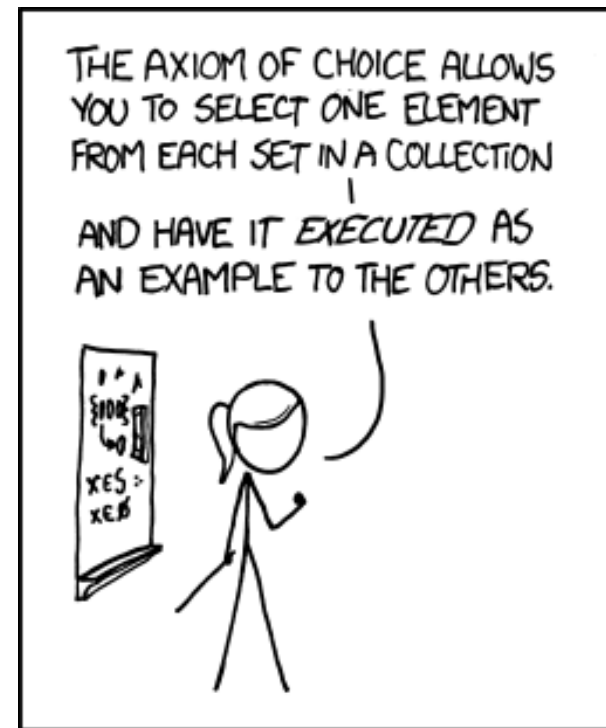


cse 311: foundations of computing

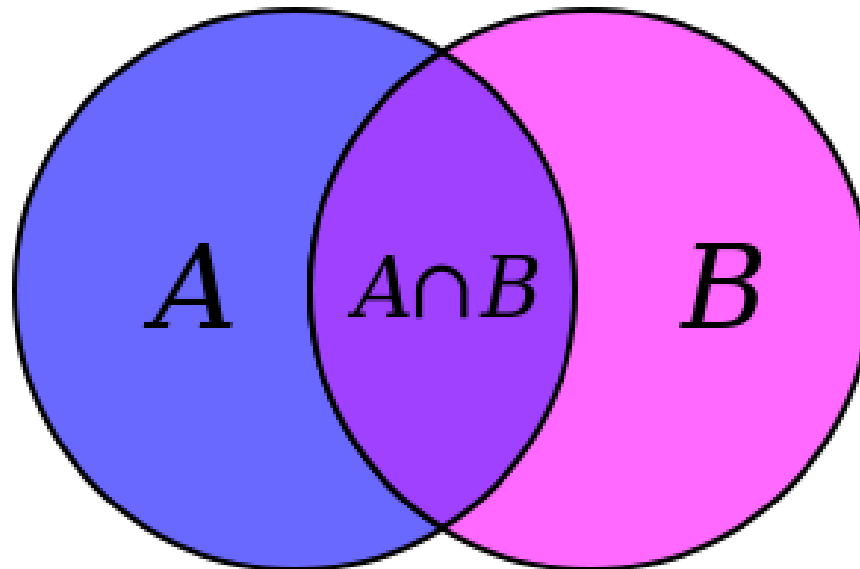
Fall 2015

Lecture 9: Set theory



MY MATH TEACHER WAS A BIG
BELIEVER IN PROOF BY INTIMIDATION.

- Formal treatment dates from late 19th century
- Direct ties between set theory and logic
- Important foundational language



some common sets

\mathbb{N} is the set of **Natural numbers**; $\mathbb{N} = \{0, 1, 2, \dots\}$

\mathbb{Z} is the set of **Integers**; $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$

\mathbb{Q} is the set of **Rational numbers**; e.g. $\frac{1}{2}$, -17, $\frac{32}{48}$

\mathbb{R} is the set of **Real numbers**; e.g. 1, -17, $\frac{32}{48}$, π

$[n]$ is the set $\{1, 2, \dots, n\}$ when n is a natural number

$\{\} = \emptyset$ is the **empty set**; the *only* set with no elements

EXAMPLES

Are these sets?

$A = \{1, 1\}$

$B = \{1, 3, 2\}$

$C = \{\square, 1\}$

$D = \{\{\}, 17\}$

$E = \{1, 2, 7, \text{cat}, \text{dog}, \emptyset, \alpha\}$

Set membership:

We write $2 \in E$; $3 \notin E$.

- A and B are *equal* if they have the same elements

$$A = B \equiv \forall x (x \in A \leftrightarrow x \in B)$$

- A is a *subset* of B if every element of A is also in B

$$A \subseteq B \equiv \forall x (x \in A \rightarrow x \in B)$$

$$A = \{1, 2, 3\}$$

$$B = \{3, 4, 5\}$$

$$C = \{3, 4\}$$

QUESTIONS

$$\emptyset \subseteq A? \quad \checkmark$$

$$A \subseteq B? \quad \times$$

$$C \subseteq B \quad \checkmark$$

- A and B are *equal* if they have the same elements

$$A = B \equiv \forall x (x \in A \leftrightarrow x \in B)$$

- A is a *subset* of B if every element of A is also in B

$$A \subseteq B \equiv \forall x (x \in A \rightarrow x \in B)$$

- Note:

$$(A = B) \equiv (A \subseteq B) \wedge (B \subseteq A)$$

building sets from predicates

- The following says “S is the set of all x’s where P(x) is true.”

$$S = \{x : P(x)\}$$

- The following says “S is the set of those elements of A for which P(x) is true.”

$$S = \{x \in A : P(x)\}$$

- “The set of all the real numbers less than one”

$$\{x \in \mathbb{R} : x < 1\}$$

- “The set of all powers of two”

$$\{x \in \mathbb{N} : \exists j (x = 2^j)\}$$

set operations

$$A \cup B = \{ x : (x \in A) \vee (x \in B) \}$$

Union

$$A \cap B = \{ x : (x \in A) \wedge (x \in B) \}$$

Intersection

$$A \setminus B = \{ x : (x \in A) \wedge (x \notin B) \}$$

Set difference

$$A = \{1, 2, 3\}$$

$$B = \{4, 5, 6\}$$

$$C = \{3, 4\}$$

QUESTIONS

Using A, B, C and set operations, make...

$$[6] = ? \quad A \cup B$$

$$\{3\} = ? \quad A \cap C$$

$$\{1, 2\} = ? \quad A \setminus C$$

$$\{1, 3\} = ? \quad (\text{Can we ever separate 1 and 2?})$$

more set operations

$$A \oplus B = \{ x : (x \in A) \oplus (x \in B) \}$$

Symmetric
difference

$$\bar{A} = \{ x : x \notin A \}$$

(with respect to universe U)

Complement

$$A = \{1, 2, 3\}$$

$$B = \{1, 4, 2, 6\}$$

$$C = \{1, 2, 3, 4\}$$

QUESTIONS

Let $S = \{1, 2\}$.

If the universe is A, then \bar{S} is...

$\{3\}$

If the universe is B, then \bar{S} is...

$\{4, 6\}$

If the universe is C, then \bar{S} is...

$\{3, 4\}$

it's Boolean algebra again! (yay...?)

- Definition for \cup based on \vee

$$A \cup B = \{x : x \in A \vee x \in B\}$$

- Definition for \cap based on \wedge

$$A \cap B = \{x : x \in A \wedge x \in B\}$$

- Complement works like \neg

$$\bar{A} = \{x : \neg(x \in A)\}$$

empty set and power set

Power set of a set A = set of all subsets of A

$$\mathcal{P}(A) = \{ B : B \subseteq A \}$$

e.g. Days = $\{M, W, F\}$

$$\begin{aligned} \mathcal{P}(\text{Days}) = \{ & \emptyset, \\ & \{M\}, \{W\}, \{F\}, \\ & \{M, W\}, \{W, F\}, \{M, F\}, \\ & \{M, W, F\} \} \end{aligned}$$

e.g. $\mathcal{P}(\emptyset) = ?$

$$(\mathcal{P}(\emptyset) = \{\emptyset\})$$

cartesian product

$$A \times B = \{ (a, b) : a \in A, b \in B \}$$

$$A = \{1, 2\}$$

$$B = \{a, b, c\}$$

$$A \times B = \{ (1, a), (1, b), (1, c), \\ (2, a), (2, b), (2, c) \}$$

$$A \times \emptyset = \emptyset$$

$$|A \times B| = |A| \cdot |B|$$

de Morgan's laws

$$\overline{A \cup B} = \bar{A} \cap \bar{B}$$

$$\begin{aligned} x \in \overline{A \cup B} &\leftrightarrow \neg(x \in A \vee x \in B) \\ &\leftrightarrow (\neg(x \in A) \wedge \neg(x \in B)) \\ &\leftrightarrow (x \in \bar{A} \wedge x \in \bar{B}) \leftrightarrow x \in \bar{A} \cap \bar{B} \end{aligned}$$

$$\overline{A \cap B} = \bar{A} \cup \bar{B}$$

$$\begin{aligned} x \in \overline{A \cap B} &\leftrightarrow \neg(x \in A \wedge x \in B) \\ &\leftrightarrow (\neg(x \in A) \vee \neg(x \in B)) \\ &\leftrightarrow (x \in \bar{A} \vee x \in \bar{B}) \\ &\leftrightarrow x \in \bar{A} \cup \bar{B} \end{aligned}$$

Proof technique:

To show $C = D$ show

$x \in C \rightarrow x \in D$ and

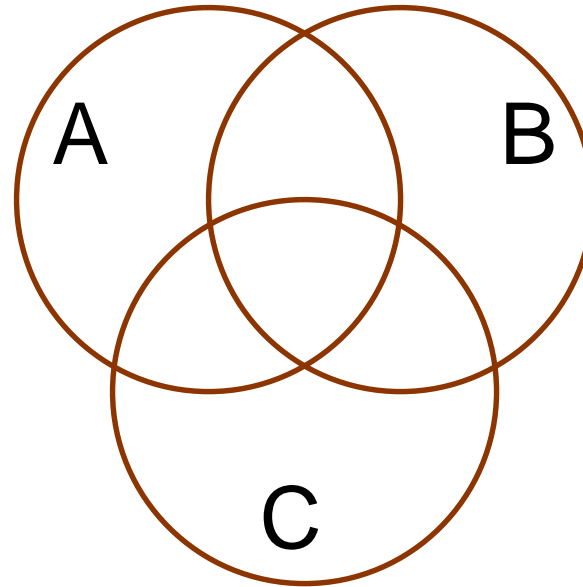
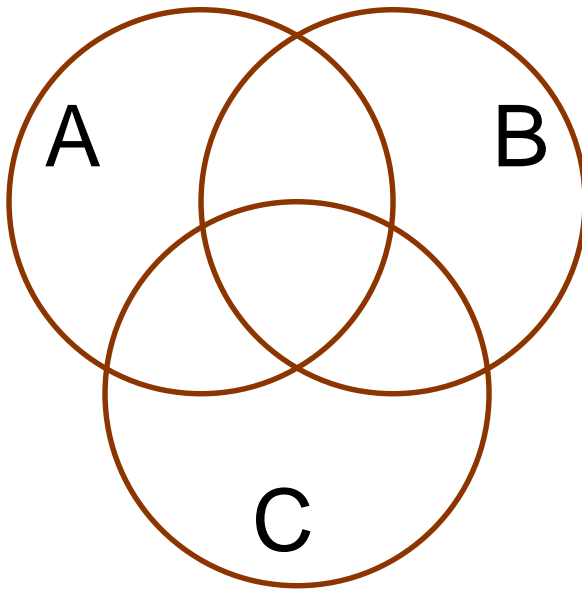
$x \in D \rightarrow x \in C$

distributive laws

just like $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

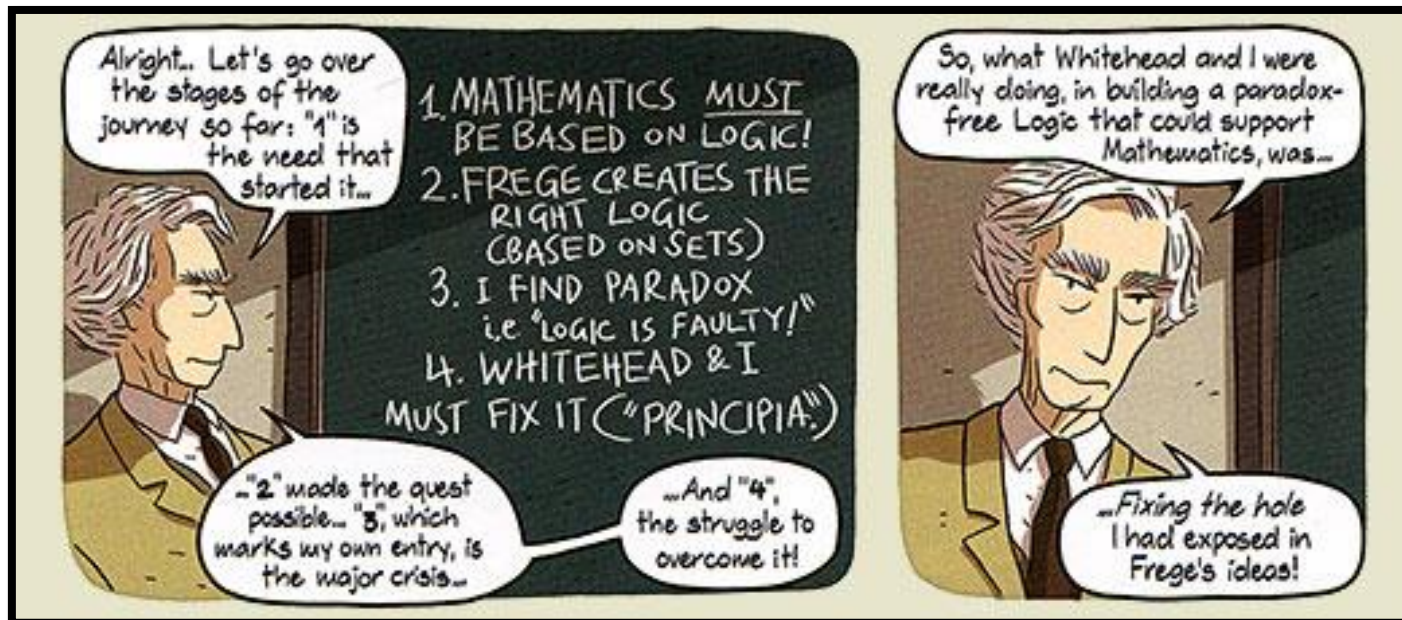
$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$



Russell's paradox

$$S = \{ x : x \notin x \}$$



representing sets using bits

- Suppose universe U is $\{1, 2, \dots, n\}$
- Can represent set $B \subseteq U$ as a vector of bits:
 $b_1 b_2 \cdots b_n$ where $b_i = 1$ when $i \in B$
 $b_i = 0$ when $i \notin B$
 - Called the *characteristic vector* of set B
- Given characteristic vectors for A and B
 - What is characteristic vector for $A \cup B$? $A \cap B$?

unix/linux file permissions

- `ls -l`
`drwxr-xr-x ... Documents/`
`-rw-r--r-- ... file1`
- Permissions maintained as bit vectors
 - Letter means bit is 1
 - “--” means bit is 0.

bitwise operations

$$\begin{array}{r} 01101101 \\ \vee \quad 00110111 \\ \hline 01111111 \end{array}$$

Java: $z = x | y$

$$\begin{array}{r} 00101010 \\ \wedge \quad 00001111 \\ \hline 00001010 \end{array}$$

Java: $z = x \& y$

$$\begin{array}{r} 01101101 \\ \oplus \quad 00110111 \\ \hline 01011010 \end{array}$$

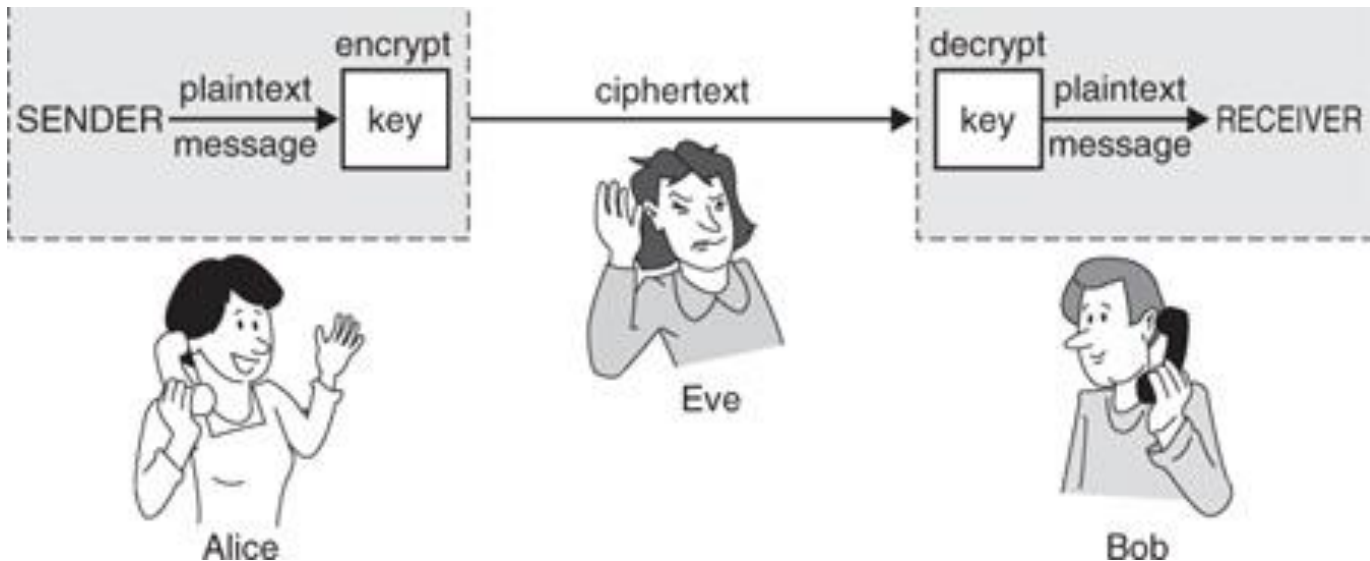
Java: $z = x \wedge y$

a useful identity

- If x and y are bits: $(x \oplus y) \oplus y = ?$ *X*
- What if x and y are bit-vectors? *Same thing,
bitwise*

private key cryptography

- Alice wants to communicate message secretly to Bob so that eavesdropper Eve who hears their conversation cannot tell what Alice's message is.
- Alice and Bob can get together and privately share a secret key **K** ahead of time.



- Alice and Bob privately share random n-bit vector K
 - Eve does not know K
- Later, Alice has n-bit message m to send to Bob
 - Alice computes $C = m \oplus K$
 - Alice sends C to Bob
 - Bob computes $m = C \oplus K$ which is $(m \oplus K) \oplus K$
- Eve cannot figure out m from C unless she can guess K

