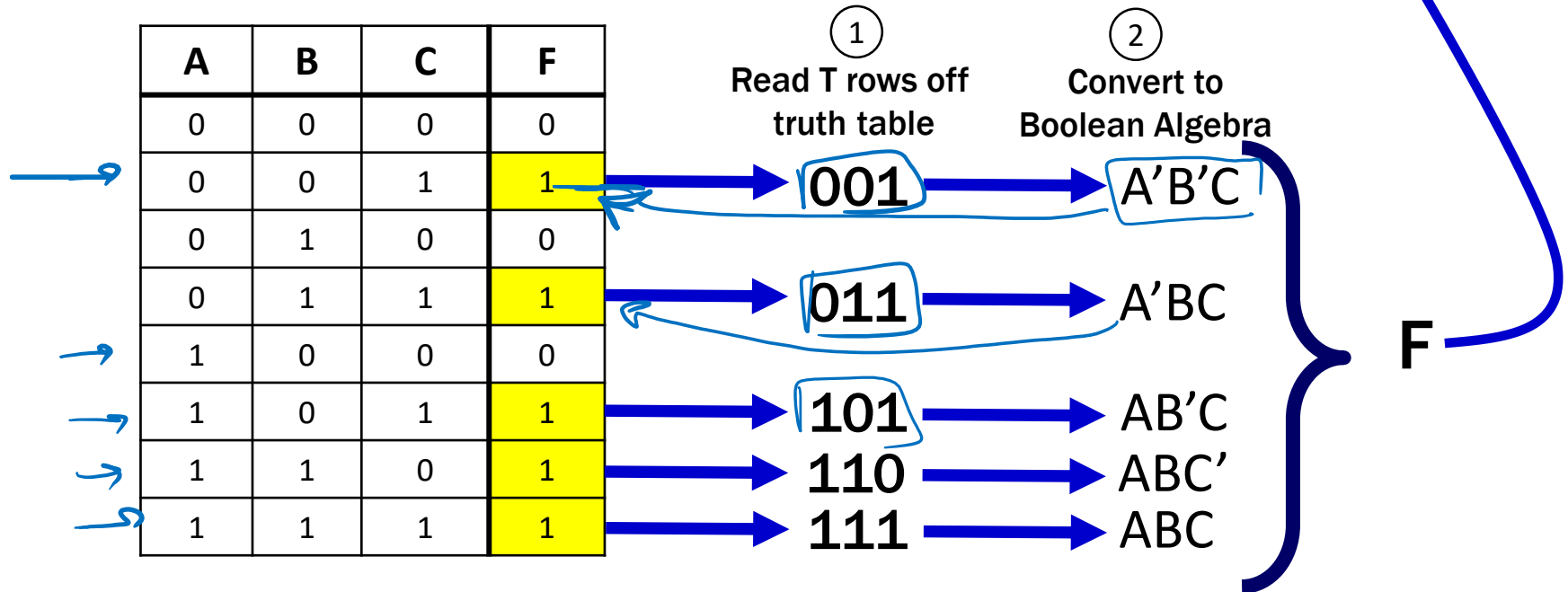
Sum-of-Products Canonical Form

- AKA **Disjunctive Normal Form (DNF)**
- AKA **Minterm Expansion**

min

③ Add the minterms together

$$F = A'B'C \text{ (min)} + A'BC + AB'C + ABC' + ABC$$



Sum-of-Products Canonical Form

Product term (or minterm)

- ANDed product of literals – input combination for which output is true
- each variable appears exactly once, true or inverted (but not both)

A	B	C	minterms
0	0	0	$A'B'C'$
0	0	1	$A'B'C$
0	1	0	$A'BC'$
0	1	1	$A'BC$
1	0	0	$AB'C'$
1	0	1	$AB'C$
1	1	0	ABC'
1	1	1	ABC

F in canonical form:

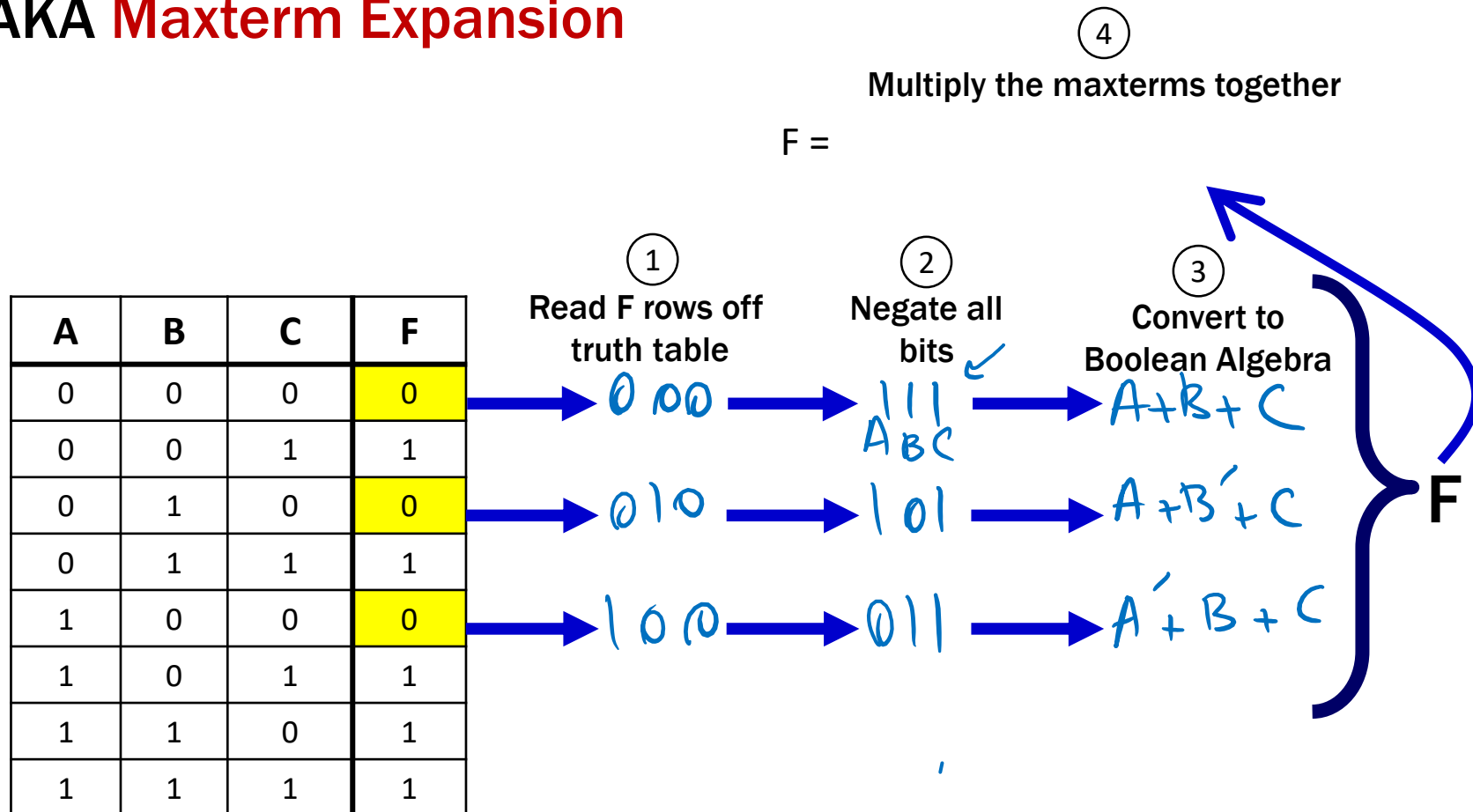
$$F(A, B, C) = A'B'C + A'BC + AB'C + ABC' + ABC$$

canonical form \neq minimal form

$$\begin{aligned} F(A, B, C) &= A'B'C + A'BC + AB'C + ABC + ABC' \\ &= (A'B' + A'B + AB' + AB)C + ABC' \\ &= ((A' + A)(B' + B))C + ABC' \\ &= C + ABC' \\ &= ABC' + C \\ &= AB + C \end{aligned}$$

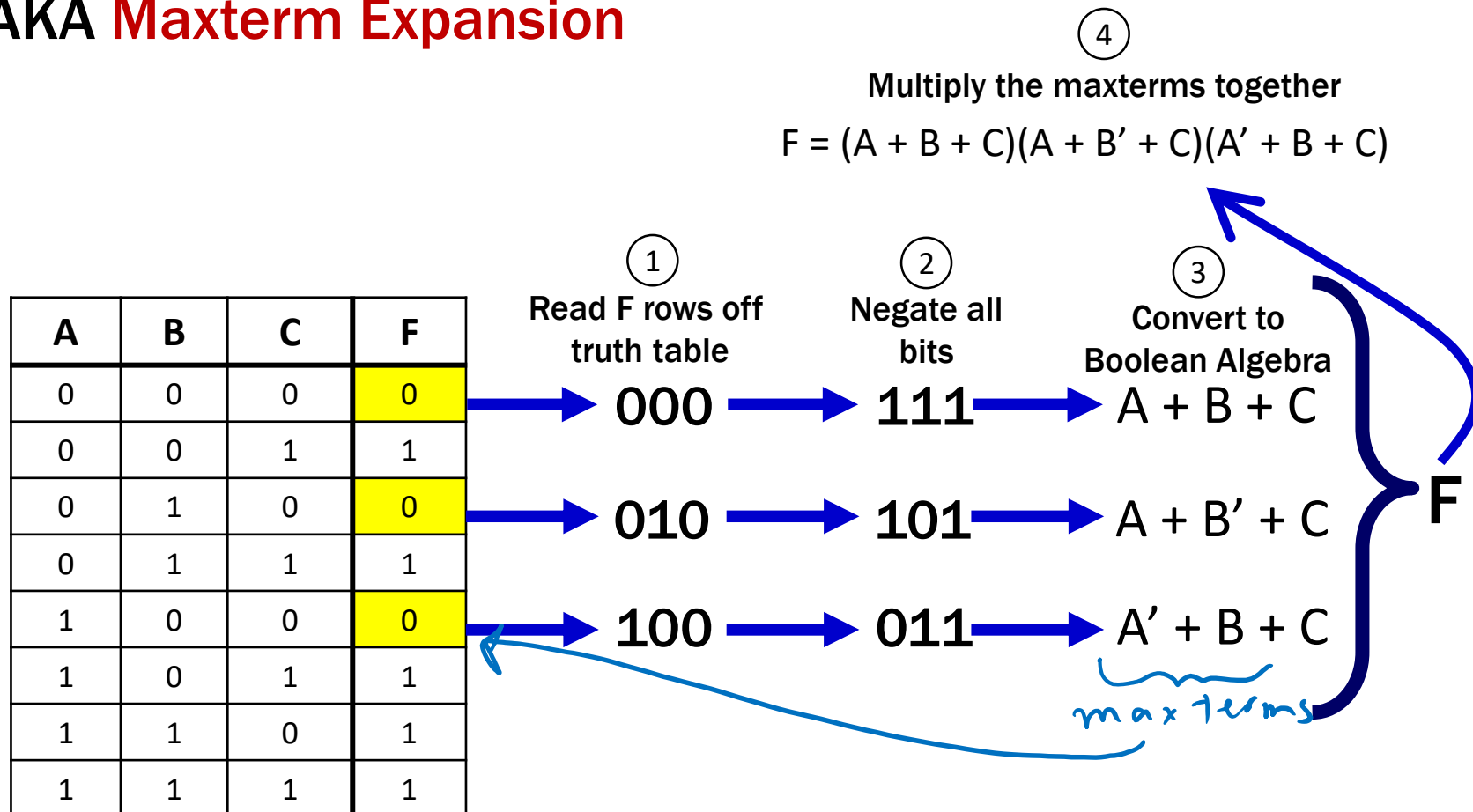
Product-of-Sums Canonical Form

- AKA **Conjunctive Normal Form (CNF)**
- AKA **Maxterm Expansion**



Product-of-Sums Canonical Form

- AKA **Conjunctive Normal Form (CNF)**
- AKA **Maxterm Expansion**



Product-of-Sums: Why does this procedure work?

Useful Facts:

- We know $(F')' = F$
- We know how to get a **minterm** expansion for F'

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1


$$F' = A'B'C' + A'BC' + AB'C'$$

Product-of-Sums: Why does this procedure work?

Useful Facts:

- We know $(F')' = F$
- We know how to get a **minterm** expansion for F'

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1


$$F' = A'B'C' + A'BC' + AB'C'$$

Taking the complement of both sides...

$$\underbrace{(F')}'_F = (A'B'C' + A'BC' + AB'C')'$$

And using DeMorgan/Comp....

$$F = (A'B'C')' (A'BC')' (AB'C')'$$

$$F = (A + B + C)(A + B' + C)(A' + B + C)$$

Product-of-Sums Canonical Form

Sum term (or maxterm)

- ORed sum of literals – input combination for which output is false
- each variable appears exactly once, true or inverted (but not both)

A	B	C	maxterms
0	0	0	$A+B+C$
0	0	1	$A+B+C'$
0	1	0	$A+B'+C$
0	1	1	$A+B'+C'$
1	0	0	$A'+B+C$
1	0	1	$A'+B+C'$
1	1	0	$A'+B'+C$
1	1	1	$A'+B'+C'$

F in canonical form:

$$F(A, B, C) = (A + B + C) (A + B' + C) (A' + B + C)$$

canonical form \neq minimal form

$$\begin{aligned} F(A, B, C) &= (A + B + C) (A + B' + C) (A' + B + C) \\ &= (A + B + C) (A + B' + C) \\ &\quad (A + B + C) (A' + B + C) \\ &= (A + C) (B + C) \end{aligned}$$

Predicate Logic

- **Propositional Logic**

“If you take the high road and I take the low road then I’ll arrive in Scotland before you.”

- **Predicate Logic**

“All positive integers x , y , and z satisfy $x^3 + y^3 \neq z^3$.”

Predicate Logic

- **Propositional Logic**

- Allows us to analyze complex propositions in terms of their simpler constituent parts (a.k.a. atomic propositions) joined by connectives

I play basketball in the morning.

- **Predicate Logic**

- Lets us analyze them at a deeper level by expressing how those propositions depend on the objects they are talking about

*I play basketball in all morning
of odd days of the month.*

Predicate Logic

Adds two key notions to propositional logic

- Predicates

- Quantifiers




Predicates

Predicate

- A function that returns a truth value, e.g.,

$\text{Cat}(x) ::= \text{"x is a cat"}$ 

$\text{Prime}(x) ::= \text{"x is prime"}$ 

$\text{HasTaken}(x, y) ::= \text{"student x has taken course y"}$ 

$\text{LessThan}(x, y) ::= \text{"x < y"}$

$\text{Sum}(x, y, z) ::= \text{"x + y = z"}$

$\text{GreaterThan5}(x) ::= \text{"x > 5"}$

$\text{HasNChars}(s, n) ::= \text{"string s has length n"}$ 

Predicates can have varying numbers of arguments and input types.

Domain of Discourse

For ease of use, we define one “type”/“domain” that we work over. This set of objects is called the “**domain of discourse**”.

For each of the following, what might the domain be?

(1) “x is a cat”, “x barks”, “x ruined my couch”

animals, pets, {cats and dogs}

(2) “x is prime”, “ $x = 0$ ”, “ $x < 0$ ”, “x is a power of two”

integers, numbers, reals

(3) “student x has taken course y” “x is a pre-req for z”

names, {student + and courses}

Domain of Discourse

For ease of use, we define one “type”/“domain” that we work over. This set of objects is called the “**domain of discourse**”.

For each of the following, what might the domain be?

(1) “x is a cat”, “x barks”, “x ruined my couch”

“mammals” or “sentient beings” or “cats and dogs” or ...

(2) “x is prime”, “ $x = 0$ ”, “ $x < 0$ ”, “x is a power of two”

“numbers” or “integers” or “integers greater than 5” or ...

(3) “student x has taken course y” “x is a pre-req for z”

“students and courses” or “university entities” or ...

Quantifiers

We use *quantifiers* to talk about collections of objects.

$\forall x P(x)$ *for all* $\{A_1, A_2, A_3, \dots\}$
 $P(A_1) \wedge P(A_2) \wedge P(A_3) \wedge \dots$
 $P(x)$ is true **for every** x in the domain
read as “**for all x , P of x** ”



$\exists x P(x)$ $\equiv P(A_1) \vee P(A_2) \vee P(A_3) \vee \dots$

There is an x in the domain for which $P(x)$ is true
read as “**there exists x , P of x** ”

Quantifiers

We use *quantifiers* to talk about collections of objects.

Universal Quantifier (“for all”): $\forall x P(x)$

$P(x)$ is true for **every** x in the domain

read as “**for all x , P of x** ”

Examples: Are these true?

- $\forall x \text{ Odd}(x)$

odd numbers.
 $\{1\}$, {All primes but 2}

- $\forall x \text{ LessThan5}(x)$

All negative, All numbers < 5
 $\{0\}$ $\{4.5\}$

Quantifiers

We use *quantifiers* to talk about collections of objects.

Universal Quantifier (“for all”): $\forall x P(x)$

$P(x)$ is true for **every** x in the domain

read as “**for all x , P of x** ”

Examples: Are these true? It depends on the domain. For example:

- $\forall x \text{ Odd}(x)$

- $\forall x \text{ LessThan4}(x)$

{1, 3, -1, -27}	Integers	Odd Integers
True	False	True
True	False	False

Quantifiers

We use *quantifiers* to talk about collections of objects.

Existential Quantifier (“exists”): $\exists x P(x)$

There is an x in the domain for which $P(x)$ is true
read as “**there exists x , P of x** ”

Examples: Are these true?

- $\exists x \text{ Odd}(x)$

All odd numbers, All numbers

- $\exists x \text{ LessThan5}(x)$

Quantifiers

We use *quantifiers* to talk about collections of objects.

Existential Quantifier (“exists”): $\exists x P(x)$

There is an x in the domain for which $P(x)$ is true
read as “**there exists x , P of x** ”

Examples: Are these true? It depends on the domain. For example:

- $\exists x \text{ Odd}(x)$

- $\exists x \text{ LessThan4}(x)$

$\{1, 3, -1, -27\}$	Integers	Positive Multiples of 5
True	True	True 5
True	True	False

Statements with Quantifiers

Just like with propositional logic, we need to define variables (this time **predicates**) before we do anything else. We must also now define a **domain of discourse** before doing anything else.

Domain of Discourse

Positive Integers

Predicate Definitions

Even(x) ::= "x is even"

Greater(x, y) ::= "x > y"

Odd(x) ::= "x is odd"

Equal(x, y) ::= "x = y"

Prime(x) ::= "x is prime"

Sum(x, y, z) ::= "x + y = z"

Statements with Quantifiers

Domain of Discourse
Positive Integers

Predicate Definitions

Even(x) ::= "x is even"	Greater(x, y) ::= "x > y"
Odd(x) ::= "x is odd"	Equal(x, y) ::= "x = y"
Prime(x) ::= "x is prime"	Sum(x, y, z) ::= "x + y = z"

Determine the truth values of each of these statements:

$\exists x \text{ Even}(x)$ T 2

$\forall x \text{ Odd}(x)$ F 2

$\forall x (\text{Even}(x) \vee \text{Odd}(x))$ T

$\exists x (\text{Even}(x) \wedge \text{Odd}(x))$ F 2

$\forall x \text{ Greater}(x+1, x)$ T

$\exists x (\text{Even}(x) \wedge \text{Prime}(x))$ T 2

Statements with Quantifiers

Domain of Discourse
Positive Integers

Predicate Definitions

Even(x) ::= "x is even"	Greater(x, y) ::= "x > y"
Odd(x) ::= "x is odd"	Equal(x, y) ::= "x = y"
Prime(x) ::= "x is prime"	Sum(x, y, z) ::= "x + y = z"

Determine the truth values of each of these statements:

- | | | |
|---|----------|--------------------------------------|
| $\exists x \text{ Even}(x)$ | T | e.g. 2, 4, 6, ... |
| $\forall x \text{ Odd}(x)$ | F | e.g. 2, 4, 6, ... |
| $\forall x (\text{Even}(x) \vee \text{Odd}(x))$ | T | every integer is either even or odd |
| $\exists x (\text{Even}(x) \wedge \text{Odd}(x))$ | F | no integer is both even and odd |
| $\forall x \text{ Greater}(x+1, x)$ | T | adding 1 makes a bigger number |
| $\exists x (\text{Even}(x) \wedge \text{Prime}(x))$ | T | Even(2) is true and Prime(2) is true |

Statements with Quantifiers

Domain of Discourse
Positive Integers

Predicate Definitions

Even(x) ::= "x is even"	Greater(x, y) ::= "x > y"
Odd(x) ::= "x is odd"	Equal(x, y) ::= "x = y"
Prime(x) ::= "x is prime"	Sum(x, y, z) ::= "x + y = z"

Translate the following statements to English

$\forall x \exists y \text{ Greater}(y, x)$

$\forall x \exists y \text{ Greater}(x, y)$

$\forall x \exists y (\text{Greater}(y, x) \wedge \text{Prime}(y))$

$\forall x (\text{Prime}(x) \rightarrow (\text{Equal}(x, 2) \vee \text{Odd}(x)))$

$\exists x \exists y (\text{Sum}(x, 2, y) \wedge \text{Prime}(x) \wedge \text{Prime}(y))$

Statements with Quantifiers (Literal Translations)

Domain of Discourse
Positive Integers

Predicate Definitions

Even(x) ::= "x is even"	Greater(x, y) ::= "x > y"
Odd(x) ::= "x is odd"	Equal(x, y) ::= "x = y"
Prime(x) ::= "x is prime"	Sum(x, y, z) ::= "x + y = z"

Translate the following statements to English

$\forall x \exists y \text{ Greater}(y, x)$

For every positive integer x, there is a positive integer y, such that $y > x$.

$\forall x \exists y \text{ Greater}(x, y)$

For every positive integer x, there is a positive integer y, such that $x > y$.

$\forall x \exists y (\text{Greater}(y, x) \wedge \text{Prime}(y))$

For every positive integer x, there is a pos. int. y such that $y > x$ and y is prime.

$\forall x (\text{Prime}(x) \rightarrow (\text{Equal}(x, 2) \vee \text{Odd}(x)))$

For each positive integer x, if x is prime, then $x = 2$ or x is odd.

$\exists x \exists y (\text{Sum}(x, 2, y) \wedge \text{Prime}(x) \wedge \text{Prime}(y))$

There exist positive integers x and y such that $x + 2 = y$ and x and y are prime.

Statements with Quantifiers (Natural Translations)

Domain of Discourse
Positive Integers

Predicate Definitions

Even(x) ::= "x is even"	Greater(x, y) ::= "x > y"
Odd(x) ::= "x is odd"	Equal(x, y) ::= "x = y"
Prime(x) ::= "x is prime"	Sum(x, y, z) ::= "x + y = z"

Translate the following statements to English

$\forall x \exists y \text{ Greater}(y, x)$

There is no greatest positive integer.

$\forall x \exists y \text{ Greater}(x, y)$

There is no least positive integer.

$\forall x \exists y (\text{Greater}(y, x) \wedge \text{Prime}(y))$

For every positive integer there is a larger number that is prime.

$\forall x (\text{Prime}(x) \rightarrow (\text{Equal}(x, 2) \vee \text{Odd}(x)))$

Every prime number is either 2 or odd.

$\exists x \exists y (\text{Sum}(x, 2, y) \wedge \text{Prime}(x) \wedge \text{Prime}(y))$

There exist prime numbers that differ by two."

English to Predicate Logic

Domain of Discourse

Mammals

Predicate Definitions

Cat(x) ::= "x is a cat"

Red(x) ::= "x is red"

LikesTofu(x) ::= "x likes tofu"

"Red cats like tofu"

"Some red cats don't like tofu"

English to Predicate Logic

Domain of Discourse

Mammals

Predicate Definitions

$\text{Cat}(x) ::= \text{"x is a cat"}$

$\text{Red}(x) ::= \text{"x is red"}$

$\text{LikesTofu}(x) ::= \text{"x likes tofu"}$

“Red cats like tofu”

$\forall x ((\text{Red}(x) \wedge \text{Cat}(x)) \rightarrow \text{LikesTofu}(x))$

“Some red cats don’t like tofu”

$\exists y ((\text{Red}(y) \wedge \text{Cat}(y)) \wedge \neg \text{LikesTofu}(y))$

English to Predicate Logic

Domain of Discourse
Mammals

Predicate Definitions

$\text{Cat}(x) ::= \text{"x is a cat"}$

$\text{Red}(x) ::= \text{"x is red"}$

$\text{LikesTofu}(x) ::= \text{"x likes tofu"}$

When putting two predicates together like this, we use an "and".

"Red cats like tofu"

When restricting to a smaller domain in a "for all" we use **implication**.

When there's no leading quantification, it means "for all".

"Some red cats don't like tofu"

When restricting to a smaller domain in an "exists" we use **and**.

"Some" means "there exists".

Negations of Quantifiers

Predicate Definitions

PurpleFruit(x) ::= “x is a purple fruit”

(*) $\forall x$ PurpleFruit(x) (“All fruits are purple”)

What is the negation of (*)?

- (a) “there exists a purple fruit”
- (b) “there exists a non-purple fruit”
- (c) “all fruits are not purple”

Try your intuition! Which one “feels” right?

Key Idea: In **every** domain, **exactly one** of a statement and its negation should be true.

Negations of Quantifiers

Predicate Definitions

PurpleFruit(x) ::= “x is a purple fruit”

(*) $\forall x$ PurpleFruit(x) (“All fruits are purple”)

What is the negation of (*)?

- (a) “there exists a purple fruit”
- (b) “there exists a non-purple fruit”
- (c) “all fruits are not purple”

Key Idea: In **every** domain, **exactly one** of a statement and its negation should be true.

Domain of Discourse

{plum}

Domain of Discourse

{apple}

Domain of Discourse

{plum, apple}

The only choice that ensures exactly one of the statement and its negation is (b).

De Morgan's Laws for Quantifiers

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

De Morgan's Laws for Quantifiers

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

“There is no largest integer”

$$\neg \exists x \forall y (x \geq y)$$

$$\equiv \forall x \neg \forall y (x \geq y)$$

$$\equiv \forall x \exists y \neg (x \geq y)$$

$$\equiv \forall x \exists y (x < y)$$

“For every integer there is a larger integer”

Scope of Quantifiers

$$\exists x (P(x) \wedge Q(x)) \quad \text{vs.} \quad \exists x P(x) \wedge \exists x Q(x)$$

Scope of Quantifiers

$$\exists x (P(x) \wedge Q(x)) \quad \text{vs.} \quad \exists x P(x) \wedge \exists x Q(x)$$

This one asserts P
and Q of the *same* x.

This one asserts P and Q
of potentially different x's.

Scope of Quantifiers

Example: $\text{NotLargest}(x) \equiv \exists y \text{ Greater}(y, x)$
 $\equiv \exists z \text{ Greater}(z, x)$

truth value:

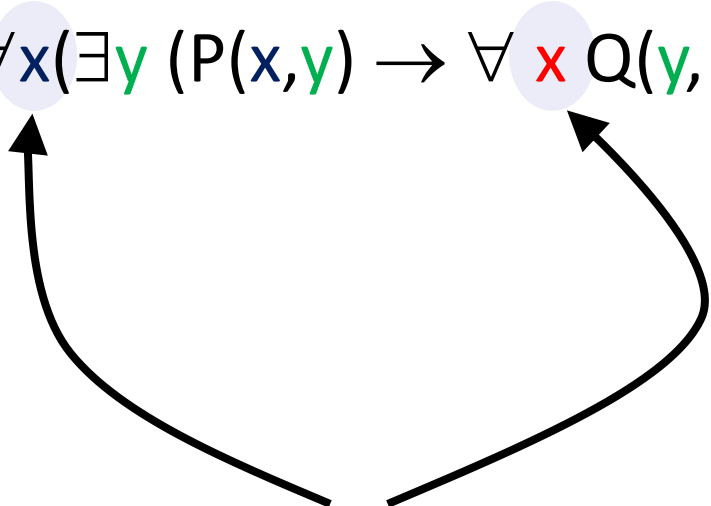
doesn't depend on y or z “**bound** variables”

does depend on x “**free** variable”

quantifiers only act on free variables of the formula
they quantify

$$\forall x (\exists y (P(x, y) \rightarrow \forall x Q(y, x)))$$

Quantifier “Style”

$$\forall x (\exists y (P(x, y) \rightarrow \forall x Q(y, x)))$$
The diagram shows the logical expression $\forall x (\exists y (P(x, y) \rightarrow \forall x Q(y, x)))$. The variable x in the first universal quantifier is highlighted with a light blue circle. The variable x in the second universal quantifier is also highlighted with a light blue circle. Two curved arrows originate from a point below the text and point to these two x variables, illustrating that the same variable name is used for two different scopes, which is a poor style choice.

This isn't “wrong”, it's just horrible style.
Don't confuse your reader by using the same
variable multiple times...there are a lot of letters...