

# CSE 311 Foundations of Computing I

Lecture 22  
Finite State Machines: Output and Minimization  
Autumn 2011

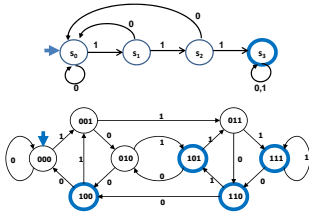
## Announcements

- Reading assignments
  - 7<sup>th</sup> Edition, Sections 13.3 and 13.4
  - 6<sup>th</sup> Edition, Section 12.3 and 12.4
  - 5<sup>th</sup> Edition, Section 11.3 and 11.4

## Last lecture highlights

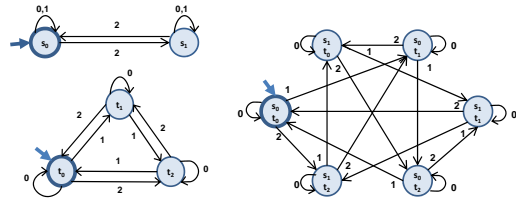
Finite state machines

- States, transitions, start state, final states
- Languages recognized by FSMs



## Last lecture highlights

- Combining FSMs to check two properties at once
  - New states record states of both FSMs



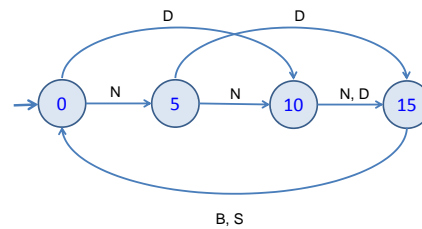
## State Machines with Output



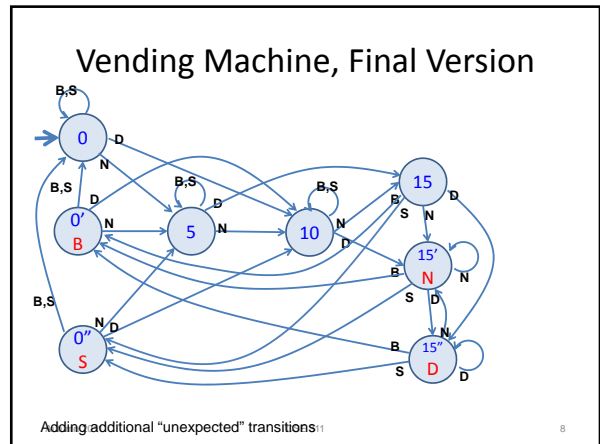
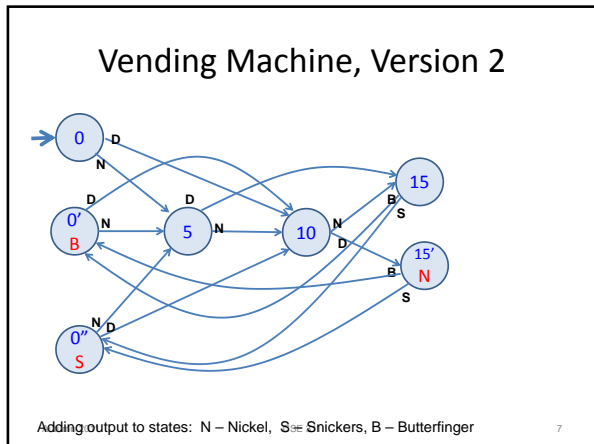
Enter 15 cents in dimes or nickels  
Press S or B for a candy bar



## Vending Machine, Version 1



Basic transitions on N (nickel), D (dime), B (butterfinger), S (snickers)



### State Minimization

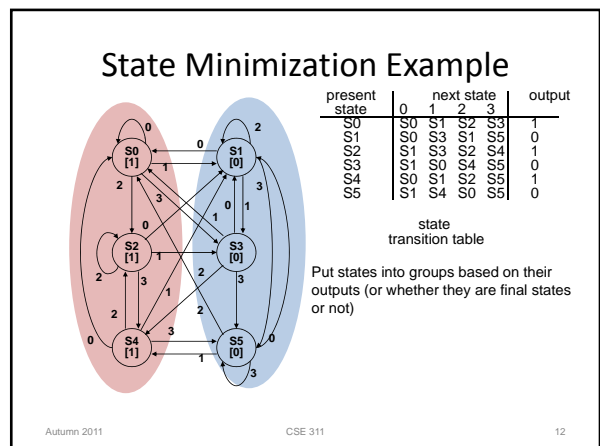
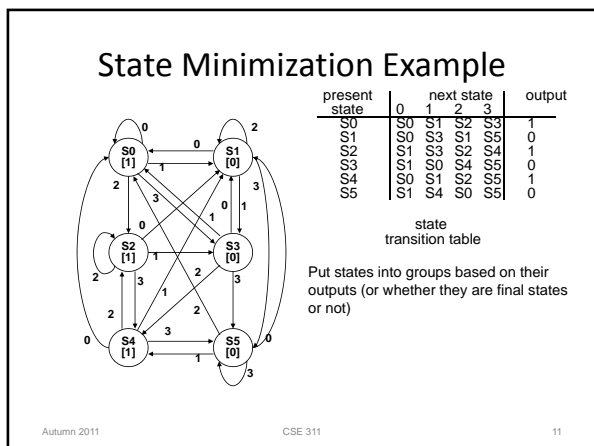
- Many different FSMs (DFAs) for the same problem
- Take a given FSM and try to reduce its state set by combining states
  - Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this

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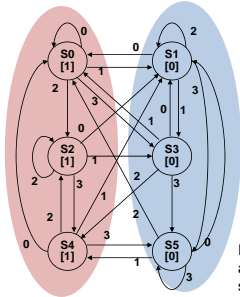
### State minimization algorithm

1. Put states into groups based on their outputs (or whether they are final states or not)
2. Repeat the following until no change happens
  - a. If there is a symbol  $s$  so that not all states in a group  $G$  agree on which group  $s$  leads to, split  $G$  into smaller groups based on which group the states go to on  $s$

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## State Minimization Example



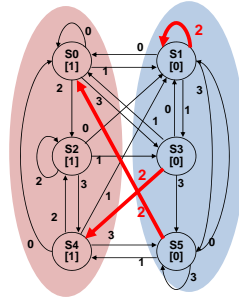
present state	0	1	2	3	output
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they are final states or not)

If there is a symbol  $s$  so that not all states in a group  $G$  agree on which group  $s$  leads to, split  $G$  based on which group the states go to on  $s$

## State Minimization Example



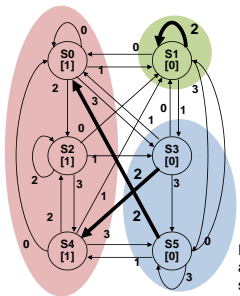
present state	0	1	2	3	output
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

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## State Minimization Example



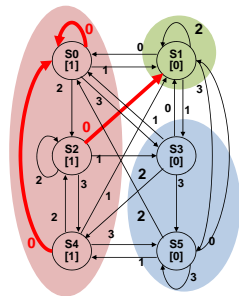
present state	0	1	2	3	output
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
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S5	S1	S4	S0	S5	0

state transition table

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## State Minimization Example



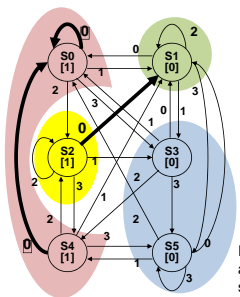
present state	0	1	2	3	output
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S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
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state transition table

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If there is a symbol  $s$  so that not all states in a group  $G$  agree on which group  $s$  leads to, split  $G$  based on which group the states go to on  $s$

## State Minimization Example



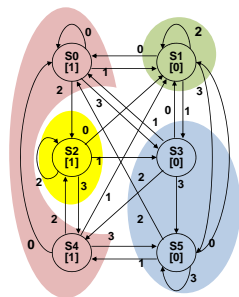
present state	0	1	2	3	output
S0	S0	S1	S2	S3	1
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state transition table

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## State Minimization Example



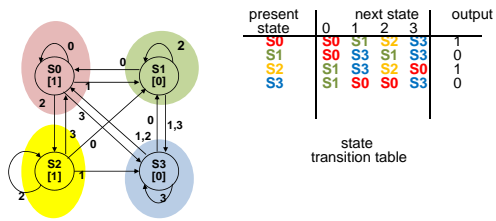
present state	0	1	2	3	output
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Can combine states S0-S4 and S3-S5.

In table replace all S4 with S0 and all S5 with S3

## Minimized Machine



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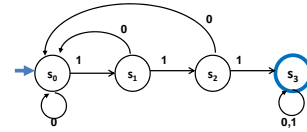
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## Another way to look at DFAs

Definition: The label of a path in a DFA is the concatenation of all the labels on its edges in order

Lemma:  $x$  is in the language recognized by a DFA iff  $x$  labels a path from the start state to some final state



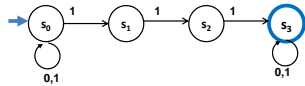
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## Nondeterministic Finite Automaton (NFA)

- Graph with start state, final states, edges labeled by symbols (like DFA) but
  - Not required to have exactly 1 edge out of each state labeled by each symbol - can have 0 or >1
  - Also can have edges labeled by empty string  $\lambda$
- Definition:  $x$  is in the language recognized by an NFA iff  $x$  labels a path from the start state to some final state



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Design an NFA to recognize the set of binary strings that contain 111 or have an even # of 1's

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