

CSE 311: Foundations of Computing

Fall 2013

Lecture 18: Structural induction, regular expressions



Announcements

Midterm back today

Graded Homework 5 back Friday

Homework 6 out later today

Review: Structural Induction

How to prove $\forall x \in S, P(x)$ is true:

Base Case: Show that $P(u)$ is true for all specific elements u of S mentioned in the *Basis step*

Inductive Hypothesis: Assume that P is true for some arbitrary values of *each* of the existing named elements mentioned in the *Recursive step*

Inductive Step: Prove that $P(w)$ holds for each of the new elements w constructed in the *Recursive step* using the named elements mentioned in the *Inductive Hypothesis*

Conclude that $\forall x \in S, P(x)$

Function Definitions on Recursively Defined Sets

Length:

$$\text{len}(\varepsilon) = 0$$

$$\text{len}(wa) = \text{len}(w) + 1 \text{ for } w \in \Sigma^*, a \in \Sigma$$

Reversal:

$$\varepsilon^R = \varepsilon$$

$$(wa)^R = aw^R \text{ for } w \in \Sigma^*, a \in \Sigma$$

Concatenation:

$$x \bullet \varepsilon = x \text{ for } x \in \Sigma^*$$

$$x \bullet wa = (x \bullet w)a \text{ for } x \in \Sigma^*, a \in \Sigma$$

$\text{len}(x \cdot y) = \text{len}(x) + \text{len}(y)$ for all $x, y \in \Sigma^*$

Let $P(y)$ be " $\text{len}(x \cdot y) = \text{len}(x) + \text{len}(y)$ for all $x \in \Sigma^*$ ".
We prove $P(y)$ for all $y \in \Sigma^*$ by structural induction.

Base Case: $y = \varepsilon$. For any $x \in \Sigma^*$, $\text{len}(x \cdot \varepsilon) = \text{len}(x) = \text{len}(x) + \text{len}(\varepsilon)$
since $\text{len}(\varepsilon) = 0$. Therefore $P(\varepsilon)$ is true

Inductive Hypothesis: Assume that $P(w)$ is true for some arbitrary
 $w \in \Sigma^*$

Inductive Step: Goal: Show that $P(wa)$ is true for every $a \in \Sigma$

Let $a \in \Sigma$. Let $x \in \Sigma^*$. Then $\text{len}(x \cdot wa) = \text{len}((x \cdot w)a)$ by defn of \cdot
 $= \text{len}(x \cdot w) + 1$ by defn of len
 $= \text{len}(x) + \text{len}(w) + 1$ by I.H.
 $= \text{len}(x) + \text{len}(wa)$ by defn of len

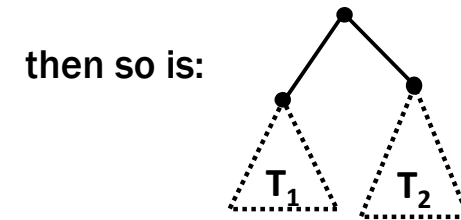
Therefore $\text{len}(x \cdot wa) = \text{len}(x) + \text{len}(wa)$ for all $x \in \Sigma^*$, so $P(wa)$ is true.

So, by induction $\text{len}(x \cdot y) = \text{len}(x) + \text{len}(y)$ for all $x, y \in \Sigma^*$

Rooted Binary Trees

• **Basis:** \bullet is a rooted binary tree

• **Recursive step:** If T_1 and T_2 are rooted binary trees



Functions Defined on Rooted Binary Trees

• $\text{size}(\bullet) = 1$

• $\text{size}(T) = 1 + \text{size}(T_1) + \text{size}(T_2)$

• $\text{height}(\bullet) = 0$

• $\text{height}(T) = 1 + \max\{\text{height}(T_1), \text{height}(T_2)\}$

Claim: For every rooted binary tree T , $\text{size}(T) \leq 2^{\text{height}(T) + 1} - 1$

Languages: sets of strings

- Sets of strings that satisfy special properties are called *languages*. Examples:
 - English sentences
 - Syntactically correct Java/C/C++ programs
 - Σ^* = All strings over alphabet Σ
 - Palindromes over Σ
 - Binary strings that don't have a 0 after a 1
 - Legal variable names. keywords in Java/C/C++
 - Binary strings with an equal # of 0's and 1's

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Regular Expressions

Regular expressions over Σ

- **Basis:**
 - \emptyset, ϵ are regular expressions
 - a is a regular expression for any $a \in \Sigma$
- **Recursive step:**
 - If **A** and **B** are regular expressions then so are:
 - $(A \cup B)$
 - (AB)
 - A^*

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Each Regular Expression is a “pattern”

ϵ matches the **empty string**

a matches the one character string a

$(A \cup B)$ matches all strings that either **A** matches or **B** matches (or both)

(AB) matches all strings that have a first part that **A** matches followed by a second part that **B** matches

A^* matches all strings that have any number of strings (even 0) that **A** matches, one after another

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Examples

- 001^*
- 0^*1^*
- $(0 \cup 1)0(0 \cup 1)0$
- $(0^*1^*)^*$
- $(0 \cup 1)^* 0110 (0 \cup 1)^*$
- $(00 \cup 11)^* (01010 \cup 10001)(0 \cup 1)^*$

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Regular Expressions in Practice

- Used to define the “tokens”: e.g., legal variable names, keywords in programming languages and compilers
- Used in **grep**, a program that does pattern matching searches in UNIX/LINUX
- Pattern matching using regular expressions is an essential feature of PHP
- We can use regular expressions in programs to process strings!

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Regular Expressions in Java

- `Pattern p = Pattern.compile("a*b");`
- `Matcher m = p.matcher("aaaaab");`
- `boolean b = m.matches();`
 - [01] a 0 or a 1 ^ start of string \$ end of string
 - [0-9] any single digit \. period \, comma \- minus
 - . any single character
 - ab a followed by b (AB)
 - (a|b) a or b (A ∪ B)
 - a? zero or one of a (A ∪ ε)
 - a* zero or more of a A*
 - a+ one or more of a AA*
- e.g. `^\[-+]?[0-9]* (\.|\\,)?[0-9]+$`
General form of decimal number e.g. 9.12¹⁴ or -9,8 (Europe)

More Examples

- All binary strings that have an even # of 1's

- All binary strings that *don't* contain 101

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