



So far:

- Propositional logic
- Logic to build circuits
- Predicates and quantifiers
- Proof systems and logical inference
- Basic set theory

[optional] Proof review session on Wed @ 6pm in EE 105

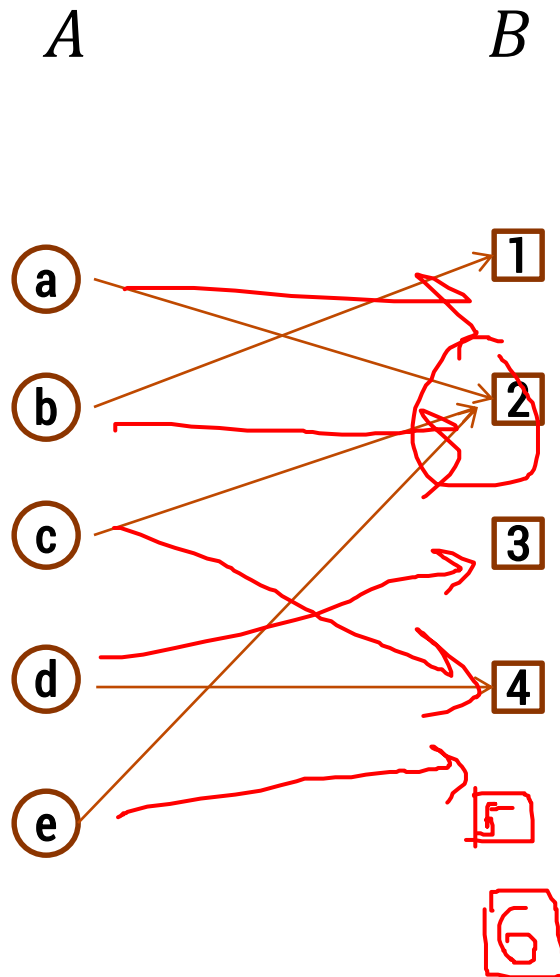
Question: If the domain of discourse is empty and  $P$  is a predicate, what is the truth value of:

$$\exists x P(x)$$

$$\forall x P(x)$$

A function from  $A$  to  $B$ :

- Every element of  $A$  is assigned to exactly one element of  $B$ .
- We write  $f : A \rightarrow B$ .
- "Image of  $X$  under  $f$ " = " $f(X)$ "  
$$= \{x : \exists y (y \in X \wedge x = f(y))\}$$
- Domain of  $f$  is  $A$
- Codomain of  $f$  is  $B$
- Image of  $f$  = Image of domain under  $f$   
= all the elements pointed to by something in the domain.



Image( $\{a\}$ ) =

Image( $\{a, e\}$ ) =

Image( $\{a, b\}$ ) =

Image( $A$ ) =

$\{1, 2\}$   
 $\{1, 2, 3\}$   
 $\{1, 2, 4\}$

# injections and surjections

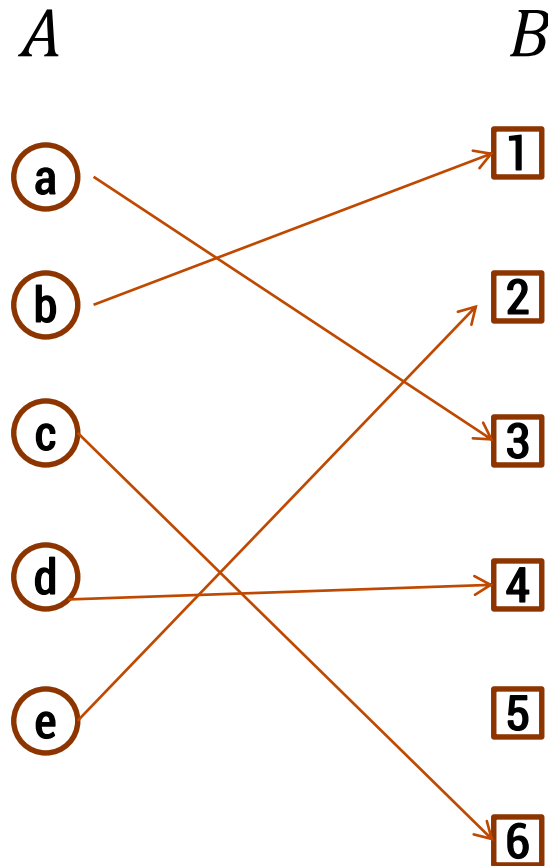
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A function  $f : A \rightarrow B$  is **one-to-one** (or, **injective**) if every output corresponds to at most one input, i.e.  $f(x) = f(x') \Rightarrow x = x'$  for all  $x, x' \in A$ .

A function  $f : A \rightarrow B$  is **onto** (or, **surjective**) if every output gets hit, i.e. for every  $y \in B$ , there exists  $x \in A$  such that  $f(x) = y$ .

# is this function one-to-one? is it onto?

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It is one-to-one, because nothing in  $B$  is pointed to by multiple elements of  $A$ .

It is not onto, because  $5$  is not pointed to by anything.

~~$x \mapsto x^2$~~

$x \mapsto x^2$

One-to-one (?)

NO  $\nexists$

Onto (?)

NO only  $x \geq 0$

$x \mapsto x^3 - x$

NO  $-1, 0, 1$



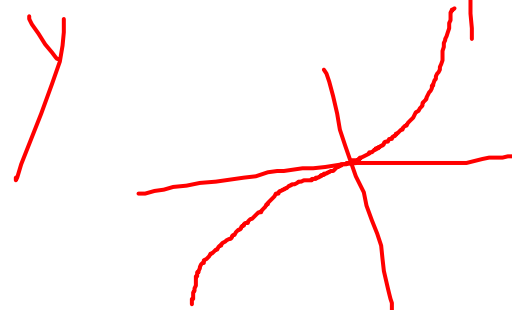
$x \mapsto e^x$

Y



$x \mapsto x^3$

Y



Domain: Reals



Dear HBO, this is a slide about digital watermarking.



# “number theory” (and applications to computing)

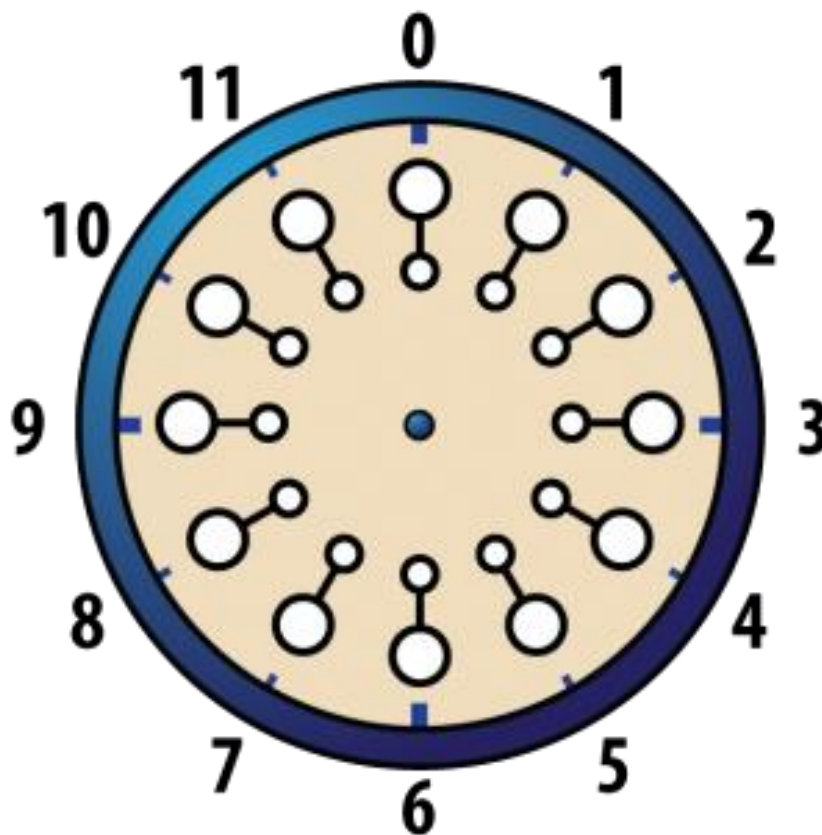
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- **How whole numbers work**  
[fascinating, deep, weird area of mathematics that no one understands, but the basics are easy and really useful]
- **Many significant applications**
  - Cryptography [this is how SSL works]
  - Hashing
  - Security
  - Error-correcting codes [this is how your bluray player works]
- **Important tool set**

```
public class Test {  
    final static int SEC_IN_YEAR = 364*24*60*60;  
    public static void main(String args[]) {  
        System.out.println(  
            "I will be alive for at least " +  
            SEC_IN_YEAR * 101 + " seconds."  
        );  
    }  
}
```

```
----jGRASP exec: java Test  
I will be alive for at least -186619904 seconds.  
----jGRASP: operation complete.
```

**Arithmetic over a finite domain: Math with wrap around**



Integers  $a$ ,  $b$ , with  $a \neq 0$ . We say that  $a$  **divides**  $b$  iff there is an integer  $k$  such that  $b = k a$ . The notation  $a \mid b$  denotes “ $a$  divides  $b$ .”

$$12 = 6 * 2.$$

$$13 = \_\_ * 2.$$

# division theorem

Let  $a$  be an integer and  $d$  a positive integer. Then there are *unique* integers  $q$  and  $r$ , with  $0 \leq r < d$ , such that  $a = d q + r$ .

$$q = a \text{ div } d$$

$$r = a \text{ mod } d$$

$$a = 17$$

$$d = 3$$

$$q = 5$$

$$r = 2$$

$$\begin{array}{ll} a = -13 & q = -5 \\ d = 3 & r = 2 \end{array}$$

Note:  $r \geq 0$  even if  $a < 0$ .  
Not quite the same as  $a \% d$ .

# arithmetic mod 7

$$9 \times 8 \equiv 2 \pmod{7}$$

$$a +_7 b = (a + b) \bmod 7$$

$$2 \times 1 \equiv 2 \pmod{7}$$

$$a \times_7 b = (a \times b) \bmod 7$$

$$5 * 5 = 25 = 7 \times 3 + 4$$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

# modular congruence

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Let  $a$  and  $b$  be integers, and  $m$  be a positive integer. We say  $a$  is **congruent** to  $b$  **modulo**  $m$  if  $m$  divides  $a - b$ . We use the notation  $a \equiv b \pmod{m}$  to indicate that  $a$  is congruent to  $b$  modulo  $m$ .

$$12 \equiv -2 \pmod{7}$$

$$12 - (-2) = 14$$



# modular arithmetic: examples

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$$A \equiv 0 \pmod{2}$$

This statement is the same as saying “A is even”; so, any A that is even (including negative even numbers) will work.

$$1 \equiv 0 \pmod{4}$$

This statement is false. If we take it mod 1 instead, then the statement is true.

16

$$\begin{aligned} A &= 17k - 1 \\ &= 17j + 16 \end{aligned}$$

$$A \equiv -1 \pmod{17}$$

If  $A = 17x - 1 = 17x + 16$  for an integer x, then it works.

Note that  $(m - 1) \bmod m$

$$\begin{aligned} &= ((m \bmod m) + (-1 \bmod m)) \bmod m \\ &= (0 + -1) \bmod m \\ &= -1 \bmod m \end{aligned}$$

# modular arithmetic can haz sense

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**Theorem:** Let  $a$  and  $b$  be integers, and let  $m$  be a positive integer. Then  $a \equiv b \pmod{m}$  if and only if  $a \bmod m = b \bmod m$ .

**Proof:** Suppose that  $a \equiv b \pmod{m}$ .

By definition:  $a \equiv b \pmod{m}$  implies  $m \mid (a - b)$

which by definition implies that  $a - b = km$  for some integer  $k$ .

Therefore  $a = b + km$ .

Taking both sides modulo  $m$  we get

$$a \bmod m = (b + km) \bmod m = b \bmod m$$

# modular arithmetic can haz sense

---

**Theorem:** Let  $a$  and  $b$  be integers, and let  $m$  be a positive integer. Then  $a \equiv b \pmod{m}$  if and only if  $a \bmod m = b \bmod m$ .

**Proof:** Suppose that  $a \bmod m = b \bmod m$ .

By the division theorem,  $a = mq + (a \bmod m)$  and

$b = ms + (b \bmod m)$  for some integers  $q, s$ .

$$\begin{aligned} a - b &= (mq + (a \bmod m)) - (ms + (b \bmod m)) \\ &= m(q - s) + (a \bmod m - b \bmod m) \\ &= m(q - s) \text{ since } a \bmod m = b \bmod m \end{aligned}$$

Therefore  $m \mid (a-b)$  and so  $a \equiv b \pmod{m}$

# consistency of addition

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Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  **$a + c \equiv b + d \pmod{m}$**

Suppose  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ .

Unrolling definitions gives us some  $k$  such that  $a - b = km$ , and some  $j$  such that  $c - d = jm$ .

Adding the equations together gives us

**$(a + c) - (b + d) = m(k + j)$** . Now, re-applying the definition of mod gives us  **$a + c \equiv b + d \pmod{m}$** .

# consistency of multiplication

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Let  $m$  be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  **$ac \equiv bd \pmod{m}$**

Suppose  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ .

Unrolling definitions gives us some  $k$  such that  $a - b = km$ , and some  $j$  such that  $c - d = jm$ .

Then,  $a = km + b$  and  $c = jm + d$ .

Multiplying both together gives us

$$ac = (km + b)(jm + d) = kjm^2 + kmd + jmb + bd$$

Rearranging gives us  $ac - bd = m(kjm + kd + jb)$ .

Using the definition of mod gives us  **$ac \equiv bd \pmod{m}$** .

# example

Let  $n$  be an integer.

Prove that  $n^2 \equiv 0 \pmod{4}$  or  $n^2 \equiv 1 \pmod{4}$

$$0^2 \equiv 0$$

$$1^2 \equiv 1$$

$$2^2 \equiv 0$$

$$3^2 \equiv 1$$

$$420 \quad n \neq 1$$

$$n = 4q + r$$

$$r = 0, 1, 2, 3$$

$$16q^2 + r^2 + 8qr$$

□