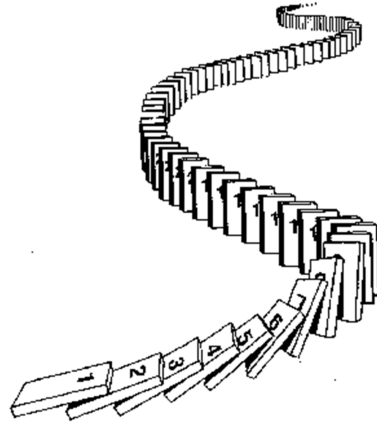


CSE 311: Foundations of Computing

Fall 2013

Lecture 14: Mathematical Induction



announcements

Reading assignment

Induction

5.1-5.2, 7th edition

4.1-4.2, 6th edition

mathematical induction

Method for proving statements about all integers

$n \geq 0$.

- Part of sound logical inference that applies only in the domain of integers

Not like scientific induction which is more like a guess from examples

- Particularly useful for reasoning about programs since the statement might be “after n times through this loop, property $P(n)$ holds”

finding a pattern

- $2^0 - 1 = 1 - 1 = 0 = 3 \cdot 0$
- $2^2 - 1 = 4 - 1 = 3 = 3 \cdot 1$
- $2^4 - 1 = 16 - 1 = 15 = 3 \cdot 5$
- $2^6 - 1 = 64 - 1 = 63 = 3 \cdot 21$
- $2^8 - 1 = 256 - 1 = 255 = 3 \cdot 85$
- ...

how do you prove it?

Want to prove $3 \mid 2^{2n} - 1$ for all integers $n \geq 0$

- $n = 0$
- $n = 1$
- $n = 2$
- $n = 3$
- ...

induction as a rule of Inference

Domain: integers ≥ 0

$$\begin{array}{l} P(0) \\ \forall k (P(k) \rightarrow P(k+1)) \\ \hline \therefore \forall n P(n) \end{array}$$

using the induction rule in a formal proof

$$\begin{array}{l} P(0) \\ \forall k (P(k) \rightarrow P(k+1)) \\ \hline \therefore \forall n P(n) \end{array}$$

1. Prove $P(0)$
2. Let k be an arbitrary integer ≥ 0
 3. Assume that $P(k)$ is true
 4. ...
 5. Prove $P(k+1)$ is true
6. $P(k) \rightarrow P(k+1)$ Direct Proof Rule
7. $\forall k (P(k) \rightarrow P(k+1))$ Intro \forall from 2-6
8. $\forall n P(n)$ Induction Rule 1&7

using the induction rule in a formal proof

$$\begin{array}{l} P(0) \\ \forall k (P(k) \rightarrow P(k+1)) \\ \hline \therefore \forall n P(n) \end{array}$$

1. Prove $P(0)$ **Base Case**
2. Let k be an arbitrary integer ≥ 0 **Inductive Hypothesis**
 3. Assume that $P(k)$ is true
 4. ...
 5. Prove $P(k+1)$ is true **Inductive Step**
6. $P(k) \rightarrow P(k+1)$ Direct Proof Rule
7. $\forall k (P(k) \rightarrow P(k+1))$ Intro \forall from 2-6
8. $\forall n P(n)$ Induction Rule 1&7

Conclusion

5 steps to inductive proofs in english

Proof:

1. "By induction we will show that $P(n)$ is true for every $n \geq 0$."
2. "Base Case:" Prove $P(0)$
3. "Inductive Hypothesis:"
Assume $P(k)$ is true for some arbitrary integer $k \geq 0$
4. "Inductive Step:" Want to prove that $P(k+1)$ is true:
Use the goal to figure out what you need.
Make sure you are using I.H. and point out where you are using it. (Don't assume $P(k+1)$!!)
5. "Conclusion: Result follows by induction"

induction example

Want to prove $3 \mid 2^{2n} - 1$ for all $n \geq 0$.

$$3 \mid 2^{2n} - 1 \text{ for all } n \geq 0.$$

geometric sum

$$1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1 \text{ for all } n \geq 0$$

$$1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1 \text{ for all } n \geq 0$$

sum of first n numbers

$$\text{For all } n \geq 1: 1 + 2 + \dots + n = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

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harmonic numbers

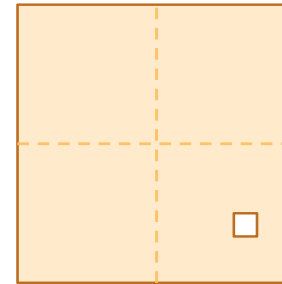
$$H_m = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{m} = \sum_{i=1}^m \frac{1}{i}$$

Prove that $H_{2^n} \geq 1 + \frac{n}{2}$ for all $n \geq 1$.

Prove that $H_{2^n} \geq 1 + \frac{n}{2}$ for all $n \geq 1$.

checkerboard tiling

Prove that a $2^n \times 2^n$ checkerboard with one square removed can be tiled with:



strong induction

$$P(0)$$
$$\forall k \left((P(0) \wedge P(1) \wedge P(2) \wedge \dots \wedge P(k)) \rightarrow P(k+1) \right)$$

$$\therefore \forall n P(n)$$

Follows from ordinary induction applied to

$$Q(n) = P(0) \wedge P(1) \wedge P(2) \wedge \dots \wedge P(n)$$

strong induction english proofs

1. By induction we will show that $P(n)$ is true for every $n \geq 0$
2. Base Case: Prove $P(0)$
3. Inductive Hypothesis:
Assume that for some arbitrary integer $k \geq 0$, $P(j)$ is true for every j from 0 to k
4. Inductive Step:
Prove that $P(k+1)$ is true using the Inductive Hypothesis (that $P(j)$ is true for all values $\leq k$)
5. Conclusion: Result follows by induction

every integer ≥ 2 is the product of primes

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recursive definitions of functions

- $F(0) = 0$; $F(n + 1) = F(n) + 1$ for all $n \geq 0$
- $G(0) = 1$; $G(n + 1) = 2 \cdot G(n)$ for all $n \geq 0$
- $0! = 1$; $(n + 1)! = (n + 1) \cdot n!$ for all $n \geq 0$
- $H(0) = 1$; $H(n + 1) = 2^{H(n)}$ for all $n \geq 0$

fibonacci numbers

$$\begin{aligned}f_0 &= 0 \\f_1 &= 1 \\f_n &= f_{n-1} + f_{n-2} \text{ for all } n \geq 2\end{aligned}$$

bounding the fibonacci numbers

Theorem: $2^{\frac{n}{2}-1} \leq f_n < 2^n$ for all $n \geq 2$.

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