

Homework #1 Due Today at 11:59pm

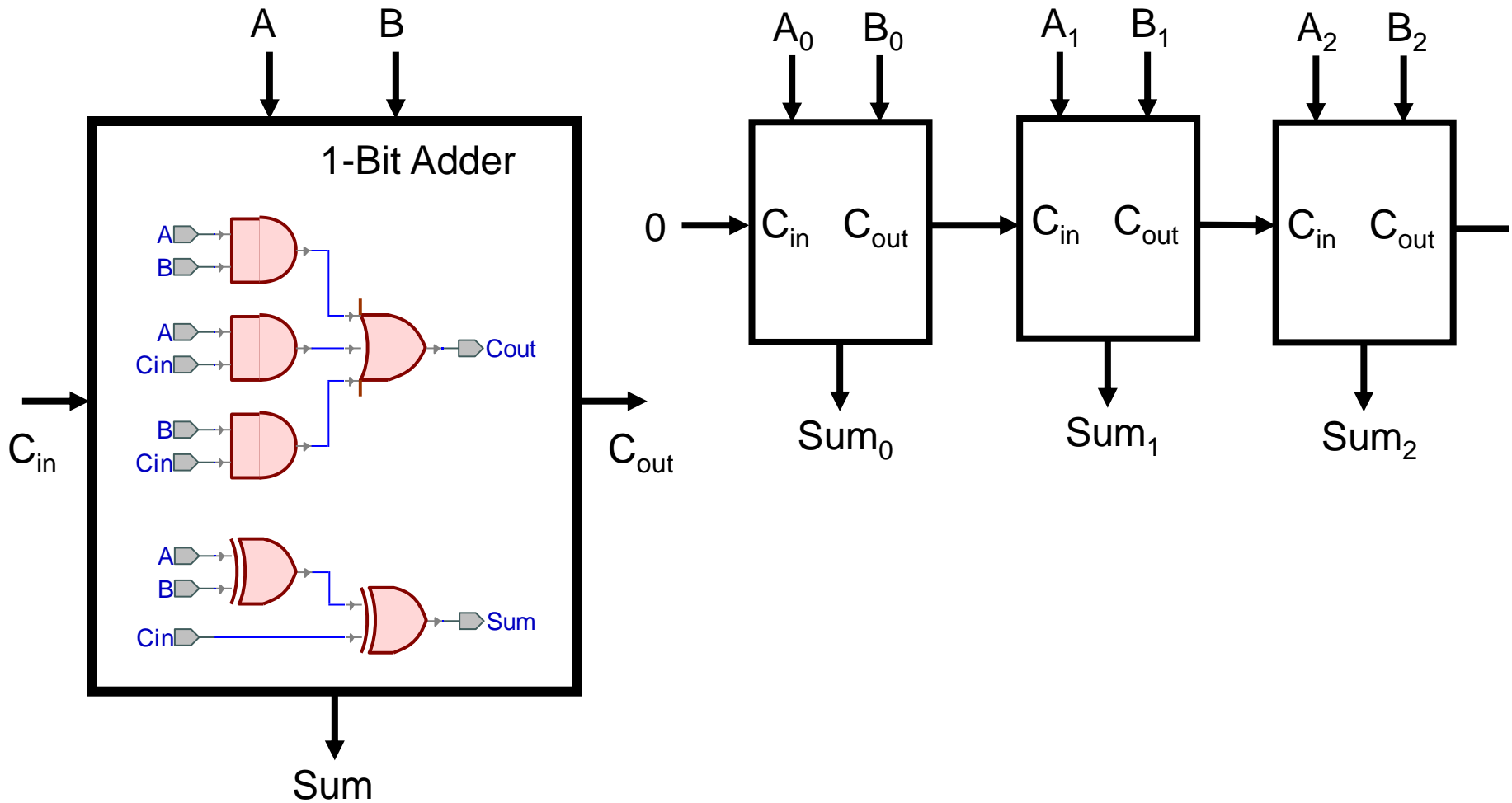
Your Gradescope account is created by your UW/CSE email address

Homework #2 will be posted today and it is due next Friday

TA Office Hours

TA	Office hours	Room
Sam Castle	Wed, 12:00-1:00	CSE 021
Jiechen Chen	Tue, 4:00-5:00	CSE 218
Rebecca Leslie	Mon, 8:30-9:30	CSE 218
Evan McCarty	Tue, 11:30-12:30	CSE 220
Tim Oleskiw	Tue, 3:00-4:00	CSE 218
Spencer Peters	Tue, 1:00-2:00	CSE 218
Robert Weber	Wed, 3:30-4:30	CSE 678 (except Oct 21st at CSE 110)
Ian Zhu	Thu, 4:30-5:30	CSE 021

a 2-bit ripple-carry adder



Fall 2015

Lecture 5: Canonical forms and predicate logic



Given a truth table:

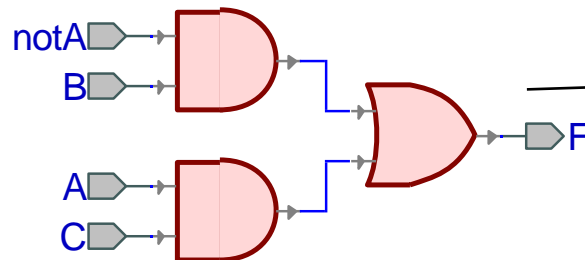
1. Write the Boolean expression
2. Minimize the Boolean expression
3. Draw as gates
4. Map to available gates

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

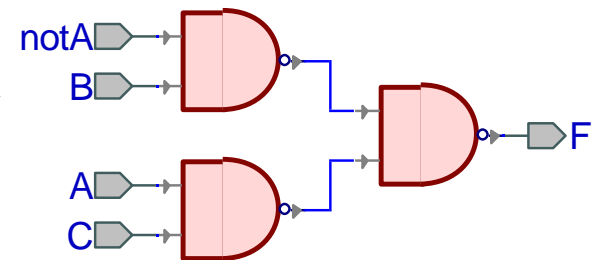
②

$$\begin{aligned}
 F &= A'BC' + A'BC + AB'C + ABC \\
 &= A'B(C' + C) + AC(B' + B) \\
 &= A'B + AC
 \end{aligned}$$

③



④



- Truth table is the unique signature of a Boolean function
- The same truth table can have many gate realizations
 - we've seen this already
 - depends on how good we are at Boolean simplification
- **Canonical forms**
 - standard forms for a Boolean expression
 - we all come up with the same expression

- also known as **Disjunctive Normal Form (DNF)**
- also known as **minterm expansion**

					$F = 001 \quad 011 \quad 101 \quad 110 \quad 111$ $F = A'B'C + A'BC + AB'C + ABC' + ABC$				
A	B	C	F	F'					
0	0	0	0	1					
0	0	1	1	0					
0	1	0	0	1					
0	1	1	1	0					
1	0	0	0	1					
1	0	1	1	0					
1	1	0	1	0					
1	1	1	1	0					

Product term (or minterm)

- ANDed product of literals – input combination for which output is true
- each variable appears exactly once, true or inverted (but not both)

A	B	C	minterms
0	0	0	$A'B'C'$
0	0	1	$A'B'C$
0	1	0	$A'BC'$
0	1	1	$A'BC$
1	0	0	$AB'C'$
1	0	1	$AB'C$
1	1	0	ABC'
1	1	1	ABC

F in canonical form:

$$F(A, B, C) = A'B'C + A'BC + AB'C + ABC' + ABC$$

canonical form \neq minimal form

$$\begin{aligned}
 F(A, B, C) &= A'B'C + A'BC + AB'C + ABC + ABC' \\
 &= (A'B' + A'B + AB' + AB)C + ABC' \\
 &= ((A' + A)(B' + B))C + ABC' \\
 &= C + ABC' \\
 &= ABC' + C \\
 &= AB + C
 \end{aligned}$$

- Also known as **Conjunctive Normal Form (CNF)**
- Also known as **maxterm expansion**

			F =		<i>000</i>	<i>010</i>	<i>100</i>
			F =		(A + B + C)	(A + B' + C)	(A' + B + C)
A	B	C	F	F'			
0	0	0	0	1			
0	0	1	1	0			
0	1	0	0	1			
0	1	1	1	0			
1	0	0	0	1			
1	0	1	1	0			
1	1	0	1	0			
1	1	1	1	0			

Complement of function in sum-of-products form:

$$- F' = A'B'C' + A'BC' + AB'C'$$

Complement again and apply de Morgan's and get the product-of-sums form:

$$- (F')' = (A'B'C' + A'BC' + AB'C')'$$

$$- F = (A + B + C) (A + B' + C) (A' + B + C)$$

Sum term (or maxterm)

- ORed sum of literals – input combination for which output is false
- each variable appears exactly once, true or inverted (but not both)

A	B	C	maxterms
0	0	0	$A+B+C$
0	0	1	$A+B+C'$
0	1	0	$A+B'+C$
0	1	1	$A+B'+C'$
1	0	0	$A'+B+C$
1	0	1	$A'+B+C'$
1	1	0	$A'+B'+C$
1	1	1	$A'+B'+C'$

F in canonical form:

$$F(A, B, C) = (A + B + C) (A + B' + C) (A' + B + C)$$

canonical form \neq minimal form

$$\begin{aligned}
 F(A, B, C) &= (A + B + C) (A + B' + C) (A' + B + C) \\
 &= (A + B + C) (A + B' + C) \\
 &\quad (A + B + C) (A' + B + C) \\
 &= (A + C) (B + C)
 \end{aligned}$$

- **Propositional Logic**

- If Pikachu doesn't wear pants, then he flies on Bieber's jet unless Taylor is feeling lonely.

- **Predicate Logic**

- If x , y , and z are positive integers, then $x^3 + y^3 \neq z^3$.



Predicate or Propositional Function

- A function that returns a truth value, e.g.,

“x is a cat”

“x is prime”

“student x has taken course y”

“ $x > y$ ”

“ $x + y = z$ ” or $\text{Sum}(x, y, z)$

“ $5 < x$ ”

Predicates will have **variables** or **constants** as arguments.

We must specify a “**domain of discourse**”, which is the possible things we’re talking about.

“x is a cat”

(e.g., **mammals**)

“x is prime”

(e.g., **positive whole numbers**)

student x has taken course y”

(e.g., **students and courses**)

$$\forall x P(x)$$

P(x) is true for **every** x in the domain

read as “**for all x, P of x**”

$$\exists x P(x)$$

There is an x in the domain for which P(x) is true

read as “**there exists x, P of x**”

- $\exists x \text{ Even}(x)$
- $\forall x \text{ Odd}(x)$
- $\forall x (\text{Even}(x) \vee \text{Odd}(x))$
- $\exists x (\text{Even}(x) \wedge \text{Odd}(x))$
- $\forall x \text{ Greater}(x+1, x)$
- $\exists x (\text{Even}(x) \wedge \text{Prime}(x))$

Domain:
Positive Integers

Even(x)
Odd(x)
Prime(x)
Greater(x,y)
 (or " $x > y$ ")
Equal(x,y)
 (or " $x = y$ ")
Sum(x,y,z)
 (or " $z = x + y$ ")

- $\forall x \exists y \text{ Greater}(y, x)$

Domain:
Positive Integers

- $\forall x \exists y \text{ Greater}(x, y)$

- $\forall x \exists y (\text{Greater}(y, x) \wedge \text{Prime}(y))$

- $\forall x (\text{Prime}(x) \rightarrow (\text{Equal}(x, 2) \vee \text{Odd}(x)))$

- $\exists x \exists y (\text{Sum}(x, 2, y) \wedge \text{Prime}(x) \wedge \text{Prime}(y))$

Even(x)
Odd(x)
Prime(x)
Greater(x,y)
(or "x>y")
Equal(x,y)
(or "x=y")
Sum(x,y,z)
(or "z=x+y")

	Prev	Now	
• $\forall x \exists y \text{ Greater}(y, x)$	T		Domain: All integers
• $\forall x \exists y \text{ Greater}(x, y)$	F		Even(x) Odd(x) Prime(x) Greater(x,y) (or " $x > y$ ") Equal(x,y) (or " $x = y$ ") Sum(x,y,z) (or " $z = x + y$ ")

Domain of quantifiers is important!

- “Red cats like tofu”

Cat(x)
Red(x)
LikesTofu(x)

- “Some red cats don’t like tofu”

- not every positive integer is prime
- some positive integer is not prime
- prime numbers do not exist
- every positive integer is not prime

$\forall x \text{ PurpleFruit}(x)$

Domain:
Fruit

$\text{PurpleFruit}(x)$

Which one is equal to $\neg \forall x \text{ PurpleFruit}(x)$?

- $\exists x \text{ PurpleFruit}(x)$?
- $\exists x \neg \text{PurpleFruit}(x)$?

$$\neg \forall x \, P(x) \equiv \exists x \, \neg P(x)$$

$$\neg \exists x \, P(x) \equiv \forall x \, \neg P(x)$$

$$\neg \forall x \ P(x) \equiv \exists x \neg P(x)$$

$$\neg \exists x \ P(x) \equiv \forall x \neg P(x)$$

“There is no largest integer.”

$$\begin{aligned} & \neg \exists x \quad \forall y \quad (x \geq y) \\ \equiv & \quad \forall x \neg \forall y \quad (x \geq y) \\ \equiv & \quad \forall x \quad \exists y \neg (x \geq y) \\ \equiv & \quad \forall x \quad \exists y \quad (y > x) \end{aligned}$$

“For every integer there is a larger integer.”

example: $\text{Notlargest}(x) \equiv \exists y \text{ Greater}(y, x)$
 $\equiv \exists z \text{ Greater}(z, x)$

truth value:

doesn't depend on y or z “**bound** variables”

does depend on x “**free** variable”

quantifiers only act on free variables of the formula they quantify

$$\forall x (\exists y (P(x, y) \rightarrow \forall x Q(y, x)))$$

$$\exists x (P(x) \wedge Q(x)) \quad \text{vs.} \quad \exists x P(x) \wedge \exists x Q(x)$$