

# CSE 311: Foundations of Computing

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Fall 2014

## Lecture 18: Structural induction, regular expressions



# Administrivia

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Midterm back today!

Average 80%

Median 81%

Range	Number
<60	
60 – 69	
70 – 79	
80 – 89	
90+	

Graded Homework 5 back Friday

Homework 6 out later today

# Structural Induction

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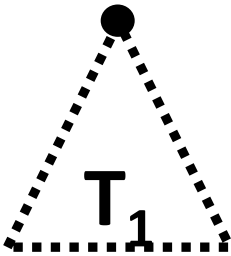
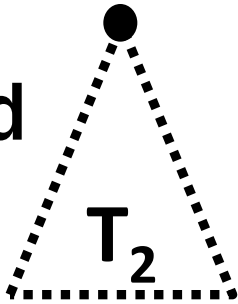
How to prove  $\forall x \in S, P(x)$  is true:

- **Base Case:** Show that  $P(u)$  is true for all specific elements of  $u \in S$  mentioned in the *Basis step*
- **Inductive Hypothesis:** Assume that  $P$  is true for some arbitrary values of each of the existing named elements mentioned in the *Recursive step*
- **Inductive Step:** Prove that  $P(w)$  holds for each of the new elements constructed in the *Recursive step* using the named elements mentioned in the Inductive Hypothesis
- **Conclude** that  $\forall x \in S, P(x)$

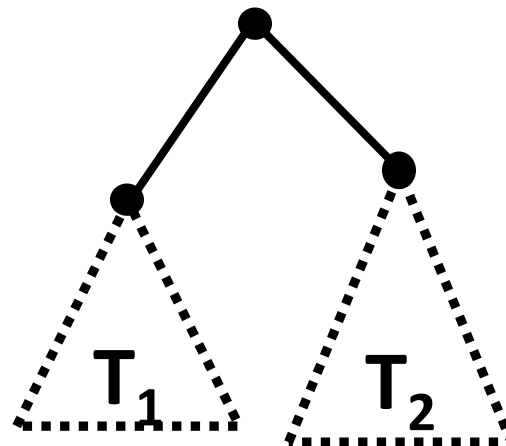
# Rooted Binary Trees

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- **Basis:** • is a rooted binary tree

- **Recursive step:** If  and  are rooted binary trees

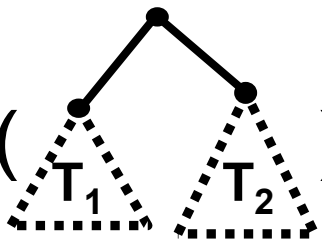
then so is:



# Functions Defined on Rooted Binary Trees

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- $\text{size}(\bullet) = 1$

- $\text{size}(\text{tree}) = 1 + \text{size}(T_1) + \text{size}(T_2)$
- 
- The diagram shows a root node at the top, connected by solid lines to two child nodes. Each child node is the root of a subtree, represented by dashed lines and labeled T1 and T2 respectively.

- $\text{height}(\bullet) = 0$

- $\text{height}(\text{tree}) = 1 + \max\{\text{height}(T_1), \text{height}(T_2)\}$
- 
- The diagram shows a root node at the top, connected by solid lines to two child nodes. Each child node is the root of a subtree, represented by dashed lines and labeled T1 and T2 respectively.

**Claim:** For every rooted binary tree  $T$ ,  $\text{size}(T) \leq 2^{\text{height}(T) + 1} - 1$

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Let  $P(T)$  be “ $\text{size}(T) \leq 2^{\text{height}(T) + 1} - 1$ ”. We go by structural induction on the definition of trees.

**Base Case:** When  $T = \bullet$ ,  $\text{size}(T) = 1$ ,  
 $\text{height}(T) = 0$ ,

Note that  $1 \leq 2^1 - 1 = 2^{0+1} - 1$ .

So,  $P(\bullet)$  is true.

**Induction Hypothesis:** Suppose  $P(T_1)$  and  $P(T_2)$  are true for some  $T_1, T_2$ .

**Induction Step:**

Note that  $\text{size}(\begin{array}{c} \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \vdots \quad \vdots \\ \text{---} \text{---} \\ \text{---} \end{array}) = 1 + \text{size}(T_1) + \text{size}(T_2)$

$$\begin{aligned} &\leq 1 + 2^{\text{height}(T_1) + 1} - 1 + 2^{\text{height}(T_2) + 1} - 1 \quad (\text{by IH}) \\ &\leq 2^{\text{height}(T_1) + 1} + 2^{\text{height}(T_2) + 1} - 1 \\ &\leq 2(2^{\max(\text{height}(T_1), \text{height}(T_2)) + 1}) - 1 \\ &\leq 2(2^{\text{height}(\begin{array}{c} \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \vdots \quad \vdots \\ \text{---} \end{array})}) - 1 \\ &\leq 2^{\text{height}(\begin{array}{c} \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \vdots \quad \vdots \\ \text{---} \end{array}) + 1} - 1 \end{aligned}$$

So, the  $P(T)$  is true for all trees by structural induction.

# Languages: Sets of Strings

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- Sets of strings that satisfy special properties are called *languages*. Examples:
  - English sentences
  - Syntactically correct Java/C/C++ programs
  - $\Sigma^*$  = All strings over alphabet  $\Sigma$
  - Palindromes over  $\Sigma$
  - Binary strings that don't have a 0 after a 1
  - Legal variable names. keywords in Java/C/C++
  - Binary strings with an equal # of 0's and 1's

# Regular Expressions

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## Regular expressions over $\Sigma$

- **Basis:**

  - $\emptyset$ ,  $\epsilon$  are regular expressions

  - $a$  is a regular expression for any  $a \in \Sigma$

- **Recursive step:**

  - If **A** and **B** are regular expressions then so are:

    - **(A  $\cup$  B)**

    - **(AB)**

    - **A\***



# Each Regular Expression is a “pattern”

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$\epsilon$  matches the **empty string**

$a$  matches the one character string  $a$

$(A \cup B)$  matches all strings that either  $A$  matches or  $B$  matches (or both)

$(AB)$  matches all strings that have a first part that  $A$  matches followed by a second part that  $B$  matches

$A^*$  matches all strings that have any number of strings (even 0) that  $A$  matches, one after another

# Examples

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- $001^*$
- $0^*1^*$
- $(0 \cup 1)0(0 \cup 1)0$
- $(0^*1^*)^*$
- $(0 \cup 1)^* 0110 (0 \cup 1)^*$
- $(00 \cup 11)^* (01010 \cup 10001)(0 \cup 1)^*$

# Regular Expressions in Practice

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- Used to define the “tokens”: e.g., legal variable names, keywords in programming languages and compilers
- Used in **grep**, a program that does pattern matching searches in UNIX/LINUX
- Pattern matching using regular expressions is an essential feature of PHP
- We can use regular expressions in programs to process strings!

# Regular Expressions in Java

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- `Pattern p = Pattern.compile("a*b");`
- `Matcher m = p.matcher("aaaaab");`
- `boolean b = m.matches();`

`[01]` a 0 or a 1    `^` start of string    `$` end of string

`[0-9]` any single digit    `\.` period    `\,` comma    `\-` minus

`.` any single character

`ab` a followed by b    **(AB)**

`(a|b)` a or b    **(A ∪ B)**

`a?` zero or one of a    **(A ∪ λ)**

`a*` zero or more of a    **A\***

`a+` one or more of a    **AA\***

- e.g. `^[\-+]?[0-9]*(\.|\,)?[0-9]+$`

General form of decimal number e.g. 9.12<sup>12</sup> or -9,8 (Europe)