## Announcements

## CSE 311 Foundations of Computing I

## Lecture 22

Finite State Machines: Output and
Minimization
Spring 2013

- Reading assignments
$-7^{\text {th }}$ Edition, Sections 13.3 and 13.4
-6 ${ }^{\text {th }}$ Edition, Section 12.3 and 12.4
- Homework 6 due today
- Homework 7 out


## Last lecture highlights

Finite state machines

- States, transitions, start state, final states
- Languages recognized by FSMs



## Last lecture highlights

- Combining FSMs to check two properties at once
- New states record states of both FSMs


State machines with output

|  | Input |  | Output |
| :---: | :---: | :---: | :---: |
| State | L | R |  |
| $\mathrm{s}_{1}$ | $\mathrm{~s}_{1}$ | $\mathrm{~s}_{2}$ | Beep |
| $\mathrm{s}_{2}$ | $\mathrm{~s}_{1}$ | $\mathrm{~s}_{3}$ |  |
| $\mathrm{~s}_{3}$ | $\mathrm{~s}_{2}$ | $\mathrm{~s}_{4}$ |  |
| $\mathrm{~s}_{4}$ | $\mathrm{~s}_{3}$ | $\mathrm{~s}_{5}$ |  |
| $\mathrm{~s}_{5}$ | $\mathrm{~s}_{4}$ | $\mathrm{~s}_{5}$ | Buzzer |

"Tug-of-war"


Vending Machine, Partial Version 1


B, S

## State Machines with Output

## [SNIIFKERSII Vending Machine

Enter 15 cents in dimes or nickels
Press S or B for a candy bar


Vending Machine, Partial Version 2



Adding additional "unexpected" transitions

## State Minimization

- Many different FSMs (DFAs) for the same problem
- Take a given FSM and try to reduce its state set by combining states
- Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this


## State minimization algorithm

1. Put states into groups based on their outputs (or whether they are final states or not)
2. Repeat the following until no change happens
a. If there is a symbol s so that not all states in a group G agree on which group $s$ leads to, split G into smaller groups based on which group the states go to on $s$


## State Minimization Example



| present |  |  |  |  |  |
| :---: | :--- | :--- | :--- | :--- | :--- |
| state | 0 | 1 | 2 | 3 | output |
| S0 | S0 | S1 | S2 | S3 | 1 |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | S0 | S1 | S2 | S5 | 1 |
| S5 | S1 | S4 | S0 | S5 | 0 |
| state |  |  |  |  |  |
| transition table |  |  |  |  |  |

Put states into groups based on their outputs (or whether they are final states or not)

## State Minimization Example



| presen |  |  | xt st |  | outpu |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state | S0 | 1 | 2 | S3 |  |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | SO | S4 | S5 | 0 |
| S4 | S0 | S1 | S2 | S5 | 1 |
| S5 | S1 | S4 | SO | S5 | 0 |

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

## State Minimization Example



| present | next state |  |  |  | output |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state |  |  |  |  |  |
| SO | S0 | 51 | S2 | S3 | 1 |
| S1 | SO | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | SO | S1 | S2 | S5 | 1 |
| S5 | S1 |  | S0 | S5 | 0 |

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol $\boldsymbol{s}$ so that not all states in a group $G$ agree on which group $\boldsymbol{s}$ leads to, split $G$ based on which group the states go to on $\boldsymbol{s}$

State Minimization Example


| present | 0 next state |  |  |  | outpu |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state |  |  |  |  |  |
| S0 | SO | 51 | S2 | 53 | 1 |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | S0 | S1 | S2 | S5 | 1 |
| S5 | S1 | S4 | S0 | S5 | 0 |

state
transition table
Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol $\boldsymbol{s}$ so that not all states in a group $G$ agree on which group $\boldsymbol{s}$ leads to, split $G$ based on which group the states go to on $\boldsymbol{s}$

State Minimization Example


| present | next state |  |  |  | output |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state | 0 | 1 |  |  |  |
| SO | S0 | 51 | S2 | S3 | 1 |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | SO | S1 | S2 | S5 |  |
| S5 | S1 | S4 | S0 |  | 0 |

state
transition table
Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol $\boldsymbol{s}$ so that not all states in a group $G$ agree on which group $s$ leads to, split G based on which group the states go to on $\boldsymbol{s}$

## State Minimization Example



| present |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| ptate | 0 | next state |  |  | output |
| state | S. | 3 |  |  |  |
| S0 | S0 | S1 | S2 | S3 | 1 |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | S0 | S1 | S2 | S5 | 1 |
| S5 | S1 | S4 | S0 | S5 | 0 |

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol $\boldsymbol{s}$ so that not all states in a group $G$ agree on which group $\boldsymbol{s}$ leads to, split $G$ based on which group the states go to on $\boldsymbol{s}$

## State Minimization Example



| present | next state |  |  |  | output |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state |  | 1 |  |  |  |
| S0 | SO | S1 | S2 | S3 | 1 |
| S1 | S0 | S3 | S1 | S5 | 0 |
| S2 | S1 | S3 | S2 | S4 | 1 |
| S3 | S1 | S0 | S4 | S5 | 0 |
| S4 | S0 | S1 | S2 | S5 | 1 |
| S5 | S1 | S4 | SO | S5 | 0 |

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol $\boldsymbol{s}$ so that not all states in a group $G$ agree on which group $\boldsymbol{s}$ leads to, split $G$ based on which group the states go to on $\boldsymbol{s}$

State Minimization Example


transition table

Can combine states S0-S4 and S3-S5.

In table replace all S4 with S0 and all S5 with S3

## Minimized Machine



| present | next state |  |  |  | output |
| :---: | :---: | :---: | :---: | :---: | :---: |
| state | 0 | 1 |  |  |  |
| S0 | S0 | S1 | S2 | S3 | 1 |
| S1 | S0 | S3 | S1 | S3 | 0 |
| S2 | S1 | S3 | S2 | S0 | 1 |
| S3 |  | S0 | S0 |  | 0 |

state
transition table

## Another way to look at DFAs

Definition: The label of a path in a DFA is the concatenation of all the labels on its edges in order

Lemma: $x$ is in the language recognized by a DFA iff $x$ labels a path from the start state to some final state

## Nondeterministic Finite Automaton (NFA)

- Graph with start state, final states, edges labeled by symbols (like DFA) but
- Not required to have exactly 1 edge out of each state labeled by each symbol - can have 0 or >1
- Also can have edges labeled by empty string $\lambda$
- Definition: $x$ is in the language recognized by an NFA iff $x$ labels a path from the start state to some final state


Design an NFA to recognize the set of binary strings that contain 111 or have an even \# of 1's

