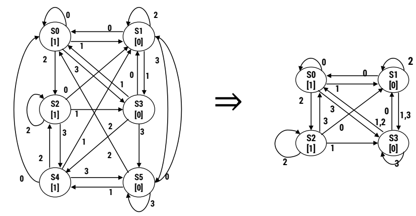
cse 311: foundations of computing

Spring 2015

Lecture 24: DFAs, NFAs, and regular expressions

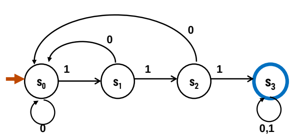
highlights

- FSMs with output at states
- State minimization



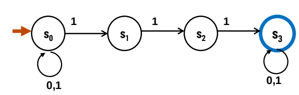
highlights

Lemma: The language recognized by a DFA is the set of strings x that label some path from its start state to one of its final states



nondeterministic finite automaton (NFA)

- Graph with start state, final states, edges labeled by symbols (like DFA) but
  - Not required to have exactly 1 edge out of each state labeled by each symbol--- can have 0 or >1
  - Also can have edges labeled by empty string ε
- **Definition:** x is in the language recognized by an NFA if and only if x labels a path from the start state to some final state



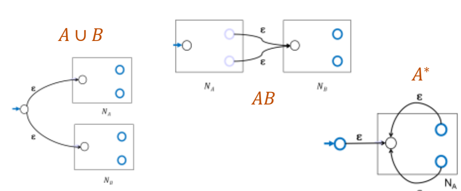
building an NFA

binary strings that have
- an even # of 1's
- or contain the substring 111 or 1000

NFAs and regular expressions

Theorem: For any set of strings (language) A described by a regular expression, there is an NFA that recognizes A.

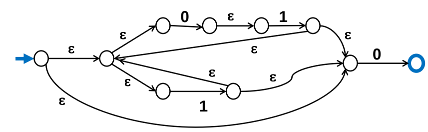
Proof idea: Structural induction based on the recursive definition of regular expressions...



build an NFA for (01 ∪1)*0

solution

(01 ∪1)*0



NFAs vs. DFAs

Every DFA **is** an NFA
- DFAs have requirements that NFAs don't have

Can NFAs recognize more languages?

NFAs vs. DFAs

Every DFA **is** an NFA
- DFAs have requirements that NFAs don't have

Can NFAs recognize more languages? No!

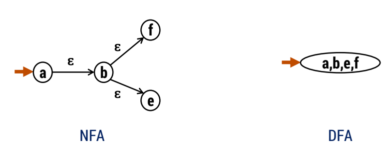
Theorem: For every NFA there is a DFA that recognizes exactly the same language.

conversion of NFAs to DFAs

Proof Idea:
- The DFA keeps track of ALL the states that the part of the input string read so far can reach in the NFA
- There will be one state in the DFA for each *subset* of states of the NFA that can be reached by some string

conversion of NFAs to a DFAs

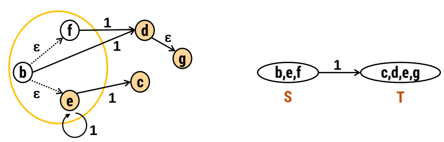
New start state for DFA
- The set of all states reachable from the start state of the NFA using only edges labeled ε



conversion of NFAs to a DFAs

For each state of the DFA corresponding to a set S of states of the NFA and each symbol **s**

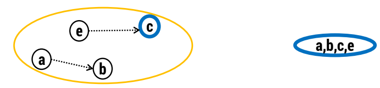
– Add an edge labeled **s** to state corresponding to T, the set of states of the NFA reached by
  - starting from some state in S, then
  - following one edge labeled by **s**, and
  - then following some number of edges labeled by ε
– T will be ∅ if no edges from S labeled **s** exist



conversion of NFAs to a DFAs

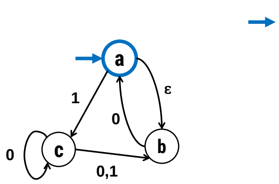
Final states for the DFA

– All states whose set contain some final state of the NFA



NFA DFA

example: NFA to DFA



NFA DFA

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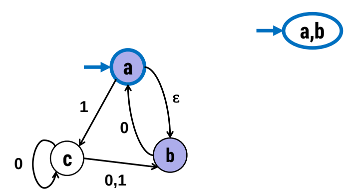
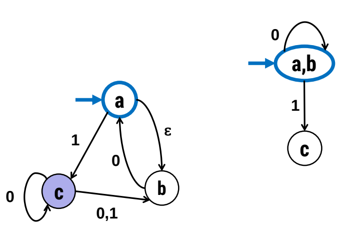


NFA DFA

example: NFA to DFA



NFA DFA

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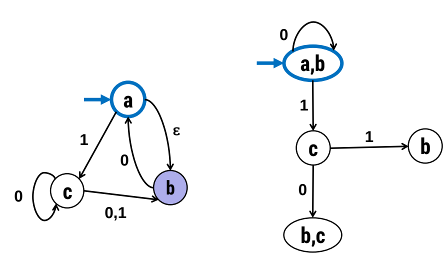
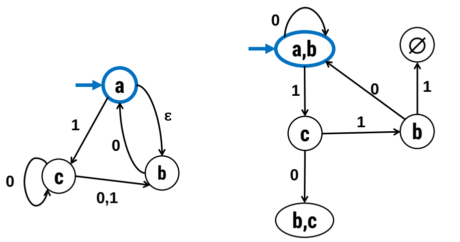


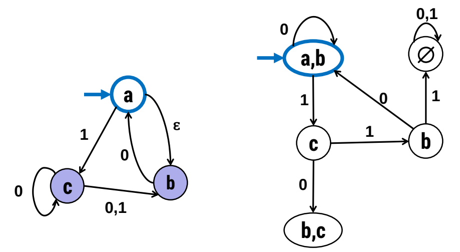
NFA DFA

example: NFA to DFA



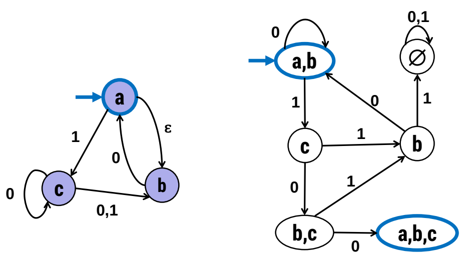
NFA DFA

example: NFA to DFA



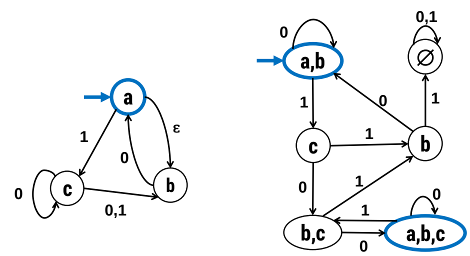
NFA DFA

example: NFA to DFA



NFA DFA

example: NFA to DFA

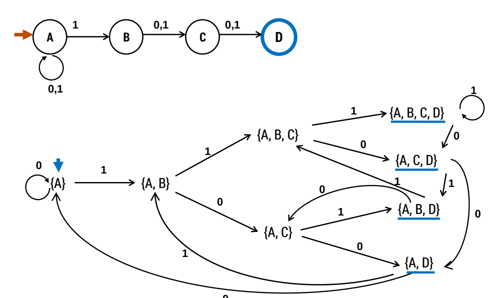


NFA DFA

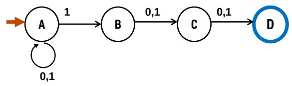
exponential blow-up in simulating mondeterminism

- In general the DFA might need a state for every subset of states of the NFA
  - Power set of the set of states of the NFA
  - n-state NFA yields DFA with at most 2^n states
  - We saw an example where roughly 2^n is necessary
    Is the n^{th} char from the end a 1?
- The famous “P=NP?” question asks whether a similar blow-up is always necessary to get rid of nondeterminism for polynomial-time algorithms

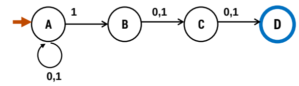
1 in third position from end

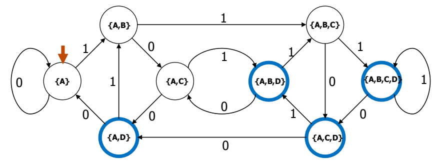


1 in third position from end



1 in third position from end





DFAs ≡ regular expressions

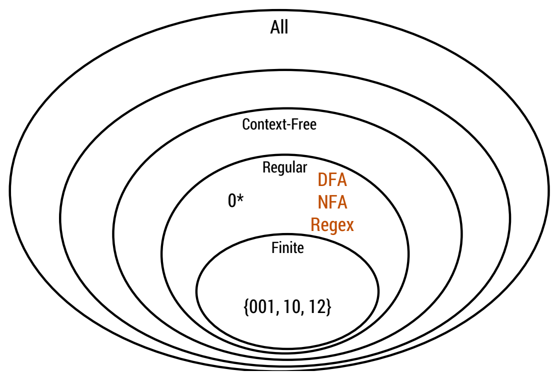
We have shown how to build an optimal DFA for every regular expression

- Build NFA
- Convert NFA to DFA using subset construction
- Minimize resulting DFA

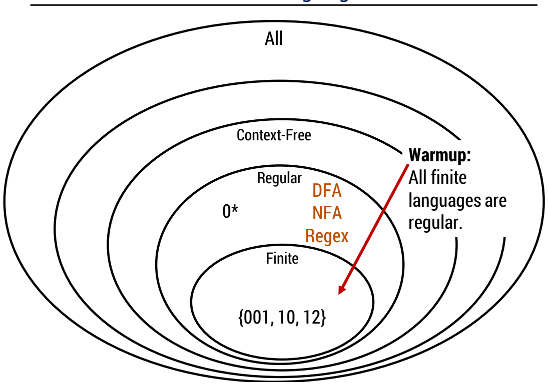
Theorem: A language is recognized by a DFA if and only if it has a regular expression.

We show the other direction of the proof at the end of these lecture slides.

languages and machines!



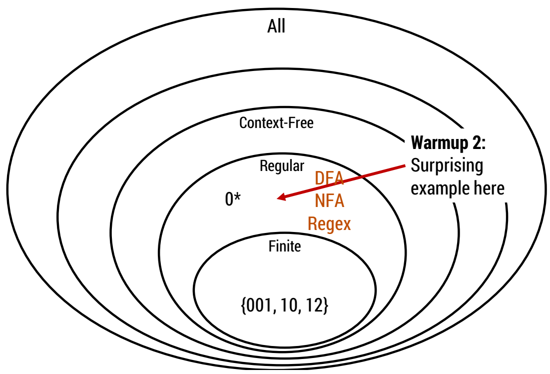
languages and machines!



DFAs recognize any finite language

Exercise: Hard code it into the DFA.

languages and machines!



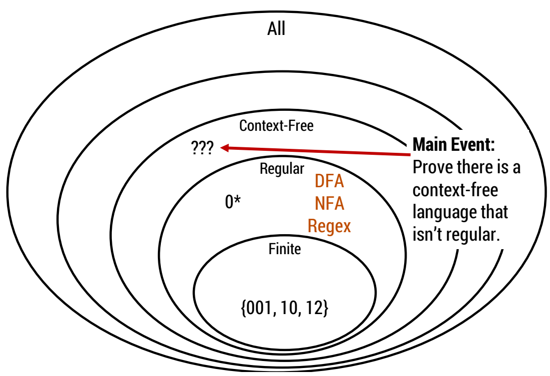
an interesting regular language

$L = \{x \in \{0,1\}^* : x \text{ has an equal number of substrings 01 and 10}\}.$

L is infinite.

L is regular.

languages and machines!



DFAs ≡ regular expressions

Theorem: A language is recognized by a DFA if and only if it has a regular expression

Proof: Last class: RegExp → NFA → DFA

Now: NFA → RegExp

(Enough to show this since every DFA is also an NFA.)

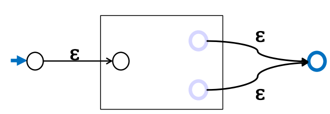
generalized NFAs

- Like NFAs but allow
  - Parallel edges
  - Regular Expressions as edge labels

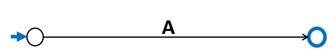
    NFAs already have edges labeled ε or ***a***
- An edge labeled by **A** can be followed by reading a string of input chars that is in the language represented by **A**
- A string x is accepted iff there is a path from start to final state labeled by a regular expression whose language contains x

starting from an NFA

Add new start state and final state



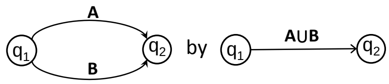
Then eliminate original states one by one, keeping the same language, until it looks like:



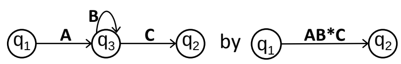
Final regular expression will be **A**

only two simplification rules

- Rule 1: For any two states q_1 and q_2 with parallel edges (possibly $q_1=q_2$), replace



- Rule 2: Eliminate non-start/final state q_3 by replacing all

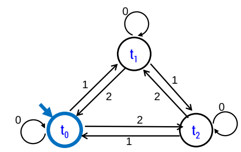


for *every* pair of states q_1, q_2 (even if $q_1=q_2$)

converting an NFA to a regular expression

Consider the DFA for the mod 3 sum

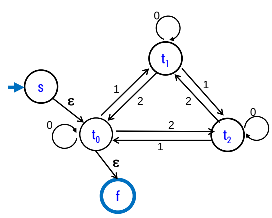
- Accept strings from {0,1,2}* where the digits mod 3 sum of the digits is 0



splicing out a node

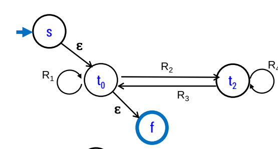
Label edges with regular expressions

$t_0 \to t_1 \to t_0$: 10*2
$t_0 \to t_1 \to t_2$: 10*1
$t_2 \to t_1 \to t_0$: 20*2
$t_2 \to t_1 \to t_2$: 20*1

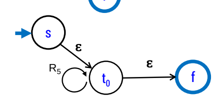


finite automaton without t_1

R_1: $0 \cup 10^*2$
R_2: $2 \cup 10^*1$
R_3: $1 \cup 20^*2$
R_4: $0 \cup 20^*1$



R_5: $R_1 \cup R_2R_4^*R_3$



Final regular expression:
$(0 \cup 10^*2 \cup (2 \cup 10^*1)(0 \cup 20^*1)^*(1 \cup 20^*2))^*$