

Power set of a set A = set of all subsets of A

$$\mathcal{P}(A) = \{ B : B \subseteq A \}$$

e.g. Days = $\{M, W, F\}$

$$A \subseteq A$$

$$\begin{aligned} \mathcal{P}(\text{Days}) = \{ & \emptyset, \\ & \{M\}, \{W\}, \{F\}, \\ & \{M, W\}, \{W, F\}, \{M, F\}, \\ & \{M, W, F\} \} \end{aligned}$$

e.g. $\mathcal{P}(\emptyset) = \{\emptyset\} \neq \emptyset$

$$\emptyset \subseteq \emptyset$$

$$\{P(A)\} \subseteq P(P(A))$$

review quiz

$$\underline{P(A) \subseteq P(P(A))}$$

false

$$\underline{P(A) \in P(P(A))}$$

true

If $|A| = n$, what is $|P(A)|$?

$$|P(A)| = 2^k \quad \text{size of } |A \times A \times A \times A|?$$

$$|P(P(A))| = 2^{(2^k)}$$

n^4

Since $B \subseteq B$,
 $B \in P(B)$

$$B = P(A)$$

$$|A \times A| = n^2$$

$$A = \{1\}$$

$$P(A) = \{\emptyset, \{1\}\}$$

$$P(P(A)) = \{\emptyset, \{\emptyset\}, \{\{1\}\}, \{\emptyset, \{1\}\}\}$$

cse 311: foundations of computing

Fall 2015

Lecture 10: Functions, Modular arithmetic



So far:

- Propositional logic
- Logic to build circuits
- Predicates and quantifiers
- Proof systems and logical inference
- Basic set theory

empty domains

Question: If the domain of discourse is empty and P is a predicate, what is the truth value of:

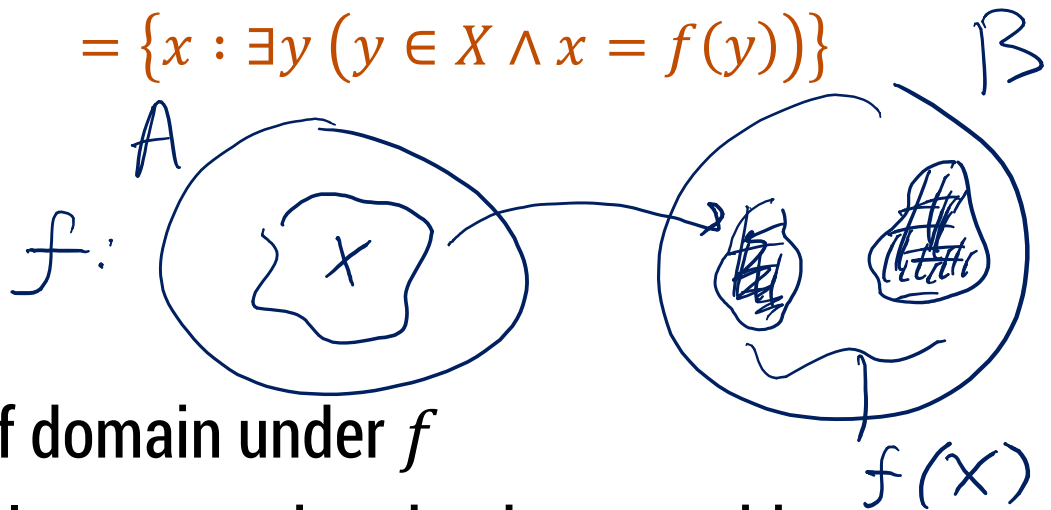
$\exists x P(x)$ F

$\forall x P(x)$ T

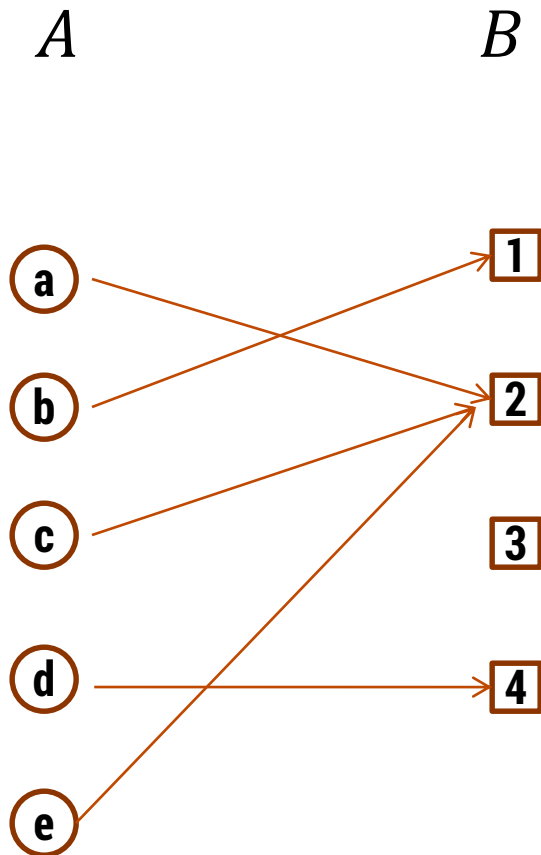
A **function** from A to B :

- Every element of A is assigned to exactly one element of B .
- We write $f : A \rightarrow B$.
- “Image of X under f ” = “ $f(X)$ ”

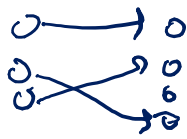
$$= \{x : \exists y (y \in X \wedge x = f(y))\}$$



- **Domain** of f is A
- **Codomain** of f is B
- **Image** of f = Image of domain under f
= all the elements pointed to by something in the domain.



$$\begin{aligned} \text{Image}(\{a\}) &= \{2\} \\ \text{Image}(\{a, e\}) &= \{2\} \\ \text{Image}(\{a, b\}) &= \{1, 2\} \\ \text{Image}(A) &= \{1, 2, 4\} \end{aligned}$$

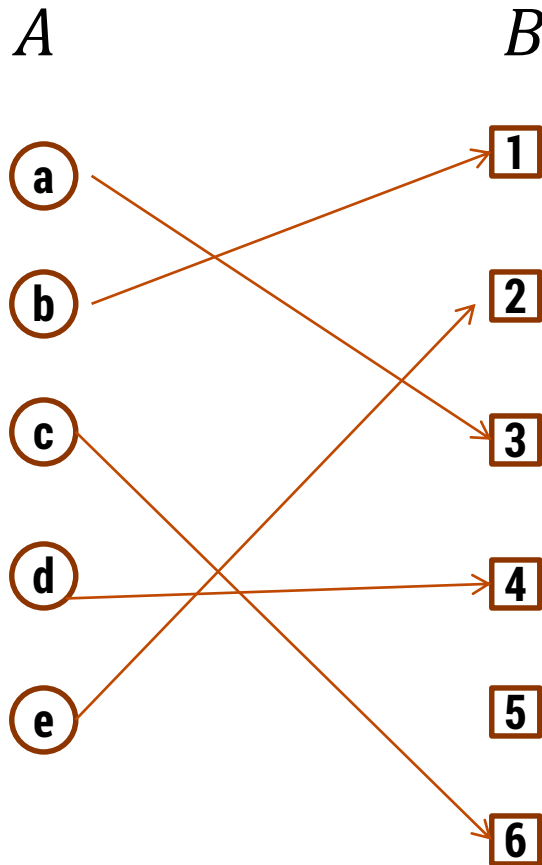


injections and surjections

A function $f : A \rightarrow B$ is **one-to-one** (or, **injective**) if every output corresponds to at most one input, i.e. $f(x) = f(x') \Rightarrow x = x'$ for all $x, x' \in A$.

A function $f : A \rightarrow B$ is **onto** (or, **surjective**) if every output gets hit, i.e. for every $y \in B$, there exists $x \in A$ such that $f(x) = y$.

is this function one-to-one? is it onto?



injective

One-to-one



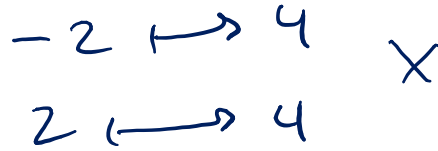
It is ~~one-to-one~~ ^{injective}, because ~~nothing~~ ^{nothing} in B is pointed to by multiple elements of A.

It is not onto, because 5 is not pointed to by anything.

One-to-one (?)

Onto (?)

$x \mapsto x^2$

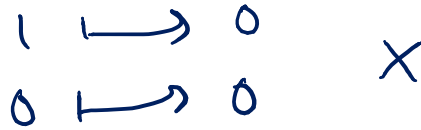


X

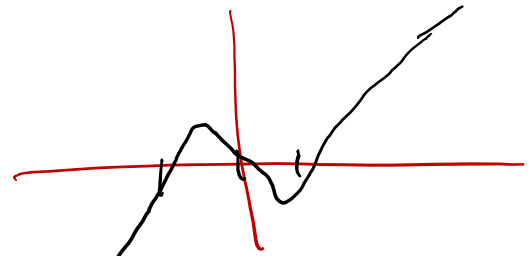
doesn't hit
neg numbers

X

$x \mapsto x^3 - x$

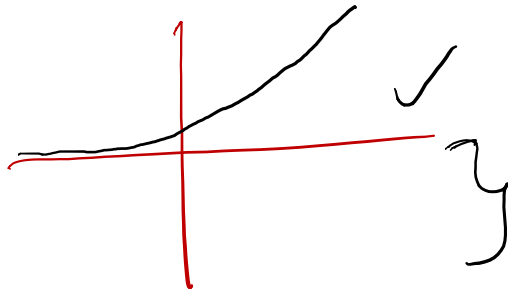


X



✓

$x \mapsto e^x$



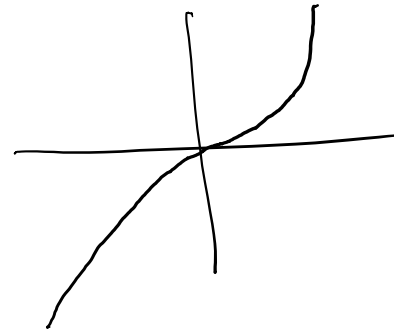
✓

X

$x \mapsto x^3$

$x \mapsto |x|^{1/3} \text{ sign}(x)$

✓



✓

Co domain

Domain: Reals

Dear HBO, this is a slide about digital watermarking.



“number theory” (and applications to computing)

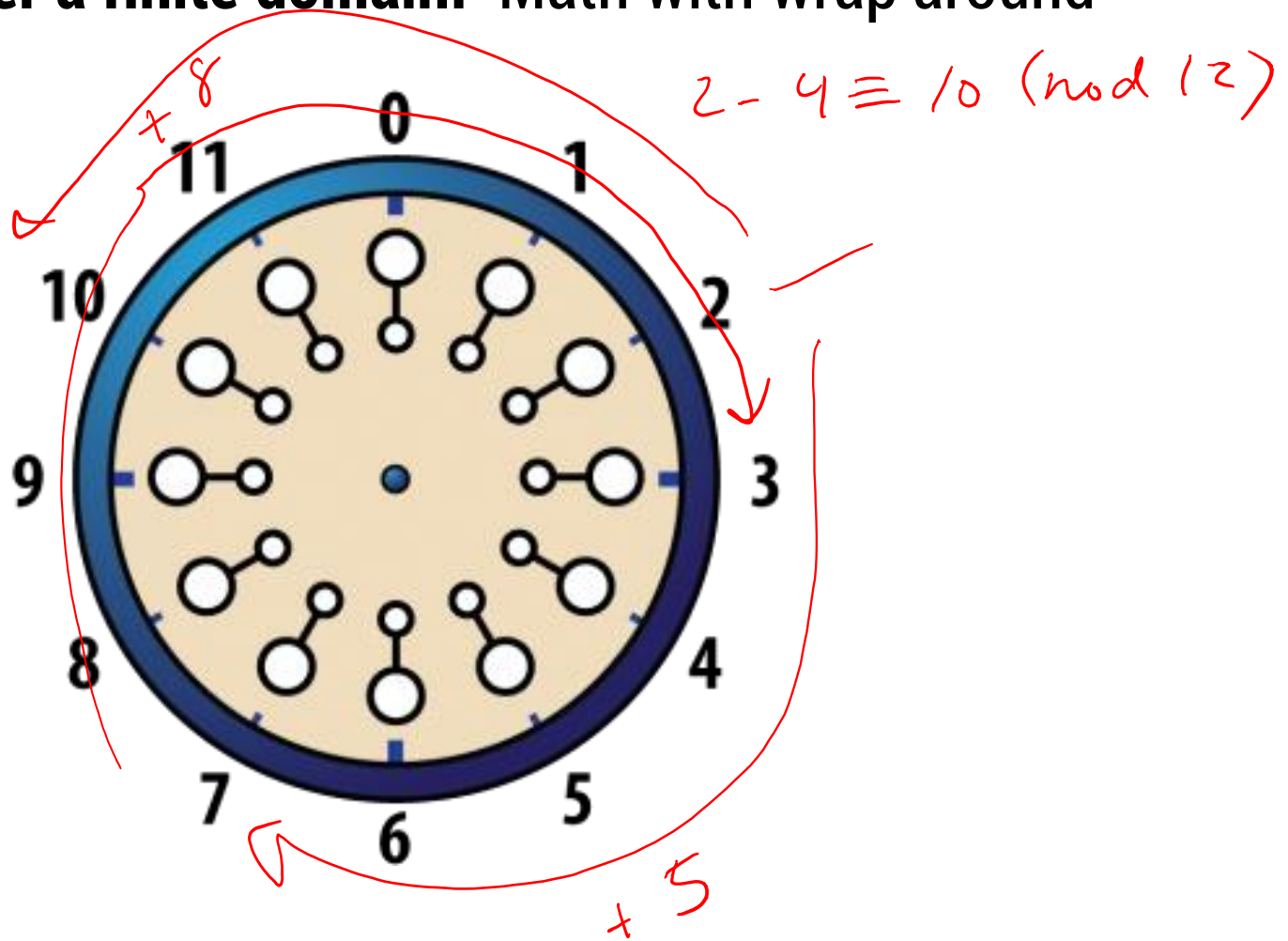
- **How whole numbers work**
[fascinating, deep, weird area of mathematics that no one understands, but the basics are easy and really useful]
- **Many significant applications**
 - Cryptography [this is how SSL works]
 - Hashing
 - Security
 - Error-correcting codes [this is how your bluray player works]
- **Important tool set**

thanks, java

```
public class Test {
    final static int SEC_IN_YEAR = 364*24*60*60;
    public static void main(String args[]) {
        System.out.println(
            "I will be alive for at least " +
            SEC_IN_YEAR * 101 + " seconds."
        );
    }
}
```

```
----jGRASP exec: java Test
I will be alive for at least -186619904 seconds.
----jGRASP: operation complete.
```

Arithmetic over a finite domain: Math with wrap around



divisibility

Integers a, b , with $a \neq 0$. We say that a **divides** b iff there is an integer k such that $b = k a$. The notation $a \mid b$ denotes “ a divides b .”

$$3 \mid 15$$

$$1 \mid 15$$

$$15 \nmid 17$$

$$\begin{array}{c} a \\ 3 \end{array} \mid \begin{array}{c} b \\ 0 \end{array} \quad ?$$

$$0 = 0 \cdot 3$$

$$b = k a$$

division theorem

Let a be an integer and d a positive integer. Then there are *unique* integers q and r , with $0 \leq r < d$, such that $a = d q + r$.

$$q = a \text{ div } d$$

$$r = a \text{ mod } d$$

$$a = 15$$

$$d = 4$$

$$q = 3$$

$$r = 3$$

$$a = -13$$

$$d = 4$$

$$q = -4$$

$$r = 3$$

Note: $r \geq 0$ even if $a < 0$.

Not quite the same as $a \% d$.

arithmetic mod 7

$$a +_7 b = (a + b) \bmod 7$$

$$a \times_7 b = (a \times b) \bmod 7$$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

x	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Handwritten annotations:
-2 -1 (above the 5 and 6 columns)
-2 (to the left of the 5 row)
-1 (to the left of the 6 row)

modular congruence

Let a and b be integers, and m be a positive integer. We say a is **congruent** to b **modulo** m if m divides $a - b$. We use the notation $a \equiv b \pmod{m}$ to indicate that a is congruent to b modulo m .

$$a \equiv b \pmod{m} \iff m \mid a - b$$

modular arithmetic: examples

$$A \equiv 0 \pmod{2}$$

This statement is the same as saying “A is even”; so, any A that is even (including negative even numbers) will work.

$$1 \equiv 0 \pmod{4}$$

This statement is false. If we take it mod 1 instead, then the statement is true.

$$A \equiv -1 \pmod{17}$$

If $A = 17x - 1 = 17(x-1) + 16$ for an integer x , then it works.

Note that $(m - 1) \pmod{m}$

$$= ((m \pmod{m}) + (-1 \pmod{m})) \pmod{m}$$

$$= (0 + -1) \pmod{m}$$

$$= -1 \pmod{m}$$

modular arithmetic can haz sense

Theorem: Let a and b be integers, and let m be a positive integer. Then $a \equiv b \pmod{m}$ if and only if $a \bmod m = b \bmod m$.

Proof: Suppose that $a \equiv b \pmod{m}$.

By definition: $a \equiv b \pmod{m}$ implies $m \mid (a - b)$

which by definition implies that $a - b = km$ for some integer k .

Therefore $a = b + km$.

Taking both sides modulo m we get

$$a \bmod m = (b+km) \bmod m = b \bmod m$$

modular arithmetic can haz sense

Theorem: Let a and b be integers, and let m be a positive integer. Then $a \equiv b \pmod{m}$ if and only if $a \bmod m = b \bmod m$.

Proof: Suppose that $a \bmod m = b \bmod m$.

By the division theorem, $a = mq + (a \bmod m)$ and

$b = ms + (b \bmod m)$ for some integers q, s .

$$a - b = (mq + (a \bmod m)) - (ms + (b \bmod m))$$

$$= m(q - s) + (a \bmod m - b \bmod m)$$

$$= m(q - s) \text{ since } a \bmod m = b \bmod m$$

Therefore $m \mid (a-b)$ and so $a \equiv b \pmod{m}$

consistency of addition

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then **$a + c \equiv b + d \pmod{m}$**

Suppose $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$.

Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Adding the equations together gives us

$(a + c) - (b + d) = m(k + j)$. Now, re-applying the definition of mod gives us $a + c \equiv b + d \pmod{m}$.

consistency of multiplication

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then **$ac \equiv bd \pmod{m}$**

Suppose $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$.

Unrolling definitions gives us some k such that $a - b = km$, and some j such that $c - d = jm$.

Then, $a = km + b$ and $c = jm + d$.

Multiplying both together gives us

$$ac = (km + b)(jm + d) = kjm^2 + kmd + jmb + bd$$

Rearranging gives us $ac - bd = m(kjm + kd + jb)$.

Using the definition of mod gives us **$ac \equiv bd \pmod{m}$** .

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$