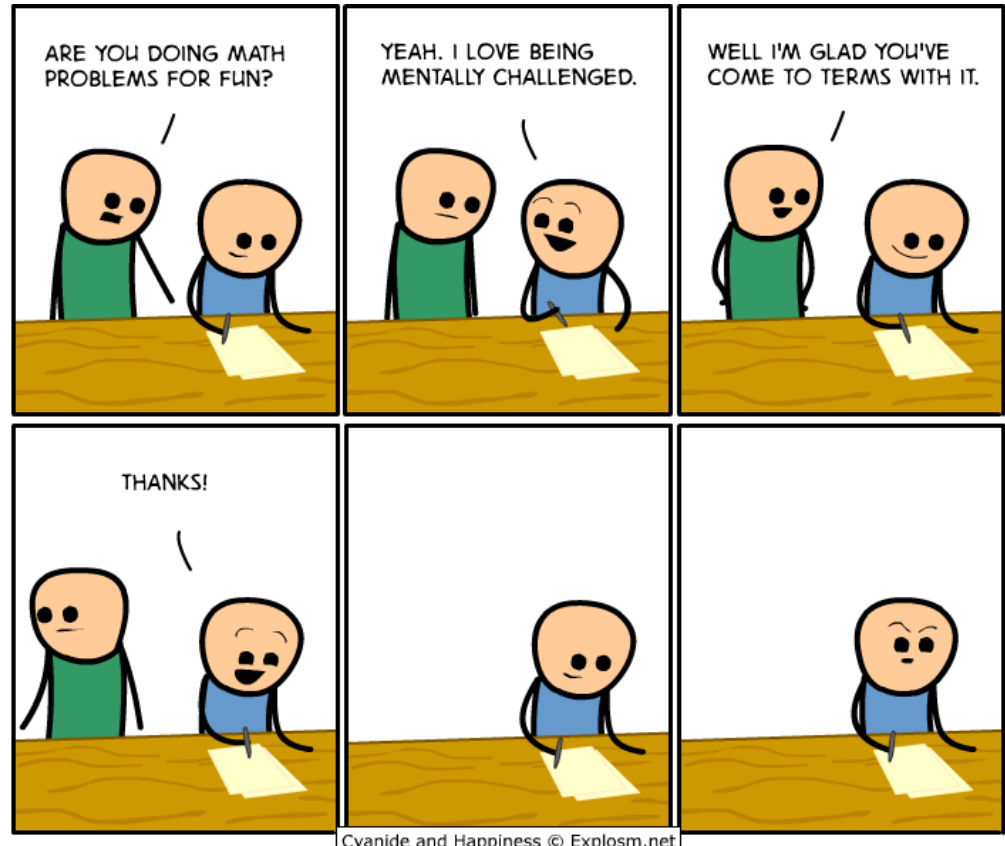
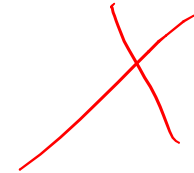


cse 311: foundations of computing

Spring 2015

Lecture 13: Primes, GCDs, modular inverses



review: repeated squaring

Since $a \bmod m \equiv a \pmod{m}$ for any a

we have $a^2 \bmod m = (a \bmod m)^2 \bmod m$

and $a^4 \bmod m = (a^2 \bmod m)^2 \bmod m$

and $a^8 \bmod m = (a^4 \bmod m)^2 \bmod m$

and $a^{16} \bmod m = (a^8 \bmod m)^2 \bmod m$

and $a^{32} \bmod m = (a^{16} \bmod m)^2 \bmod m$

Can compute $a^k \bmod m$ for $k = 2^i$ in only i steps

review: general algorithm

ModPow(a, k, m) should compute $a^k \bmod m$.

If $k == 0$ then

return 1

If $(k \bmod 2 == 0)$ then

return ModPow($a^2 \bmod m$, $k/2$, m) ←

else

return ($a \times$ ModPow(a , $k - 1$, m)) mod m ←

$$\begin{aligned} k &= 81453 \\ &= (10011111000101101)_2 \\ &= 2^{16} + 2^{13} + 2^{12} + 2^{11} + 2^{10} + 2^9 + 2^5 + 2^3 + 2^2 + 2^0 \end{aligned}$$

Total # of arithmetic operations $\sim 4 \times 16 = 64$

An integer p greater than 1 is called *prime* if the only positive factors of p are 1 and p .

$$p = 13$$

A positive integer that is greater than 1 and is not prime is called *composite*.

$$26 = 13 \times 2$$

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A positive integer that is greater than 1 and is not prime is called *composite*.

FUNDAMENTAL THEOREM OF ARITHMETIC

Every positive integer greater than 1 has a unique prime factorization

$$48 = 2 \cdot 2 \cdot 2 \cdot 2 \cdot 3$$

$$591 = 3 \cdot 197$$

$$45,523 = 45,523$$

$$321,950 = 2 \cdot 5 \cdot 5 \cdot 47 \cdot 137$$

$$1,234,567,890 = 2 \cdot 3 \cdot 3 \cdot 5 \cdot 3,607 \cdot 3,803$$

$$1 = \underbrace{\hspace{10em}}_{\text{empty product}}$$

FACTORIZATION

If n is composite, it has a factor of size at most \sqrt{n} .

$$n = p_1 \cdot p_2 \cdots p_k, \quad k \geq 2$$

$$\text{If } p_1 > \sqrt{n}, \quad p_2 > \sqrt{n}$$

$$\implies p_1 \cdot p_2 > n$$
$$n \geq p_1 \cdot p_2 > n$$

contradiction.

EUCCLAD'S THEOREM

There are an infinite number of primes.

7. F
8. $F \rightarrow P$
9. P

Proof by contradiction:

Suppose that there are only a finite number of primes:

p_1, p_2, \dots, p_n

$$p_1 p_2 \dots p_n + 1 = p_{i_1} p_{i_2} \dots p_{i_k}$$

$$1 \equiv 0 \pmod{p_{i_1}}$$

contradiction.

p_1, \dots, p_n
first n primes

$$p_1 \dots p_n + 1 = q_1 \dots q_m \quad \square$$

FAMOUS ALGORITHMIC PROBLEMS

- **Primality Testing**

- Given an integer n , determine if n is prime

- Fermat's little theorem test:

If p is prime and $a \neq 0$, then $a^{p-1} \equiv 1 \pmod{p}$

- **Factoring**

- Given an integer n , determine the prime factorization of n

Factor the following 232 digit number [RSA768]:

123018668453011775513049495838496272077285
356959533479219732245215172640050726365751
874520219978646938995647494277406384592519
255732630345373154826850791702612214291346
167042921431160222124047927473779408066535
1419597459856902143413



123018668453011775513049495838496272077285356959533479219
732245215172640050726365751874520219978646938995647494277
406384592519255732630345373154826850791702612214291346167
042921431160222124047927473779408066535141959745985690214
3413

=

334780716989568987860441698482126908177047949837
13768568912431388982883793878 17
43087737814467999489

×

3674604366679959042824463379 643
4308764267603228381573966651 068
10270092798736308917



GREATEST COMMON DIVISOR

GCD(a, b):

Largest integer d such that $d \mid a$ and $d \mid b$

– $\text{GCD}(100, 125) = 25$

– $\text{GCD}(17, 49) = 1$

– $\text{GCD}(11, 66) = 11$

– $\text{GCD}(13, 0) = 13$

– $\text{GCD}(180, 252) = 36$

$2^2 \cdot 3^2 \cdot 5$ $2^2 \cdot 3^2 \cdot 7$

$2^2 \cdot 3^2 = 36$

GCD AND FACTORING

$$a = 2^3 \cdot 3 \cdot 5^2 \cdot 7 \cdot 11 = 46,200$$

$$b = 2 \cdot 3^2 \cdot 5^3 \cdot 7 \cdot 13 = 204,750$$

$$\text{GCD}(a, b) = 2^{\min(3,1)} \cdot 3^{\min(1,2)} \cdot 5^{\min(2,3)} \cdot 7^{\min(1,1)} \cdot 11^{\min(1,0)} \cdot 13^{\min(0,1)}$$



Factoring is expensive!

Can we compute **GCD(a,b)** without factoring?

If a and b are positive integers, then
 $\gcd(a, b) = \gcd(b, a \bmod b)$

Proof:

$a \equiv c \pmod{d}$

$\quad \quad \quad b \quad \quad \quad c$

By definition $a = \cancel{(a \operatorname{div} b) \cdot b} + \underbrace{(a \bmod b)}_c \pmod{d}$

If $d \mid a$ and $d \mid b$ then $d \mid (a \bmod b)$.

If $d \mid b$ and $d \mid (a \bmod b)$ then $\underline{d \mid a}$.

EUCCLAD'S ALGORITHM

Repeatedly use the GCD fact to reduce numbers
until you get $\text{GCD}(x, 0) = x$.

$$\begin{aligned}\text{GCD}(660, 126) &= \text{GCD}(126, 30) \\ &= \text{GCD}(30, 6) \\ &= \text{GCD}(6, 0) \\ &= 6.\end{aligned}$$

$$\begin{aligned}a &> b \\ \text{GCD}(a, b) \\ &= \\ \text{GCD}(b, a \bmod b)\end{aligned}$$

EUCLEID'S ALGORITHM

$$\text{GCD}(x, y) = \text{GCD}(y, x \bmod y)$$

```
int GCD(int a, int b){ /* a >= b, b > 0 */
    int tmp;
    while (b > 0) {
        tmp = a % b;
        a = b;
        b = tmp;
    }
    return a;
}
```

Example: GCD(660, 126)

BEZOUT'S THEOREM

If a and b are positive integers, then there exist integers s and t such that

$$\gcd(a,b) = sa + tb$$

EXTENDED EUCLIDEAN ALGORITHM

$$\underline{13(27x) \equiv 4 \pmod{35}}$$
$$x \equiv 17$$

$$1 \equiv 13 \cdot 27 \pmod{35}$$

- Can use Euclid's Algorithm to find s, t such that

$$\gcd(a, b) = sa + tb$$

$$1 = \gcd(35, 27) = 13 \cdot 27 + (-10) \cdot 35$$

- e.g. $\gcd(35, 27)$: $35 = 1 \cdot 27 + 8$ $35 - 1 \cdot 27 = 8$

$$27 = 3 \cdot 8 + 3$$
$$27 - 3 \cdot 8 = 3$$

$$27x \equiv 4 \pmod{35}$$

$$8 = 2 \cdot 3 + 2$$
$$8 - 2 \cdot 3 = 2$$

\Rightarrow

$$3 = 1 \cdot 2 + 1$$
$$3 - 1 \cdot 2 = 1$$

$$x \equiv 17 \pmod{35}$$

$$2 = 2 \cdot 1 + 0$$

- Substitute back from the bottom

$$1 = 3 - 1 \cdot 2 = 3 - 1(8 - 2 \cdot 3) = (-1) \cdot 8 + 3 \cdot 3$$

$$= (-1) \cdot 8 + 3(27 - 3 \cdot 8) = 3 \cdot 27 + (-10) \cdot 8$$

=

MULTIPLICATIVE INVERSE mod m

Suppose $\text{GCD}(a, m) = 1$

By Bézout's Theorem, there exist integers s and t such that $sa + tm = 1$.

$s \bmod m$ is the multiplicative inverse of a :

$$1 = (sa + tm) \bmod m = sa \bmod m$$

SOLVING MODULAR EQUATIONS

Solving $ax \equiv b \pmod{m}$ for unknown x when $\gcd(a, m) = 1$.

1. Find s such that $sa + tm = 1$
2. Compute $a^{-1} = s \pmod{m}$, the multiplicative inverse of a modulo m
3. Set $x = (a^{-1} \cdot b) \pmod{m}$

EXAMPLE

Solve: $7x \equiv 1 \pmod{26}$