administrivia

Homework #1 Due Friday at 11:59pm

Please try out Gradescope before then!

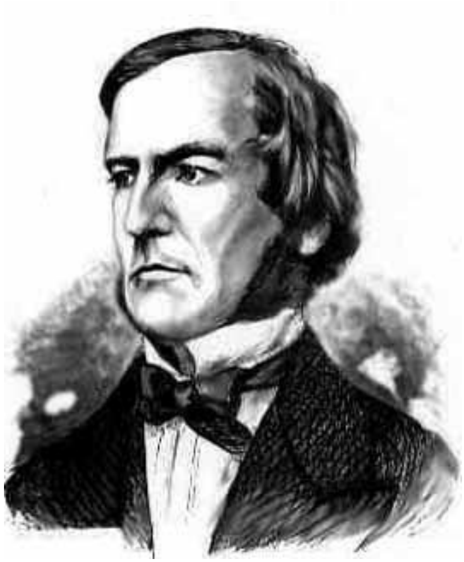
(You can submit multiple times, so do a test run on the first homework.)

Sections start this week:

Section	Day/Time	Room
AA Sam	Th, 830-920	MGH 242
AB Rebecca	Th, 930-1020	MGH 234
AC Robert	Th, 1030-1120	JHN 075
BA Jiechen	Th, 1230-120	MGH 228
BB Tim	Th, 130-220	MGH 242
BC Evan	Th, 230-320	MEB 242

Boolean algebra

- Boolean algebra to circuit design
- Boolean algebra
  - a set of elements B containing {0, 1}
  - binary operations { + , • }
  - and a unary operation { ’ }
  - such that the following axioms hold:



1. The set B contains at least two elements: 0, 1

For any a, b, c in B:

2. closure:	a + b is in B	a • b is in B
3. commutativity:	a + b = b + a	a • b = b • a
4. associativity:	a + (b + c) = (a + b) + c	a • (b • c) = (a • b) • c
5. identity:	a + 0 = a	a • 1 = a
6. distributivity:	a + (b • c) = (a + b) • (a + c)	a • (b + c) = (a • b) + (a • c)
7. complementarity:	a + a’ = 1	a • a’ = 0

axioms and theorems of Boolean algebra

identity:

1. $X + 0 = X$ 1D. $X \cdot 1 = X$

null:

2. $X + 1 = 1$ 2D. $X \cdot 0 = 0$

idempotency:

3. $X + X = X$ 3D. $X \cdot X = X$

involution:

4. $(X')' = X$

complementarity:

5. $X + X' = 1$ 5D. $X \cdot X' = 0$

commutativity:

6. $X + Y = Y + X$ 6D. $X \cdot Y = Y \cdot X$

associativity:

7. $(X + Y) + Z = X + (Y + Z)$ 7D. $(X \cdot Y) \cdot Z = X \cdot (Y \cdot Z)$

distributivity:

8. $X \cdot (Y + Z) = (X \cdot Y) + (X \cdot Z)$ 8D. $X + (Y \cdot Z) = (X + Y) \cdot (X + Z)$

axioms and theorems of Boolean algebra

uniting:

9. $X \cdot Y + X \cdot Y' = X$

9D. $(X + Y) \cdot (X + Y') = X$

absorption:

10. $X + X \cdot Y = X$

10D. $X \cdot (X + Y) = X$

11. $(X + Y') \cdot Y = X \cdot Y$

11D. $(X \cdot Y') + Y = X + Y$

factoring:

12. $(X + Y) \cdot (X' + Z) = X \cdot Z + X' \cdot Y$

12D. $X \cdot Y + X' \cdot Z = (X + Z) \cdot (X' + Y)$

consensus:

13. $(X \cdot Y) + (Y \cdot Z) + (X' \cdot Z) = X \cdot Y + X' \cdot Z$

13D. $(X + Y) \cdot (Y + Z) \cdot (X' + Z) = (X + Y) \cdot (X' + Z)$

de Morgan's:

14. $(X + Y + \ldots)' = X' \cdot Y' \cdot \ldots$

14D. $(X \cdot Y \cdot \ldots)' = X' + Y' + \ldots$

proving theorems (rewriting)

Using the laws of Boolean Algebra:

$X + 1 = 1$

prove the theorem: $X \bullet Y + X \bullet Y' = X$

distributivity (8)
complementarity (5)
identity (1D)

$$
\begin{aligned}
X \bullet Y + X \bullet Y' &= X \bullet (Y + Y') \\
&= X \bullet (1) \\
&= X
\end{aligned}
$$

prove the theorem: $X + X \bullet Y = X$

identity (1D)
distributivity (8)
null (2)
identity (1D)

$$
\begin{aligned}
X + X \bullet Y &= X \bullet 1 + X \bullet Y \\
&= X \bullet (1 + Y) \\
&= X \bullet (1) \\
&= X
\end{aligned}
$$

Using complete truth table:

For example, de Morgan's Law:

$(X + Y)' = X' \bullet Y'$
NOR is equivalent to AND
with inputs complemented

X	Y	X′	Y′	(X + Y)′	X′ • Y′
0	0	1	1		
0	1	1	0		
1	0	0	1	0	0
1	1	0	0		

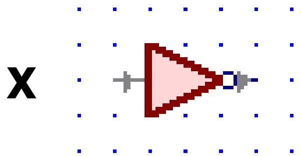
$(X \bullet Y)' = X' + Y'$
NAND is equivalent to OR
with inputs complemented

X	Y	X′	Y′	(X • Y)′	X′ + Y′
0	0	1	1		
0	1	1	0		
1	0	0	1		
1	1	0	0		

more gates

NOT

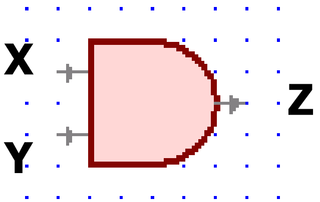
X' $\bar{X}$ $\neg X$



X	Y
0	1
1	0

AND

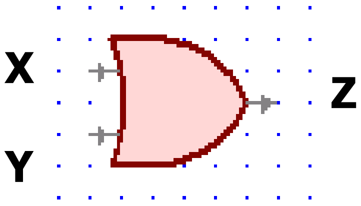
$X \cdot Y$ XY $X \wedge Y$



X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1

OR

$X + Y$ $X \vee Y$

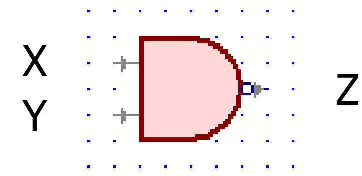


X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1

more gates

NAND

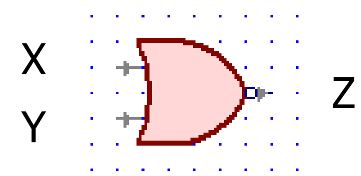
$\neg(X \wedge Y) \quad (XY)'$



X	Y	Z
0	0	1
0	1	1
1	0	1
1	1	0

NOR

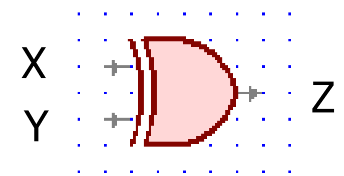
$\neg(X \vee Y) \quad (X + Y)'$



X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	0

XOR

$X \oplus Y \quad X'Y + XY'$

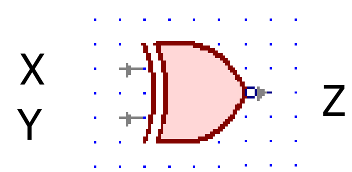


X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	0

XNOR

$X \leftrightarrow Y \quad (X'Y + XY')'$

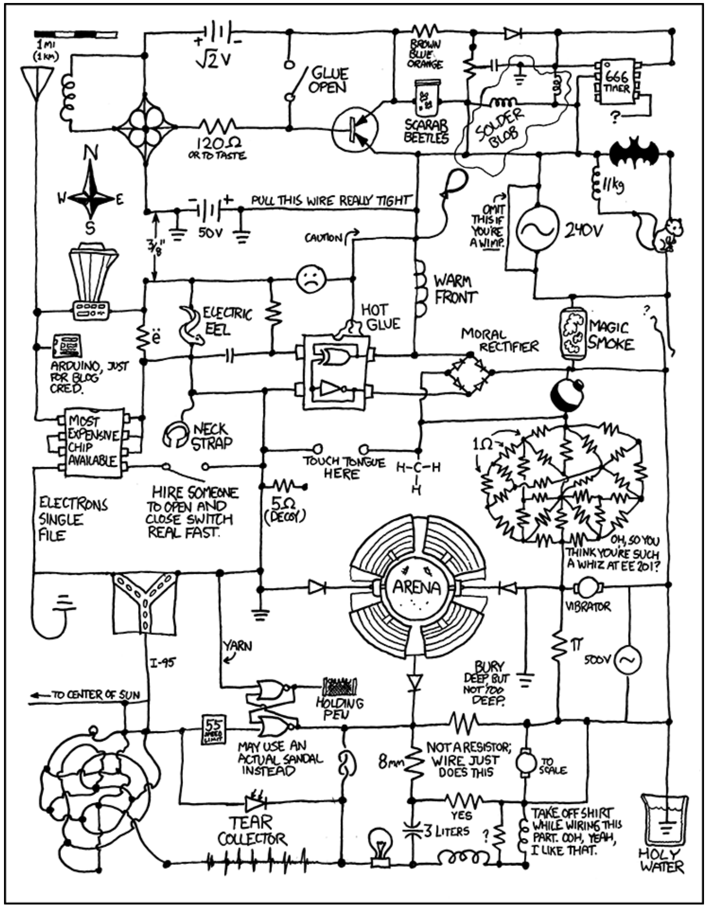
$\equiv XY + X'Y'$



X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	1

Spring 2015

Lecture 4: Boolean Algebra and Circuits



Sessions of class:

We would like to compute the number of lectures or quiz sections remaining *at the start* of a given day of the week.

- Inputs: Day of the Week, Lecture/Section flag
- Output: Number of sessions left

Examples: Input: (Wednesday, Lecture) Output: 2

Input: (Monday, Section) Output: 1

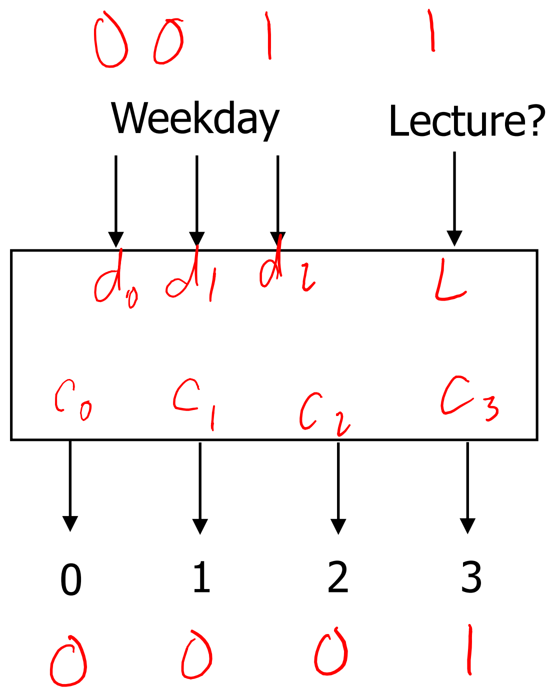
implementation in software

```
public int classesLeft (weekday, lecture_flag) {
    switch (day) {
        case SUNDAY:
        case MONDAY:
            return lecture_flag ? 3 : 1;
        case TUESDAY:
        case WEDNESDAY:
            return lecture_flag ? 2 : 1;
        case THURSDAY:
            return lecture_flag ? 1 : 1;
        case FRIDAY:
            return lecture_flag ? 1 : 0;
        case SATURDAY:
            return lecture_flag ? 0 : 0;
   }
}
```

implementation with combinational logic

Encoding:

- How many bits for each input/output?
- Binary number for weekday
- One bit for each possible output



Sunday – 0
M – 1
T – 2
W – 3
⋮

} output 3

defining our inputs

```
public int classesLeft (weekday, lecture_flag) {
    switch (day) {
        case SUNDAY:
        case MONDAY:
            return lecture_flag ? 3 : 1;
        case TUESDAY:
        case WEDNESDAY:
             return lecture_flag ? 2 : 1;
        case THURSDAY:
             return lecture_flag ? 1 : 1;
        case FRIDAY:
             return lecture_flag ? 1 : 0;
        case SATURDAY:
             return lecture_flag ? 0 : 0;
    }
}
```

Weekday	Number	Binary
Sunday	0	$(000)_2$
Monday	1	$(001)_2$
Tuesday	2	$(010)_2$
Wednesday	3	$(011)_2$
Thursday	4	$(100)_2$
Friday	5	$(101)_2$
Saturday	6	$(110)_2$

converting to a truth table

Weekday	Number	Binary
Sunday	0	$(000)_2$
Monday	1	$(001)_2$
Tuesday	2	$(010)_2$
Wednesday	3	$(011)_2$
Thursday	4	$(100)_2$
Friday	5	$(101)_2$
Saturday	6	$(110)_2$

Weekday	Lecture?	c0	c1	c2	c3
000	0	0	1	0	0
000	1	0	0	0	1
001	0	0	1	0	0
001	1	0	0	0	1
010	0	0	1	0	0
010	1	0	0	1	0
011	0	0	1	0	0
011	1	0	0	1	0
100	-	0	1	0	0
101	0	1	0	0	0
101	1	0	1	0	0
110	-	1	0	0	0
111	-	-	-	-	-

truth table ⇒ logic (part one)

DAY	d2d1d0	L	c0	c1	c2	c3
SunS	000	0	0	1	0	0
SunL	000	1	0	0	0	1
MonS	001	0	0	1	0	0
MonL	001	1	0	0	0	1
TueS	010	0	0	1	0	0
TueL	010	1	0	0	1	0
WedS	011	0	0	1	0	0
WedL	011	1	0	0	1	0
Thu	100	-	0	1	0	0
FriS	101	0	1	0	0	0
FriL	101	1	0	1	0	0
Sat	110	-	1	0	0	0
-	111	-	-	-	-	-

c3 = (DAY == SUN and LEC) or (DAY == MON and LEC)

c3 = (d2 == 0 && d1 == 0 && d0 == 0 && L == 1) ||
(d2 == 0 && d1 == 0 && d0 == 1 && L == 1)

c3 = d2'•d1'•d0'•L + d2'•d1'•d0•L

0'0'0'·1 = 1·1·1·1 = 1

truth table ⇒ logic (part two)

c3 = d2'•d1'•d0'•L + d2'•d1'•d0•L

c2 = (DAY == TUE and LEC) or
(DAY == WED and LEC)

c2 = d2'•d1•d0'•L + d2'•d1•d0•L

DAY	d2d1d0	L	c0	c1	c2	c3
SunS	000	0	0	1	0	0
SunL	000	1	0	0	0	1
MonS	001	0	0	1	0	0
MonL	001	1	0	0	0	1
TueS	010	0	0	1	0	0
TueL	010	1	0	0	1	0
WedS	011	0	0	1	0	0
WedL	011	1	0	0	1	0
Thu	100	-	0	1	0	0
FriS	101	0	1	0	0	0
FriL	101	1	0	1	0	0
Sat	110	-	1	0	0	0
-	111	-	-	-	-	-

truth table ⇒ logic (part three)

DAY	d2d1d0	L	c0	c1	c2	c3
SunS	000	0	0	1	0	0
SunL	000	1	0	0	0	1
MonS	001	0	0	1	0	0
MonL	001	1	0	0	0	1
TueS	010	0	0	1	0	0
TueL	010	1	0	0	1	0
WedS	011	0	0	1	0	0
WedL	011	1	0	0	1	0
Thu	100	-	0	1	0	0
FriS	101	0	1	0	0	0
FriL	101	1	0	1	0	0
Sat	110	-	1	0	0	0
-	111	-	-	-	-	-

c3 = d2’•d1’•d0’•L + d2’•d1’•d0•L

c2 = d2’•d1 •d0’•L + d2’•d1 •d0•L

c1 =

[you do this one]

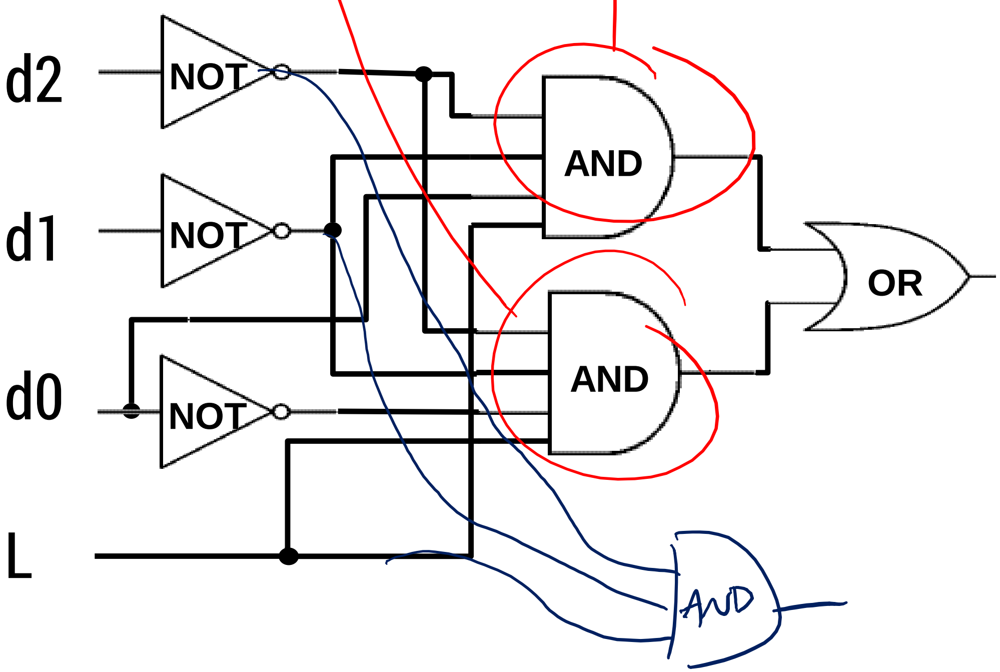
c0 = d2•d1’• d0 •L’ + d2•d1 •d0’

logic ⇒ gates

$c3 = d2' d1' L (d0 + d0')$

$= d2' d1' L$

$c3 = d2' \cdot d1' \cdot d0' \cdot L + d2' \cdot d1' \cdot d0 \cdot L$

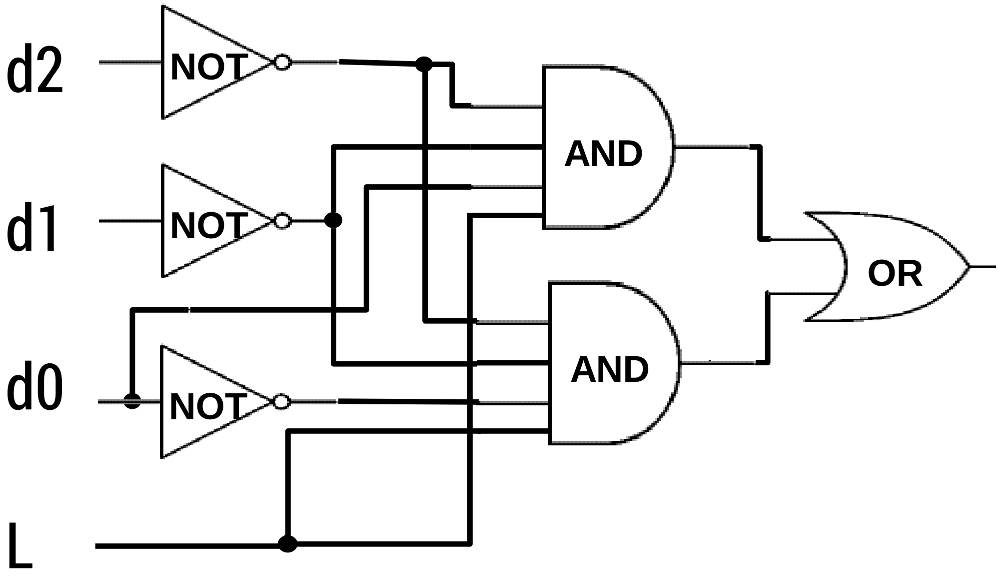


(multiple input AND gates)

DAY	d2d1d0	L	c0	c1	c2	c3
SunS	000	0	0	1	0	0
SunL	000	1	0	0	0	1
MonS	001	0	0	1	0	0
MonL	001	1	0	0	0	1
TueS	010	0	0	1	0	0
TueL	010	1	0	0	1	0
WedS	011	0	0	1	0	0
WedL	011	1	0	0	1	0
Thu	100	-	0	1	0	0
FriS	101	0	1	0	0	0
FriL	101	1	0	1	0	0
Sat	110	-	1	0	0	0
-	111	-	-	-	-	-

simplifying using Boolean algebra

$$c3 = d2' \cdot d1' \cdot d0' \cdot L \; + \; d2' \cdot d1' \cdot d0 \cdot L$$
$$= d2' \cdot d1' \cdot (d0' \; + \; d0) \cdot L$$
$$= d2' \cdot d1' \cdot (1) \cdot L$$
$$= d2' \cdot d1' \cdot L$$



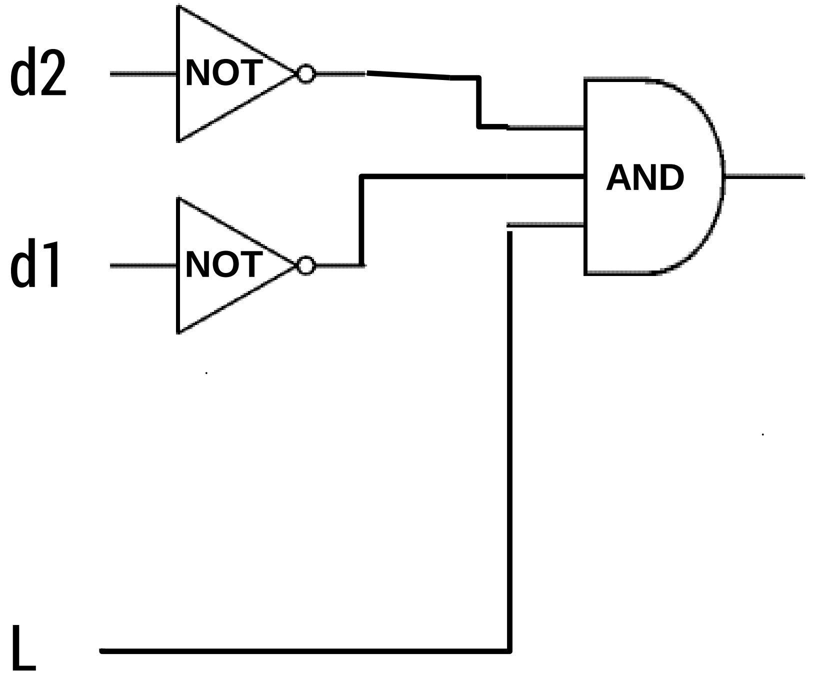
simplifying using Boolean algebra

c3 = d2’•d1’•d0’•L + d2’•d1’•d0•L

= d2’•d1’•(d0’ + d0)•L

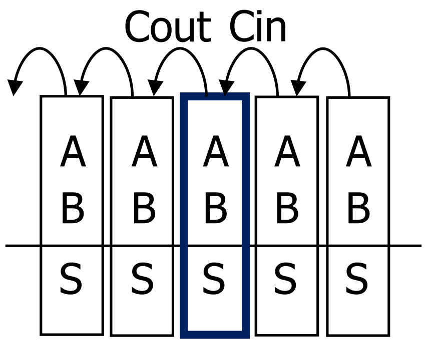
= d2’•d1’•(1)•L

= d2’•d1’•L



1-bit binary adder

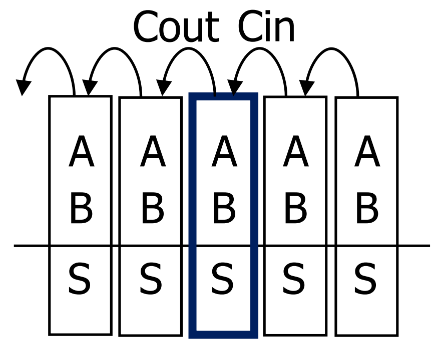
- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



A	B	Cin	Cout	S
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

1-bit binary adder

- **Inputs:** A, B, Carry-in
- **Outputs:** Sum, Carry-out



A	B	Cin	Cout	S	T	T′
0	0	0	0	0	1	0
0	0	1	0	1	1	0
0	1	0	0	1	1	0
0	1	1	1	0	1	0
1	0	0	0	1	1	0
1	0	1	1	0	1	0
1	1	0	1	0	0	1
1	1	1	1	1	1	0

$T' = A B C_{in}'$

$S = A' B' Cin + A' B Cin' + A B' Cin' + A B Cin$

$Cout = A' B Cin + A B' Cin + A B Cin' + A B Cin$

$T = (A B C_{in}')' = A' + B' + C_{in}$

apply theorems to simplify expressions

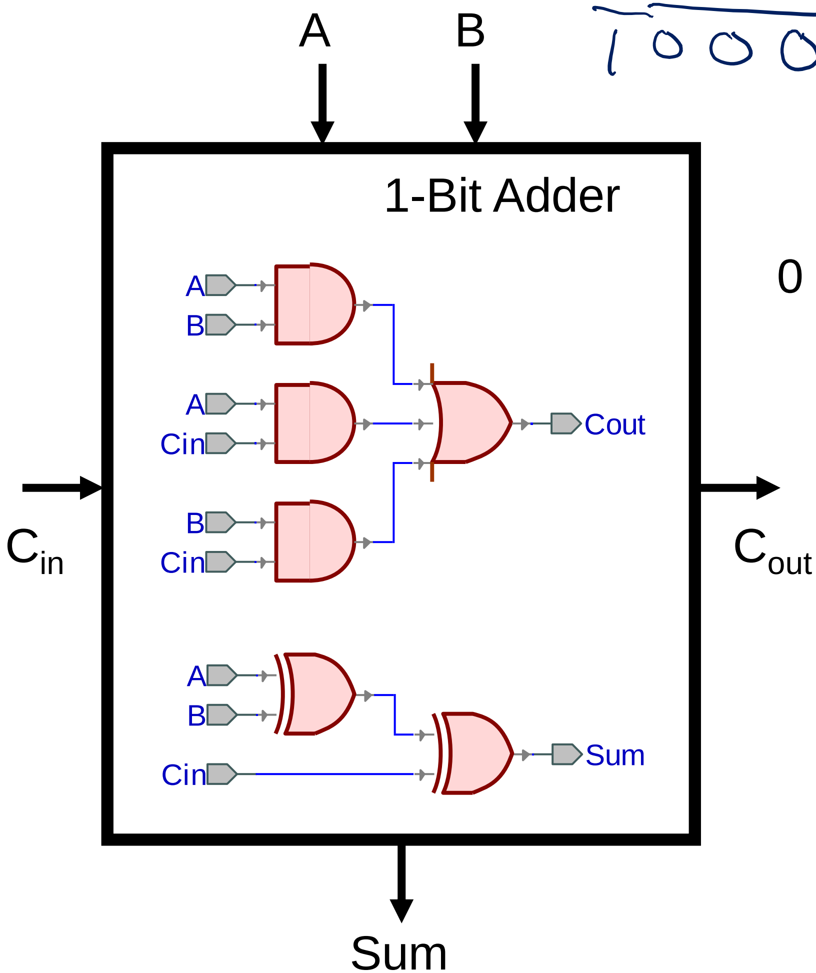
The theorems of Boolean algebra can simplify expressions

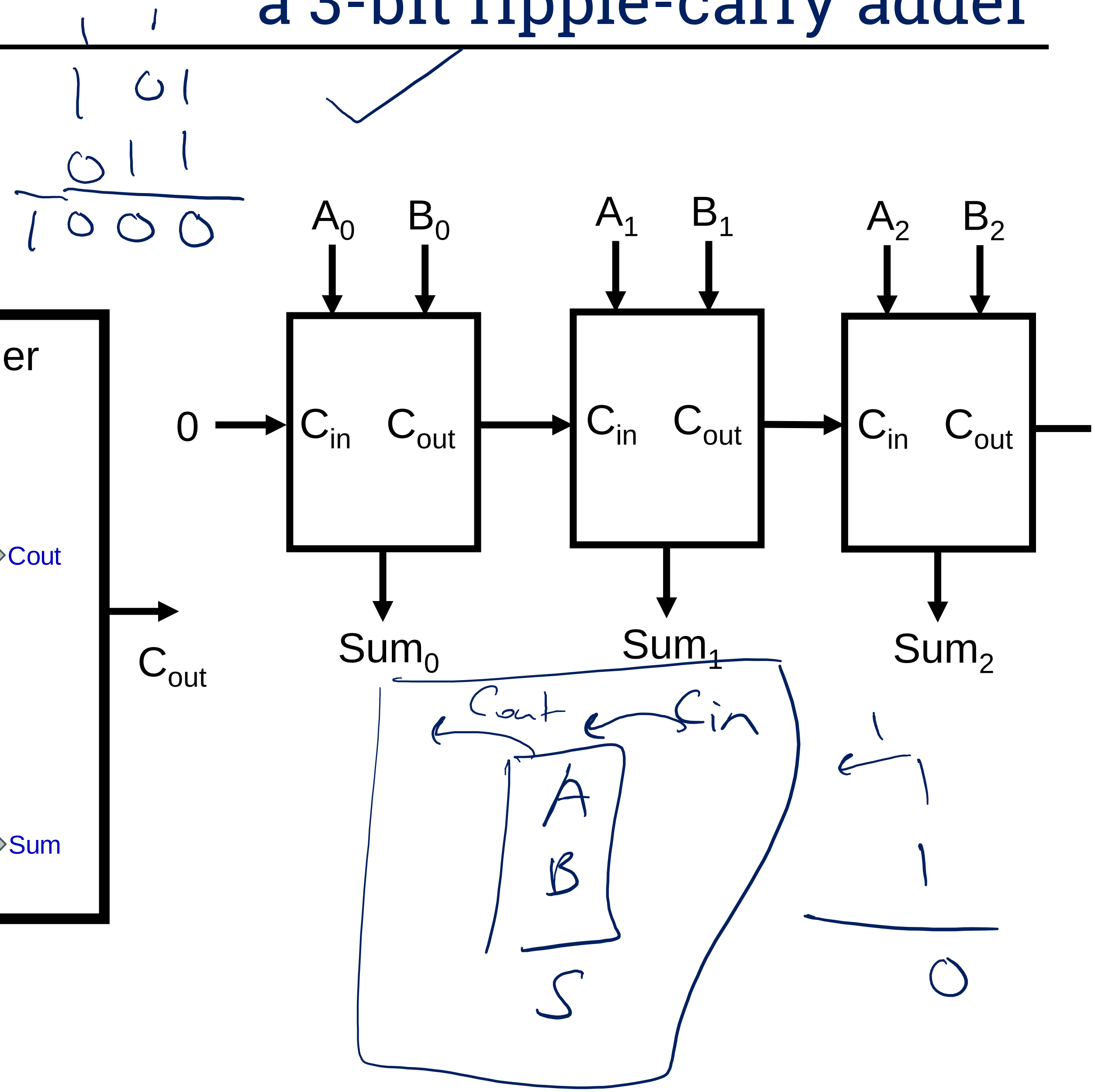
– e.g., full adder's carry-out function

$$
\begin{aligned}
\text{Cout} &= A'\,B\,Cin + A\,B'\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= A'\,B\,Cin + A\,B'\,Cin + A\,B\,Cin' + \boxed{A\,B\,Cin + A\,B\,Cin} \\
&= A'\,B\,Cin + A\,B\,Cin + A\,B'\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= (A' + A)\,B\,Cin + A\,B'\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= (1)\,B\,Cin + A\,B'\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= B\,Cin + A\,B'\,Cin + A\,B\,Cin' + \boxed{A\,B\,Cin + A\,B\,Cin} \\
&= B\,Cin + A\,B'\,Cin + A\,B\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= B\,Cin + A\,(B' + B)\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= B\,Cin + A\,(1)\,Cin + A\,B\,Cin' + A\,B\,Cin \\
&= B\,Cin + A\,Cin + A\,B\,(Cin' + Cin) \\
&= B\,Cin + A\,Cin + A\,B\,(1) \\
&= B\,Cin + A\,Cin + A\,B
\end{aligned}
$$

adding extra terms creates new factoring opportunities

a 3-bit ripple-carry adder





mapping truth tables to logic gates

Given a truth table:

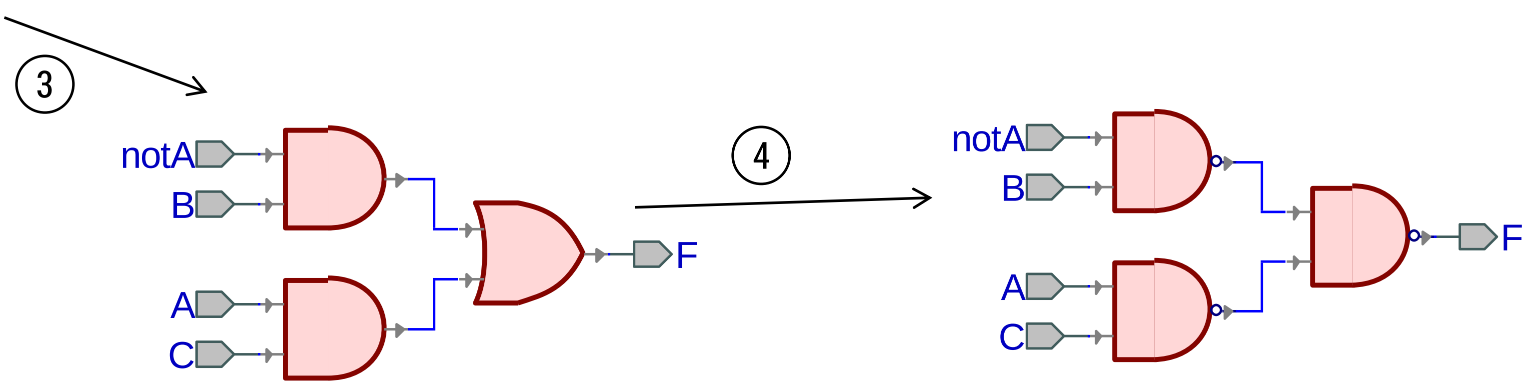
1. Write the Boolean expression
2. Minimize the Boolean expression
3. Draw as gates
4. Map to available gates

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

(1)

(2)

$$F = A'BC' + A'BC + AB'C + ABC$$
$$= A'B(C' + C) + AC(B' + B)$$
$$= A'B + AC$$



canonical forms

- Truth table is the unique signature of a Boolean function
- The same truth table can have many gate realizations
  - we've seen this already
  - depends on how good we are at Boolean simplification
- **Canonical forms**
  - standard forms for a Boolean expression
  - we all come up with the same expression

sum-of-products canonical form

- also known as Disjunctive Normal Form (DNF)
- also known as minterm expansion

F = *001* *011* *101* *110* *111*

$F = A'B'C + A'BC + AB'C + ABC' + ABC$

A	B	C	F	F′
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

sum-of-products canonical form

Product term (or minterm)

- ANDed product of literals – input combination for which output is true
- each variable appears exactly once, true or inverted (but not both)

A	B	C	minterms
0	0	0	A′B′C′
0	0	1	A′B′C
0	1	0	A′BC′
0	1	1	A′BC
1	0	0	AB′C′
1	0	1	AB′C
1	1	0	ABC′
1	1	1	ABC

F in canonical form:

$$F(A, B, C) = A'B'C + A'BC + AB'C + ABC' + ABC$$

canonical form ≠ minimal form

$$\begin{aligned} F(A, B, C) &= A'B'C + A'BC + AB'C + ABC + ABC' \\ &= (A'B' + A'B + AB' + AB)C + ABC' \\ &= ((A' + A)(B' + B))C + ABC' \\ &= C + ABC' \\ &= ABC' + C \\ &= AB + C \end{aligned}$$

product-of-sums canonical form

- Also known as Conjunctive Normal Form (CNF)
- Also known as maxterm expansion

F = *000* *010* *100*

$F = (A + B + C)\,(A + B' + C)\,(A' + B + C)$

A	B	C	F	F′
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

Complement of function in sum-of-products form:

– $F' = A'B'C' + A'BC' + AB'C'$

Complement again and apply de Morgan's and get the product-of-sums form:

– $(F')' = (A'B'C' + A'BC' + AB'C')'$

– $F = (A + B + C)\ (A + B' + C)\ (A' + B + C)$

product-of-sums canonical form

Sum term (or maxterm)

- ORed sum of literals – input combination for which output is false
- each variable appears exactly once, true or inverted (but not both)

A	B	C	maxterms
0	0	0	$A+B+C$
0	0	1	$A+B+C'$
0	1	0	$A+B'+C$
0	1	1	$A+B'+C'$
1	0	0	$A'+B+C$
1	0	1	$A'+B+C'$
1	1	0	$A'+B'+C$
1	1	1	$A'+B'+C'$

F in canonical form:

$$F(A, B, C) = (A + B + C)\ (A + B' + C)\ (A' + B + C)$$

canonical form ≠ minimal form

$$\begin{aligned} F(A, B, C) &= (A + B + C)\ (A + B' + C)\ (A' + B + C) \\ &= (A + B + C)\ (A + B' + C) \\ &\quad\ (A + B + C)\ (A' + B + C) \\ &= (A + C)\ (B + C) \end{aligned}$$