

CSE 311 Foundations of Computing I

Lecture 20
Relations, Graphs, Finite State Machines
Autumn 2011

Announcements

- Reading assignments
 - 7th Edition, Sections 9.3 and 13.3
 - 6th Edition, Section 8.3 and 12.3
 - 5th Edition, Section 7.3 and 11.3

Lecture highlights



Let A and B be sets,
A **binary relation from A to B** is a subset of $A \times B$

$S \circ R = \{(a, c) \mid \exists b \text{ such that } (a,b) \in R \text{ and } (b,c) \in S\}$

$R^1 = R; \quad R^{n+1} = R^n \circ R$

n-ary relations

Let A_1, A_2, \dots, A_n be sets. An n-ary relation on these sets is a subset of $A_1 \times A_2 \times \dots \times A_n$.

Relational databases

Student_Name	ID_Number	Major	GPA
Knuth	328012098	CS	4.00
Von Neuman	481080220	CS	3.78
Von Neuman	481080220	Mathematics	3.78
Russell	238082388	Philosophy	3.85
Einstein	238001920	Physics	2.11
Newton	1727017	Mathematics	3.61
Karp	348882811	CS	3.98
Newton	1727017	Physics	3.61
Bernoulli	2921938	Mathematics	3.21
Bernoulli	2921939	Mathematics	3.54

Alternate Approach

Student_ID	Name	GPA	Student_ID	Major
328012098	Knuth	4.00	328012098	CS
481080220	Von Neuman	3.78	481080220	CS
238082388	Russell	3.85	481080220	Mathematics
238001920	Einstein	2.11	238082388	Philosophy
1727017	Newton	3.61	238001920	Physics
348882811	Karp	3.98	1727017	Mathematics
2921938	Bernoulli	3.21	348882811	CS
2921939	Bernoulli	3.54	1727017	Physics
			2921938	Mathematics
			2921939	Mathematics

Database Operations

Projection

Join

Select

Matrix representation

Relation R on $A=\{a_1, \dots, a_p\}$

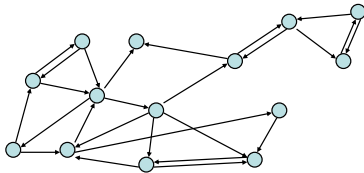
$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, a_j) \in R, \\ 0 & \text{if } (a_i, a_j) \notin R. \end{cases}$$

$\{(1, 1), (1, 2), (1, 4), (2, 1), (2, 3), (3, 2), (3, 3), (4, 2), (4, 3)\}$

Directed graphs

$G = (V, E)$
V – vertices
E – edges, order pairs of vertices

Path: v_1, v_2, \dots, v_k , with (v_i, v_{i+1}) in E
Simple Path
Cycle
Simple Cycle



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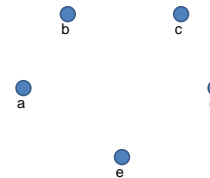
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Representation of relations

Directed Graph Representation (Digraph)

$\{(a, b), (a, a), (b, a), (c, a), (c, d), (c, e), (d, e)\}$



Paths in relations

Let R be a relation on a set A. There is a path of length n from a to b if and only if $(a, b) \in R^n$

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Connectivity relation

Let R be a relation on a set A. The connectivity relation R^* consists of the pairs (a,b) such that there is a path from a to b in R.

$$R^* = \bigcup_{k=0}^{\infty} R^k$$

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Properties of Relations

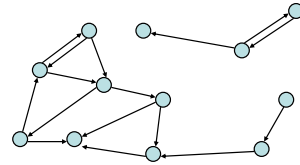
Let R be a relation on A

R is **reflexive** iff $(a,a) \in R$ for every $a \in A$

R is **symmetric** iff $(a,b) \in R$ implies $(b, a) \in R$

R is **transitive** iff $(a,b) \in R$ and $(b, c) \in R$ implies $(a, c) \in R$

Transitive-Reflexive Closure



Add the minimum possible number of edges to make the relation transitive and reflexive

The transitive-reflexive closure of a relation R is the connectivity relation R^*

Finite state machines

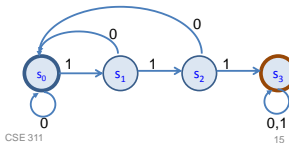
States

Transitions on inputs

Start state and final states

The language recognized by a machine is the set of strings that reach a final state

State	0	1
s_0	s_0	s_1
s_1	s_0	s_2
s_2	s_0	s_3
s_3	s_3	s_3



What language does this machine recognize?

