

# CSE 311 Foundations of Computing I

Lecture 5, Boolean Logic and Predicates  
Autumn 2011

## Announcements

- Reading assignments
  - Boolean Algebra
    - 12.1 – 12.3 7<sup>th</sup> Edition
    - 11.1 – 11.3 6<sup>th</sup> Edition
    - 10.1 – 10.3 5<sup>th</sup> Edition
  - Predicates and Quantifiers
    - 1.4, 1.5 7<sup>th</sup> Edition
    - 1.3, 1.4 5<sup>th</sup> and 6<sup>th</sup> Edition

## Highlights from last lecture

- Boolean algebra to circuit design
- Boolean algebra
  - a set of elements  $B = \{0, 1\}$
  - binary operations  $\{+, \cdot\}$
  - and a unary operation  $\{\prime\}$
  - such that the following axioms hold:



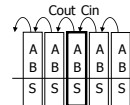
George Boole – 1854

1. the set  $B$  contains at least two elements:  $a, b$
2. closure:  $a + b$  is in  $B$
3. commutativity:  $a + b = b + a$
4. associativity:  $a + (b + c) = (a + b) + c$
5. identity:  $a + 0 = a$
6. distributivity:  $a + (b \cdot c) = (a + b) \cdot (a + c)$
7. complementarity:  $a + a' = 1$

- $a \cdot b$  is in  $B$
- $a \cdot b = b \cdot a$
- $a \cdot (b \cdot c) = (a \cdot b) \cdot c$
- $a \cdot 1 = a$
- $a \cdot (b + c) = (a \cdot b) + (a \cdot c)$
- $a \cdot a' = 0$

## A simple example: 1-bit binary adder

- Inputs: A, B, Carry-in
- Outputs: Sum, Carry-out



A	B	Cin	Cout	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$Cout = B Cin + A Cin + AB$$

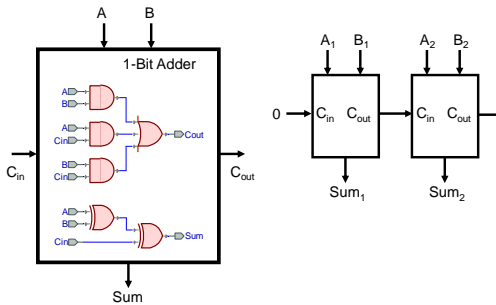
$$S = A' B' Cin + A' B Cin' + A B' Cin' + A B Cin$$

$$= A' (B' Cin + B Cin') + A (B' Cin' + B Cin)$$

$$= A' Z + A Z'$$

$$= A \text{ xor } Z = A \text{ xor } (B \text{ xor } Cin)$$

## Preview: A 2-bit ripple-carry adder



## Mapping truth tables to logic gates

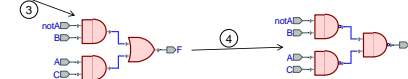
1. Write the Boolean expression
2. Minimize the Boolean expression
3. Draw as gates
4. Map to available gates

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$F = A'BC' + A'BC + AB'C + ABC$$

$$= A'B(C'+C) + AC(B'+B)$$

$$= A'B + AC$$



## Canonical forms

- Truth table is the unique signature of a Boolean function
- The same truth table can have many gate realizations
  - we've seen this already
  - depends on how good we are at Boolean simplification
- Canonical forms
  - standard forms for a Boolean expression
  - we all come up with the same expression

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## Sum-of-products canonical forms

- Also known as disjunctive normal form
- Also known as minterm expansion

A	B	C	F	F'
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	1	0
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$F = 001 \ 011 \ 101 \ 110 \ 111$   
 $F = A'B'C' + A'BC' + AB'C' + ABC' + ABC$   
 $F' = A'B'C' + A'BC' + AB'C'$

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## Sum-of-products canonical form (cont'd)

- Product term (or minterm)
  - ANDed product of literals – input combination for which output is true
  - each variable appears exactly once, true or inverted (but not both)

A	B	C	minterms
0	0	0	A'B'C' m0
0	0	1	A'B'C m1
0	1	0	A'BC' m2
0	1	1	A'BC m3
1	0	0	AB'C' m4
1	0	1	AB'C m5
1	1	0	ABC' m6
1	1	1	ABC m7

$F$  in canonical form:  
 $F(A, B, C) = \sum m(1,3,5,6,7)$   
 $= m1 + m3 + m5 + m6 + m7$   
 $= A'B'C' + A'BC' + AB'C' + ABC' + ABC$

canonical form  $\neq$  minimal form  
 $F(A, B, C) = A'B'C' + A'BC' + AB'C' + ABC' + ABC$   
 $= (A'B' + A'B + AB' + AB)C + ABC'$   
 $= ((A' + A)(B' + B))C + ABC'$   
 $= C + ABC'$   
 $= ABC' + C$   
 $= AB + C$

short-hand notation for minterms of 3 variables

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## Product-of-sums canonical form

- Also known as conjunctive normal form
- Also known as maxterm expansion

A	B	C	F	F'
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	1	0
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$F = 000 \ 010 \ 100$   
 $F = (A + B + C)(A + B' + C)(A' + B + C)$   
 $F' = (A + B + C')(A + B' + C')(A' + B + C')$

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## Product-of-sums canonical form (cont'd)

- Sum term (or maxterm)
  - ORed sum of literals – input combination for which output is false
  - each variable appears exactly once, true or inverted (but not both)

A	B	C	maxterms
0	0	0	A+B+C M0
0	0	1	A+B+C' M1
0	1	0	A+B'+C M2
0	1	1	A+B'+C' M3
1	0	0	A'+B+C M4
1	0	1	A'+B+C' M5
1	1	0	A'+B'+C M6
1	1	1	A'+B'+C' M7

$F$  in canonical form:  
 $F(A, B, C) = \prod M(0,2,4)$   
 $= M0 \cdot M2 \cdot M4$   
 $= (A + B + C)(A + B' + C)(A' + B + C)$

canonical form  $\neq$  minimal form  
 $F(A, B, C) = (A + B + C)(A + B' + C)(A' + B + C)$   
 $= (A + B + C)(A + B' + C)$   
 $= (A + C)(B + C)$

short-hand notation for maxterms of 3 variables

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## S-o-P, P-o-S, and de Morgan's theorem

- Complement of function in sum-of-products form
  - $F' = A'B'C' + A'BC' + AB'C'$
- Complement again and apply de Morgan's and get the product-of-sums form
  - $(F')' = (A'B'C' + A'BC' + AB'C')'$
  - $F = (A + B + C)(A + B' + C)(A' + B + C)$
- Complement of function in product-of-sums form
  - $F' = (A + B + C')(A + B' + C')(A' + B + C')(A' + B' + C')$
- Complement again and apply de Morgan's and get the sum-of-product form
  - $(F')' = ((A + B + C')(A + B' + C')(A' + B + C')(A' + B' + C'))'$
  - $F = A'B'C' + A'BC' + AB'C' + ABC'$

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## Predicate Calculus

- *Predicate or Propositional Function*  
– A function that returns a truth value
- “x is a cat”
- “x is prime”
- “student x has taken course y”
- “ $x > y$ ”
- “ $x + y = z$ ”

## Quantifiers

- $\forall x P(x)$  :  $P(x)$  is true for every  $x$  in the domain
- $\exists x P(x)$  : There is an  $x$  in the domain for which  $P(x)$  is true

## Statements with quantifiers

- $\exists x \text{ Even}(x)$
- $\forall x \text{ Odd}(x)$
- $\forall x (\text{Even}(x) \vee \text{Odd}(x))$
- $\exists x (\text{Even}(x) \wedge \text{Odd}(x))$
- $\forall x \text{ Greater}(x+1, x)$
- $\exists x (\text{Even}(x) \wedge \text{Prime}(x))$

Domain:  
Positive Integers

Even(x)  
Odd(x)  
Prime(x)  
Greater(x,y)  
Equal(x,y)

## Statements with quantifiers

- $\forall x \exists y \text{ Greater}(y, x)$
- $\forall x \exists y \text{ Greater}(x, y)$
- $\forall x \exists y (\text{Greater}(y, x) \wedge \text{Prime}(y))$
- $\forall x (\text{Prime}(x) \rightarrow (\text{Equal}(x, 2) \vee \text{Odd}(x)))$
- $\exists x \exists y (\text{Equal}(x, y + 2) \wedge \text{Prime}(x) \wedge \text{Prime}(y))$

Domain:  
Positive Integers

Even(x)  
Odd(x)  
Prime(x)  
Greater(x,y)  
Equal(x,y)

## Statements with quantifiers

- “There is an odd prime”
- “If x is greater than two, x is not an even prime”
- $\forall x \forall y \forall z ((\text{Equal}(z, x+y) \wedge \text{Odd}(x) \wedge \text{Odd}(y)) \rightarrow \text{Even}(z))$
- “There exists an odd integer that is the sum of two primes”

Domain:  
Positive Integers

Even(x)  
Odd(x)  
Prime(x)  
Greater(x,y)  
Equal(x,y)

## Goldbach’s Conjecture

- Every even integer greater than two can be expressed as the sum of two primes

Even(x)  
Odd(x)  
Prime(x)  
Greater(x,y)  
Equal(x,y)

Domain:  
Positive Integers