



Recursion

[Sections 6.1, 6.3-6.7]

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Recursion

- ◆ A recursive definition is one which is expressed in terms of itself
- ◆ Examples:
 - ◆ "A horse is a four-legged animal which is the progeny of two horses"
 - ◆ Compound interest: "The value after 10 years is equal to the interest rate times the value after 9 years."
 - ◆ An arithmetic expression is either a number, or two arithmetic expressions with a +, -, * or / between them

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Computer Science Examples

- ◆ We've already seen recursive data structures:

```
struct Node {  
    int data;  
    Node *link;  
};
```

- ◆ Functions (and methods) are also allowed to be defined recursively
 - ◆ **Recursive function**: a function that calls itself (possibly indirectly)
 - ◆ Recursive functions are fundamental in computer science

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The Essence of Recursive Computation

- ◆ Recursive functions only work when there's a well-defined notion of **making progress**
- ◆ If a function calls itself to solve a subproblem, two things had better happen:
 - ◆ The subproblem has to be simpler than the original problem
 - ◆ Some subproblems have to be so simple that you can solve them without recursion (bottoming out)

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What to Look for

```
int myRecursiveFunction( Data someInput )  
{  
    if( isReallySimple( someInput ) ) {  
        int solution = computeEasyAnswer( someInput );  
        return solution;  
    } else {  
        Data sub1 = findSubproblem1( someInput );  
        Data subN = findSubproblemN( someInput );  
        int val1 = myRecursiveFunction( sub1 );  
        int valN = myRecursiveFunction( subN );  
        int solution = combineSubproblems( val1, ..., valN );  
        return solution;  
    }  
}
```

Check if we've bottomed out

If so, compute the answer without recursion

If not,

Break the problem into simpler subproblems

Use recursion to solve the subproblems

Combine the solutions to the subproblems into a solution to the original problem

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Factorial

- ◆ $n!$ ("n factorial") can be defined in two ways:

- ◆ Non-recursive definition

$$n! = n * (n-1) * (n-2) \dots * 2 * 1$$

- ◆ Recursive definition

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n(n-1)! & \text{if } n > 0 \end{cases}$$

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The Recursive Factorial Function

```
int factorial( int n )
{
    if( n == 0 ) {
        return 1;
    } else {
        return n * factorial( n - 1 );
    }
}
```

- ◆ Note that the factorial function invokes itself.
- ◆ How can this work?

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Automatic Variables

- ◆ Recall:
 - ◆ Static memory is set aside ahead of time
 - ◆ Dynamic memory is explicitly managed at runtime
- ◆ Automatic memory holds local variables and parameters to functions
 - ◆ Organized into a stack
 - ◆ Size changes dynamically at runtime
 - ◆ Managed automatically by the compiler
 - Block entered: push stack
 - Block exited: pop stack

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Activation Records

- ◆ An activation record is a struct that holds information about a function call
 - ◆ Function parameters, local variables, other stuff
- ◆ A single function can have multiple activation records on the stack simultaneously
 - ◆ This is how recursive functions are possible
 - ◆ Each recursive call gets its own activation record

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Example

```
int factorial( int n )
{
    if( n == 0 ) {
        return 1;
    } else {
        int sub = factorial( n - 1 );
        return n * sub;
    }
}

int main( void )
{
    int x = factorial( 4 );
    cout << "4! = " << x << endl;
}
```

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Infinite Recursion

- ◆ Remember to always check for bottoming out first
 - ◆ Look for "base case"
- ◆ If not...

```
int badFactorial( int n )
{
    int sub = badFactorial( n - 1 );
    if( sub == 1 ) {
        return 1;
    } else {
        return n * sub;
    }
}
```

- ◆ What is the value of `badFactorial(2)`?

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Recursive Algorithms for Recursive Data Structures

- ◆ Many pointer-based data structures use the pointer-to-another-instance idea
- ◆ This recursive structure allows for elegant recursive algorithms

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Printing a Linked List

```
void IntList::print()
{
    printFrom( front );
}

void IntList::printFrom( Node *from )
{
    if( from != NULL ) {
        cout << from->data << " ";
        printFrom( from->link );
    }
}
```

- ◆ How many recursive calls do we make when printing the list <1,2,3,4>?

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Printing in Reverse Order

- ◆ Difficult problem: in our singly-linked list, links only point forward
- ◆ How can we modify recursive list printing to print in reverse?

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Other Recursive List Operations

- ◆ Can also use recursion to, e.g. sum a list

```
int IntList::sumFrom( Node *from )
{
    if( from == NULL ) {
        return 0;
    } else {
        return from->data + sumFrom( from->link );
    }
}
```

- ◆ How would you modify this to count the length of a list? Add N to each element of a list? Find the maximum in the list?

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Recursion: Not Just for Lists

- ◆ Works fine on arrays:

```
void selectionSort( int data[], int size, int from )
{
    if( (size - from) > 1 ) {
        int mi = 0;
        for( int idx = from + 1; idx < size; idx++ ) {
            if( data[ idx ] < data[ mi ] ) {
                mi = idx;
            }
        }
        swap( data[ from ], data[ mi ] );
        selectionSort( data, size, from + 1 );
    }
}
```

- ◆ Another example: trees

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Challenge Problem

- ◆ What does this function do?

```
int mystery( int x )
{
    assert( x > 0 );
    if( x == 1 ) {
        return 0;
    } else {
        return 1 + mystery( x / 2 );
    }
}
```

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Recursion vs. Iteration

- ◆ When to use recursion?
 - ◆ Processing recursive data structures
 - ◆ Some algorithms can be expressed more elegantly in recursive form
 - “Divide & Conquer” algorithms
- ◆ When to use iteration instead?
 - ◆ Nonrecursive data structures
 - ◆ Problems without obvious recursive structure
 - ◆ Sometimes because it's faster

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Which is Better?

- ◆ If a single recursive call is at the very end of the function:
 - ◆ Known as tail recursion
 - ◆ Easy to rewrite iteratively (many good compilers can actually do this for you...)
- ◆ Recursive problems that are not tail recursive are harder to write nonrecursively
 - ◆ Usually have to simulate recursion with a stack

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In Theory...

- ◆ Some programming languages provide no loops (loops are implemented through if and recursion)
 - ◆ "Functional" languages
- ◆ Any iteration can be rewritten using recursion, and vice-versa
- ◆ But that doesn't mean they're equally easy to use in all situations!

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Summary

- ◆ Recursion is something defined in terms of itself
 - ◆ Recursive procedures
 - ◆ Recursive data structures
- ◆ Activation records make it work
- ◆ Two parts of all recursive functions
 - ◆ Base case(s) - bottom out and solve non-recursively
 - ◆ Recursive case(s) - solve a simpler recursive subproblem
 - ◆ Base case always checked first

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