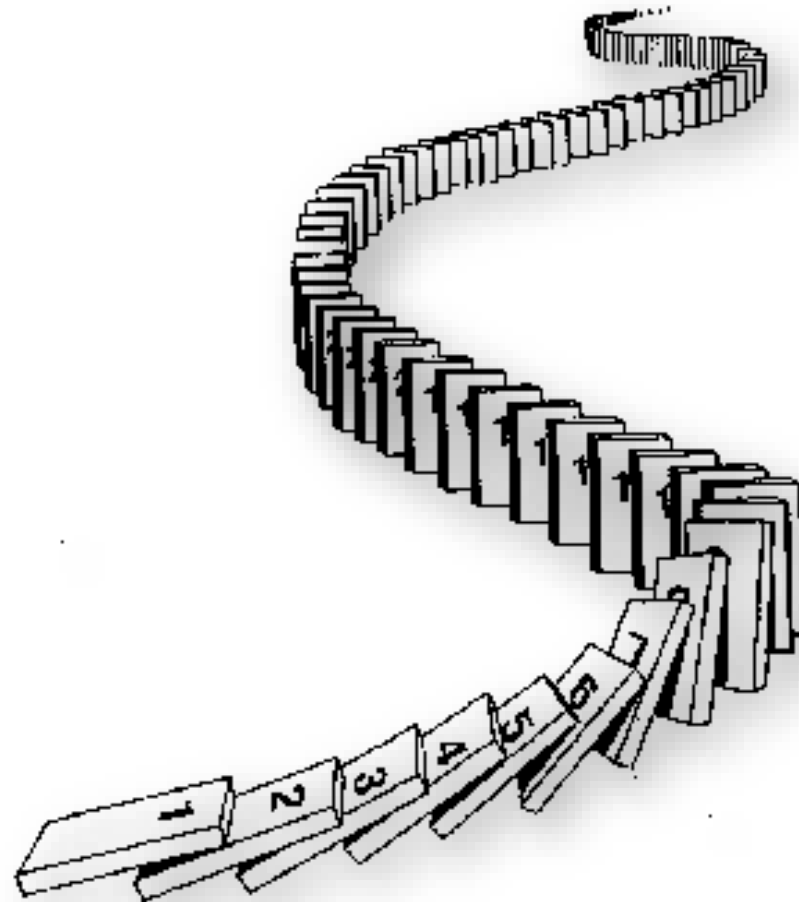


# cse 311: foundations of computing

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Fall 2015

## Lecture 15: Induction



prove: for all  $n \geq 0$ ,  $a$  is odd  $\rightarrow a^n$  is odd

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Let  $n > 0$  be arbitrary.

Suppose that  $a$  is odd. We know that if  $a, b$  are odd, then  $ab$  is also odd.

So:  $(\dots ((a \cdot a) \cdot a) \dots \cdot a) = a^n$  [ $n$  times]

Those “...”s are a problem! We’re trying to say “we can use the same argument over and over...”

We’ll come back to this.

# mathematical induction

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## Method for proving statements about all integers $\geq 0$

- A new logical inference rule!
  - It only applies over the natural numbers
  - The idea is to **use** the special structure of the naturals to prove things more easily
- Particularly useful for reasoning about programs!

```
for(int i=0; i < n; n++) { ... }
```

- Show  $P(i)$  holds after  $i$  times through the loop

```
public int f(int x) {  
    if (x == 0) { return 0; }  
    else { return f(x-1)+1; }  
}
```

- $f(x) = x$  for all values of  $x \geq 0$  naturally shown by induction.

# induction is a rule of inference

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Domain: Natural Numbers

$$\begin{array}{l} P(0) \\ \forall k (P(k) \rightarrow P(k + 1)) \end{array}$$

---

$$\therefore \forall n P(n)$$

# using the induction rule in a formal proof

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$$\begin{array}{l} P(0) \\ \forall k (P(k) \rightarrow P(k + 1)) \end{array}$$

---

$$\therefore \forall n P(n)$$

1. Prove  $P(0)$
2. Let  $k$  be an arbitrary integer  $\geq 0$ 
  3. Assume that  $P(k)$  is true
  4. ...
  5. Prove  $P(k+1)$  is true
6.  $P(k) \rightarrow P(k+1)$
7.  $\forall k (P(k) \rightarrow P(k+1))$
8.  $\forall n P(n)$

Direct Proof Rule

Intro  $\forall$  from 2-6

Induction Rule 1&7

# format of an induction proof

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$$P(0)$$

$$\forall k (P(k) \rightarrow P(k + 1))$$

---

$$\therefore \forall n P(n)$$

1. Prove  $P(0)$

**Base Case**

2. Let  $k$  be an arbitrary integer  $\geq 0$

3. Assume that  $P(k)$  is true

**Inductive Hypothesis**

4. ...

5. Prove  $P(k+1)$  is true

**Inductive Step**

6.  $P(k) \rightarrow P(k+1)$

Direct Proof Rule

7.  $\forall k (P(k) \rightarrow P(k+1))$

Intro  $\forall$  from 2-6

8.  $\forall n P(n)$

Induction Rule 1&7

**Conclusion**

$$1 + 2 + 4 + 8 + \dots + 2^n$$

---

- $1 = 1$
- $1 + 2 = 3$
- $1 + 2 + 4 = 7$
- $1 + 2 + 4 + 8 = 15$
- $1 + 2 + 4 + 8 + 16 = 31$

$$1 + \dots + 2^n = 2^{n+1} - 1$$

Can we describe the pattern?

$$1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1$$

# proving $1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1$

---

- We could try proving it normally...
  - We want to show that  $1 + 2 + 4 + \dots + 2^n = 2^{n+1}$ .
  - So, what do we do now? We can sort of explain the pattern, but that's not a proof...
- We could prove it for  $n=1, n=2, n=3, \dots$  (individually), but that would literally take forever...



# inductive proof in five easy steps

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## Proof:

1. “We will show that  $P(n)$  is true for every  $n \geq 0$  by **induction.**”

2. “Base Case:” Prove  $P(0)$

3. “Inductive Hypothesis:”

Assume  $P(k)$  is true for some arbitrary integer  $k \geq 0$ ”

4. “Inductive Step:” Want to prove that  $P(k+1)$  is true:

Use the goal to figure out what you need.

Make sure you are using I.H. and point out where you are using it. (Don't assume  $P(k+1)$  !)

5. “Conclusion: Result follows by induction.”

proving  $1 + 2 + \dots + 2^n = 2^{n+1} - 1$

---

$$P(n) = "1 + 2 + \dots + 2^n = 2^{n+1} - 1"$$

Goal:  $\forall$  integers  $n \geq 0$ ,  $P(n)$

Base case:  $P(0) = "1 = 2^1 - 1"$  true b/c  $2 - 1 = 1$ .

IH: Assume  $1 + 2 + \dots + 2^k = 2^{k+1} - 1$   $P(k)$   
for some integer  $k \geq 0$ .

IS: From  $P(k)$ , we know  
 $1 + 2 + \dots + 2^k + 2^{k+1} = 2^{k+1} - 1 + 2^{k+1}$   
 $= 2^{k+2} - 1$ .

Hence  $P(k+1)$  holds.

Conclusion: Therefore by induction  $\forall n P(n)$ .

# proving $1 + 2 + \dots + 2^n = 2^{n+1} - 1$

---

1. Let  $P(n)$  be " $1 + 2 + \dots + 2^n = 2^{n+1} - 1$ ". We will show  $P(n)$  is true for all natural numbers by induction.
2. Base Case ( $n=0$ ):  $2^0 = 1 = 2 - 1 = 2^{0+1} - 1$
3. Induction Hypothesis: Suppose that  $P(k)$  is true for some arbitrary  $k \geq 0$ .
4. Induction Step:

Goal: Show  $P(k+1)$ , i.e. show  $1 + 2 + \dots + 2^k + 2^{k+1} = 2^{k+2} - 1$

$$1 + 2 + \dots + 2^k = 2^{k+1} - 1 \quad \text{by IH}$$

Adding  $2^{k+1}$  to both sides, we get:

$$1 + 2 + \dots + 2^k + 2^{k+1} = 2^{k+1} + 2^{k+1} - 1$$

Note that  $2^{k+1} + 2^{k+1} = 2(2^{k+1}) = 2^{k+2}$ .

So, we have  $1 + 2 + \dots + 2^k + 2^{k+1} = 2^{k+2} - 1$ , which is exactly  $P(k+1)$ .

5. Thus  $P(k)$  is true for all  $k \in \mathbb{N}$ , by induction.

## another example of a pattern

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- $2^0 - 1 = 1 - 1 = 0 = 3 \cdot 0$
- $2^2 - 1 = 4 - 1 = 3 = 3 \cdot 1$
- $2^4 - 1 = 16 - 1 = 15 = 3 \cdot 5$
- $2^6 - 1 = 64 - 1 = 63 = 3 \cdot 21$
- $2^8 - 1 = 256 - 1 = 255 = 3 \cdot 85$
- ...

$2^{2n} - 1$  is a multiple of 3

$$\forall n \geq 0$$

~~$$2^{2n} - 1 = (2 - 1)(2^{2n-1} + 2^{2n-2} + \dots + 1)$$~~

prove:  $3 \mid 2^{2n} - 1$  for all  $n \geq 0$

---

$$P(n) = "3 \mid 2^{2n} - 1"$$

Pf by induction that  $\forall n \geq 0$   $P(n)$

Base case:  $P(0)$ :  $3 \mid \underbrace{2^{2 \cdot 0} - 1}_0$  is true b/c  
 $0 = 3 \cdot 0$ .

IH: Assume that  $3 \mid 2^{2k} - 1$  for some  
int.  $k \geq 0$

IS: From (IH),  $2^{2k} - 1 = 3j$  for some int  $j$

$$\text{Thus } 2^2(2^{2k} - 1) = 2^2 \cdot 3j$$

$$\Rightarrow 2^2(2^{2k} - 1) = 12j \quad P(k+1)$$

$$\Rightarrow 2^{2(k+1)} - 4 = 12j \quad \uparrow\uparrow$$

$$\Rightarrow 2^{2(k+1)} - 1 = 12j + 3 = 3(4j + 1)$$

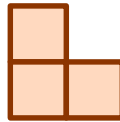
(Conclusion:  $3j$  incl  $\forall n P(n)$ )

$$\text{For all } n \geq 1: 1 + 2 + \cdots + n = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

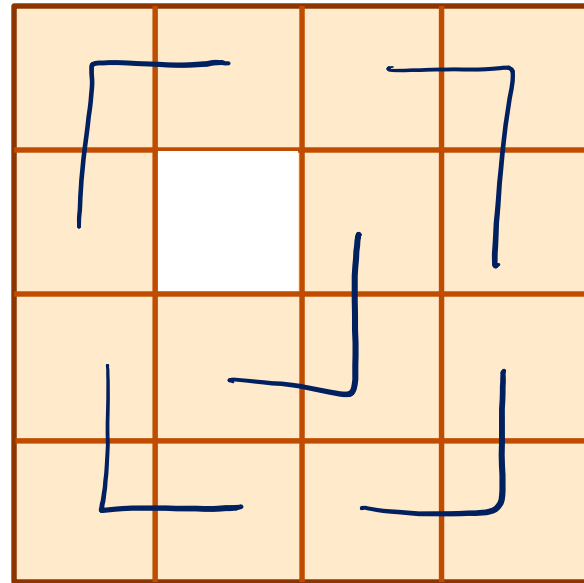
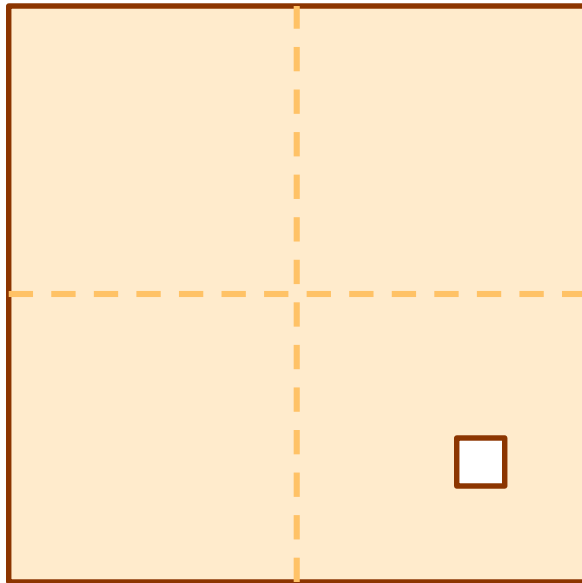
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# checkerboard tiling

Prove that a  $2^n \times 2^n$  checkerboard with one square removed can be tiled with:

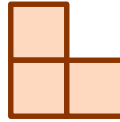



$P(0)$



# checkerboard tiling

Prove that a  $2^n \times 2^n$  checkerboard with one square removed can be tiled with:



$P(n)$  = "A  $2^n \times 2^n$  checkerboard w/ one square removed can be tiled by ."

Base case:  $P(0)$  We just have an empty  $1 \times 1$  board which is tiled w/out doing anything.

IH: Assume  $P(n)$  holds for some integer  $n \geq 0$ .

IS: Consider a  $2^{n+1} \times 2^{n+1}$  board with one square removed.

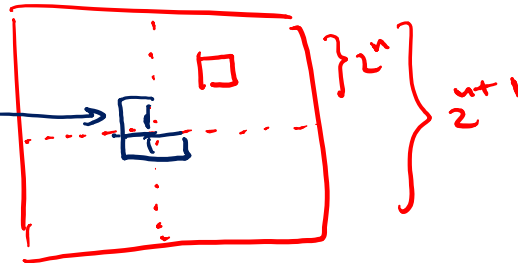
By symmetry, we can assume:

Now place a block at the intersection 

We now have four  $2^n \times 2^n$  squares

$S_1, S_2, S_3, S_4$

each with one block removed.



Thus the  $2^n \times 2^n$  square can be tiled.

$\Rightarrow P(n+1)$

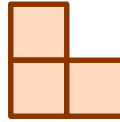
By IH, each of  $S_1, S_2, S_3, S_4$  can be tiled.

(cont.)

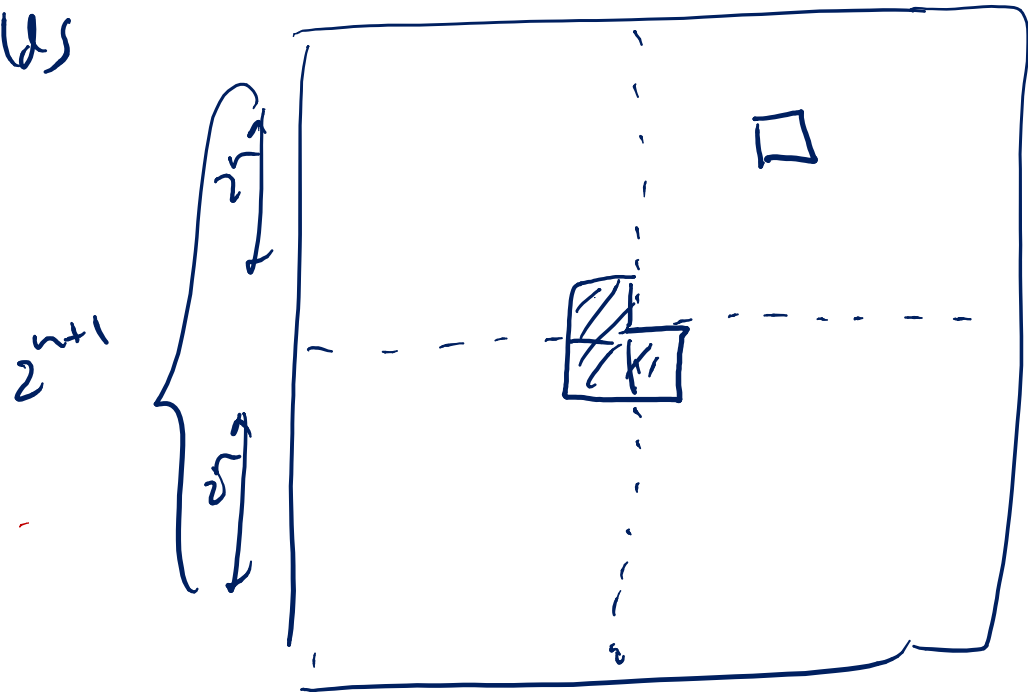
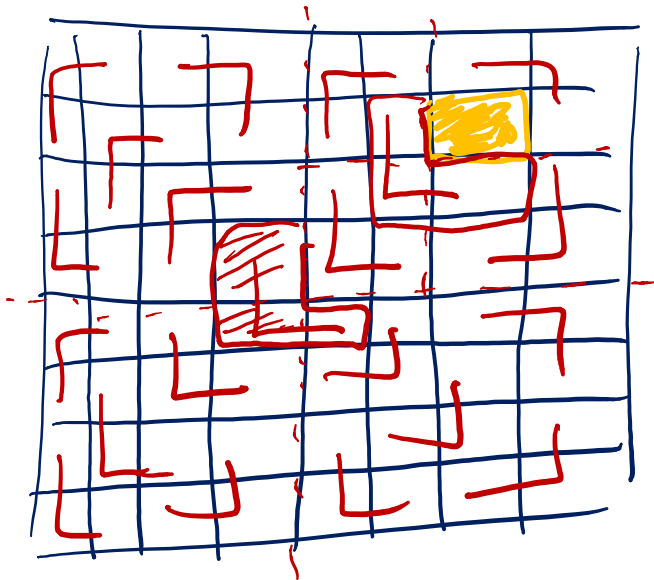


# checkerboard tiling

Prove that a  $2^n \times 2^n$  checkerboard with one square removed can be tiled with:



Assume it holds  
for  $2^n \times 2^n$



$\forall n \ P(n) \rightarrow P(n+1)$

prove:  $n^n \geq n!$  for all  $n \geq 1$

$$P(n) = "n^n \geq n!"$$

$$n! = n(n-1) \cdots 3 \cdot 2 \cdot 1$$

Base case:  $P(1)$  is true:

$$1 = 1! \geq 1! = 1$$

IH: Assume that  $k^k \geq k!$   
for some  $k \geq 1$ .

IS:  $k^k \geq k! \Rightarrow (k+1)k^k \geq (k+1)k! = (k+1)!$

$$(k+1)^{k+1} = (k+1) \cdot (k+1)^k \geq (k+1)k^k$$

Combine w/ prev. line  $(k+1)^{k+1} \geq k!$

Conclusion: By induction  $\forall n P(n)$ .

~~$Q(n) = P(n+1)$~~

~~$\forall n \geq 0 Q(n)$~~

~~$\Rightarrow \forall n \geq 1 P(n)$~~