Power set of a set A = set of all subsets of A

$$\boxed{\mathcal{P}(A) = \{\, B : B \subseteq A \,\}}$$

e.g. $\text{Days} = \{M, W, F\}$

$$\begin{aligned}\mathcal{P}(\text{Days}) = \{\, &\varnothing, \\ &\{M\}, \{W\}, \{F\}, \\ &\{M, W\}, \{W, F\}, \{M, F\}, \\ &\{M, W, F\} \quad \}\end{aligned}$$

e.g. $\mathcal{P}(\varnothing) = \{\varnothing\}$

Fall 2015

Lecture 10: Functions, Modular arithmetic

[this special lecture was given by a 5-year-old]

a little recap

So far:

- Propositional logic
- Logic to build circuits
- Predicates and quantifiers
- Proof systems and logical inference
- Basic set theory

Question: If the domain of discourse is empty and P is a predicate, what is the truth value of:

$\exists x\ P(x)$

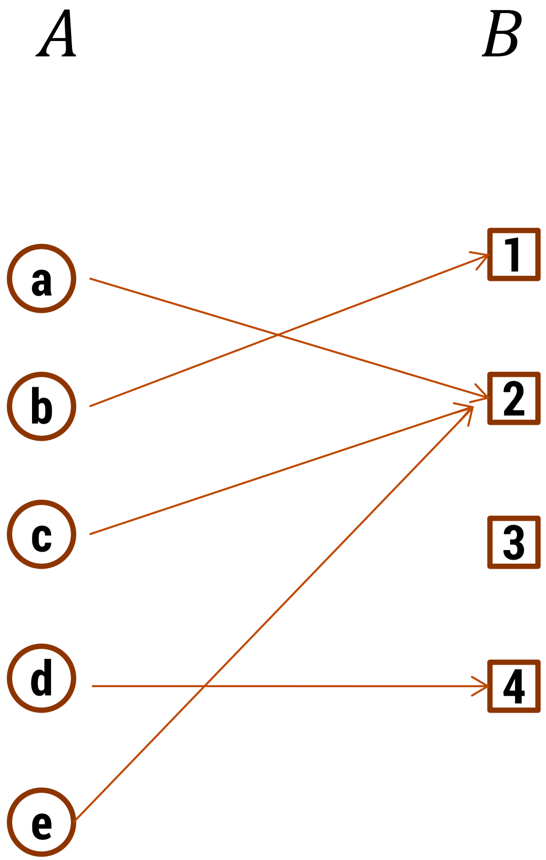
$\forall x\ P(x)$

A **function** from A to B:

- Every element of A is assigned to exactly one element of B.
- We write $f : A \to B$.
- "Image of X under f" $= "f(X)"$
  $= \{x : \exists y \left(y \in X \land x = f(y)\right)\}$

- **Domain** of f is A
- **Codomain** of f is B
- **Image** of f = Image of domain under f
  = all the elements pointed to by something
  in the domain.

image



Image({a}) =

Image({a, e}) =

Image({a, b}) =

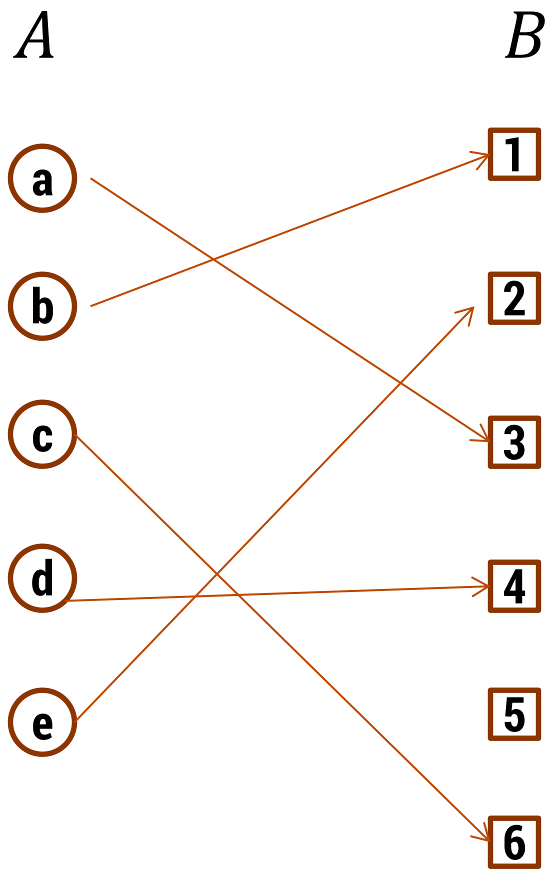
Image(A) =

injections and surjections

A **function** $f : A \to B$ is **one-to-one** (or, **injective)** if every output corresponds to at most one input, i.e. $f(x) = f(x') \Rightarrow x = x'$ for all $x, x' \in A$.

A **function** $f : A \to B$ is **onto** (or, **surjective)** if every output gets hit, i.e. for every $y \in B$, there exists $x \in A$ such that $f(x) = y$.

is this function one-to-one? is it onto?



It is one-to-one, because nothing in B is pointed to by multiple elements of A.

It is not onto, because 5 is not pointed to by anything.

QUIZ!

	One-to-one (?)	Onto (?)
$x \mapsto x^2$		
$x \mapsto x^3 - x$		
$x \mapsto e^x$		
$x \mapsto x^3$		

Domain: Reals

a harder quiz

Dear HBO, this is a slide about digital watermarking.

“number theory” (and applications to computing)

- How whole numbers work

  [fascinating, deep, weird area of mathematics that no one understands, but the basics are easy and really useful]

- Many significant applications
  - Cryptography [this is how SSL works]
  - Hashing
  - Security
  - Error-correcting codes [this is how your bluray player works]

- Important tool set

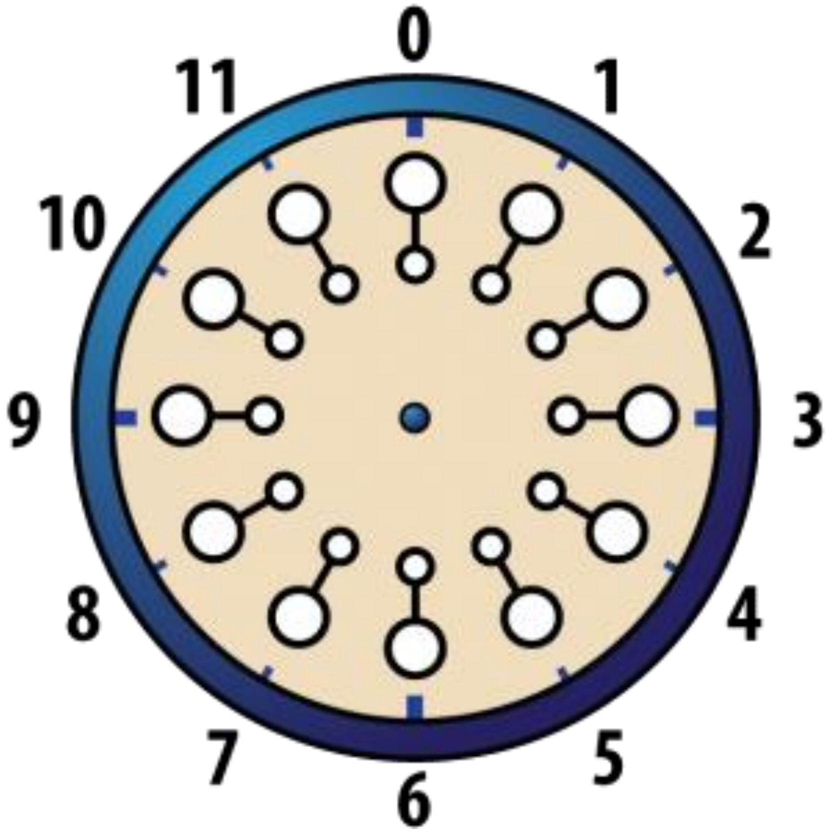
thanks, java

```
public class Test {
    final static int SEC_IN_YEAR = 364*24*60*60;
    public static void main(String args[]) {
        System.out.println(
            “I will be alive for at least ” +
            SEC_IN_YEAR * 101 + “ seconds.”
        );
    }
}
```

```
 ----jGRASP exec: java Test
I will be alive for at least -186619904 seconds.

 ----jGRASP: operation complete.
```

Arithmetic over a finite domain: Math with wrap around



Integers a, b, with $a \neq 0$. We say that a **divides** b iff there is an integer k such that $b = k\,a$. The notation $a \mid b$ denotes "a divides b."

Let a be an integer and d a positive integer. Then there are *unique* integers q and r, with $0 \leq r < d$, such that $a = d\,q + r$.

$q = a$ **div** d $r = a$ **mod** d

Note: r ≥ 0 even if a < 0.
Not quite the same as `a % d`.

arithmetic mod 7

$$a +_7 b = (a + b) \bmod 7$$

$$a \times_7 b = (a \times b) \bmod 7$$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

modular congruence

Let a and b be integers, and m be a positive integer.
We say *a* is **congruent** to *b* **modulo** *m* if *m* divides *a* – *b*.
We use the notation a ≡ b (mod m) to indicate that a is congruent to b modulo m.

modular arithmetic: examples

A ≡ 0 (mod 2)

This statement is the same as saying "A is even"; so, any A that is even (including negative even numbers) will work.

1 ≡ 0 (mod 4)

This statement is false. If we take it mod 1 instead, then the statement is true.

A ≡ -1 (mod 17)

If A = 17x – 1 = 17(x-1) + 16 for an integer x, then it works.

Note that (m – 1) mod m

= ((m mod m) + (-1 mod m)) mod m

= (0 + -1) mod m

= -1 mod m

modular arithmetic can haz sense

Theorem: Let a and b be integers, and let m be a positive integer. Then $a \equiv b \pmod{m}$ if and only if $a \bmod m = b \bmod m$.

Proof: Suppose that $a \equiv b \pmod{m}$.
By definition: $a \equiv b \pmod{m}$ implies $m \mid (a - b)$
which by definition implies that $a - b = km$ for some integer k.
Therefore $a = b + km$.
Taking both sides modulo m we get
$a \bmod m = (b+km) \bmod m = b \bmod m$

modular arithmetic can haz sense

Theorem: Let a and b be integers, and let m be a positive integer. Then $a \equiv b \pmod{m}$ if and only if a mod m = b mod m.

Proof: Suppose that a mod m = b mod m.

By the division theorem, $a = mq + (a \bmod m)$ and
$b = ms + (b \bmod m)$ for some integers q,s.

$$
\begin{aligned}
a - b &= (mq + (a \bmod m)) - (mr + (b \bmod m)) \\
&= m(q - r) + (a \bmod m - b \bmod m) \\
&= m(q - r) \quad \text{since } a \bmod m = b \bmod m
\end{aligned}
$$

Therefore $m \mid (a-b)$ and so $a \equiv b \pmod{m}$

consistency of addition

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $\mathbf{a + c \equiv b + d \pmod{m}}$

Suppose $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$.

Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Adding the equations together gives us $(a + c) - (b + d) = m(k + j)$. Now, re-applying the definition of mod gives us $a + c \equiv b + d \pmod{m}$.

consistency of multiplication

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $\mathbf{ac \equiv bd \pmod{m}}$

Suppose $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$.

Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Then, $a = km + b$ and $c = jm + d$.

Multiplying both together gives us

$$ac = (km + b)(jm + d) = kjm^2 + kmd + jmb + bd$$

Rearranging gives us $ac - bd = m(kjm + kd + jb)$.

Using the definition of mod gives us $ac \equiv bd \pmod{m}$.

Let n be an integer.
Prove that $n^2 \equiv 0 \pmod 4$ or $n^2 \equiv 1 \pmod 4$