CSE 311 Foundations of Computing I

Lecture 22 Finite State Machines: Output and Minimization Spring 2013

Announcements

- Reading assignments

 7th Edition, Sections 13.3 and 13.4
 6th Edition, Section 12.3 and 12.4
- Homework 6 due today
- Homework 7 out

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Last lecture highlights

Finite state machines

- States, transitions, start state, final states
- Languages recognized by FSMs



Last lecture highlights

- Combining FSMs to check two properties at once
 - New states record states of both FSMs



State machines with output

	Inp	Output	
State	L	R	
S ₁	S ₁	s ₂	Веер
s ₂	s ₁	s ₃	
s ₃	s ₂	S ₄	
S ₄	S ₃	S ₅	
s ₅	s ₄	s ₅	Buzzer

"Tug-of-war"

(S₃ ,

 S_4

 S_5

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State Machines with Output



SWIFTCERS Vending Machine



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Enter 15 cents in dimes or nickels Press S or B for a candy bar



Vending Machine, Partial Version 1

(S₂



 $\left(\mathbf{S}_{1} \right)$

Vending Machine, Partial Version 2





Adding additional "unexpected" transitions

State Minimization

- Many different FSMs (DFAs) for the same problem
- Take a given FSM and try to reduce its state set by combining states
 - Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this

State minimization algorithm

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- 1. Put states into groups based on their outputs (or whether they are final states or not)
- 2. Repeat the following until no change happens
 - a. If there is a symbol s so that not all states in a group
 G agree on which group s leads to, split G into
 smaller groups based on which group the states go





State Minimization Example



present state	0	ne: 1	xt sta 2	ate 3	output
S0 S1 S2 S3 S4 S5	S0 S0 S1 S1 S0 S1	S1 S3 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S4 S5 S5 S5 S5	1 0 1 0 1 0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)



S0 [1] [0] 2 S2 [1] S3 [0] S4 [1] S5 [0]

State Minimization Example

present state	0	ne» 1	kt sta 2	ate 3	output		
S0 S1 S2 S3 S4 S5	S0 S0 S1 S1 S0 S1	S1 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5	1 0 1 0 1 0		
state transition table							
Put states into groups based on their outputs (or whether they are final states or not)							

State Minimization Example



present state	0	ne: 1	xt sta 2	ate 3	output
S0 S1 S2 S34 S5	S0 S0 S1 S1 S0 S1	S1 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5	1 0 1 0 1 0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group G agree on which group s leads to, split G based on which group the states go to on s

State Minimization Example



			-		
present state	0	nex 1	kt sta 2	ate 3	output
S0 S1 S2 S3 S4 S5	S0 S1 S1 S0 S1	S1 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5 S5	1 0 1 0 1 0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group G agree on which group s leads to, split G based on which group the states go to on s

State Minimization Example



present state	0	ne: 1	xt sta 2	ate 3	output
S0 S1 S2 S3 S4 S5	S0 S0 S1 S1 S0 S1	S1 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5 S5	1 0 1 0 1 0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol **s** so that not all states in a group G agree on which group s leads to, split G based on which group the states go to on s

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State Minimization Example



	present	next state			output			
	state	0	1	2	3			
	S0 S1 S2 S3 S4 S5	S0 S0 S1 S1 S0 S1	S1 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5 S5	1 0 1 0 1 0		
	state transition table							
	Put states into groups based on their outputs (or whether they are final states or not)							
If there is a symbol s so that not all states in								

a group G agree on which group s leads to, split G based on which group the states go to on s

State Minimization Example



present state	0	ne: 1	xt sta 2	ate 3	output
S0 S1 S2 S3 S4 S5	S0 S0 S1 S1 S0 S1	S1 S3 S3 S0 S1 S4	S2 S1 S2 S4 S2 S0	S3 S5 S4 S5 S5 S5 S5	1 0 1 0 1 0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol s so that not all states in a group G agree on which group s leads to, split G based on which group the states go to on s

State Minimization Example



			•				
present state	0	ne: 1	xt sta 2	ate 3	output		
S0 S1 S2 S3 S4 S5	S0 S1 S1 S0 S1 S1	51 53 53 50 51 54	S2 S1 S2 S4 S2 S0	S3 S5 S5 S5 S5 S5 S5	1 0 1 0 1 0		
state transition table							
Can combine states S0-S4 and S3-S5.							

In table replace all S4 with S0 and all S5 with S3

Minimized Machine





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Another way to look at DFAs

Definition: The label of a path in a DFA is the concatenation of all the labels on its edges in order

Lemma: x is in the language recognized by a DFA iff x labels a path from the start state to some final state



Nondeterministic Finite Automaton (NFA)

- Graph with start state, final states, edges labeled by symbols (like DFA) but
 - Not required to have exactly 1 edge out of each state labeled by each symbol - can have 0 or >1
 - Also can have edges labeled by empty string λ
- Definition: x is in the language recognized by an NFA iff x labels a path from the start state to some final state



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Design an NFA to recognize the set of binary strings that contain 111 or have an even # of 1's