

Building Java Programs

Chapter 13
binary search and complexity

reading: 13.1-13.2

Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```

- Gauss also has a way of solving this

```
public static int sum2(int n) {  
    return n * (n + 1) / 2;  
}
```

- Which one is more efficient?

Runtime Efficiency (13.2)

- **efficiency**: measure of computing resources used by code.
 - can be relative to speed (time), memory (space), etc.
 - most commonly refers to run time
- We want to be able to compare different algorithms to see which is more efficient

Efficiency Try 1

- Let's time the methods!

n = 1	sum1 took	0ms,	sum2 took	0ms
n = 5	sum1 took	0ms,	sum2 took	0ms
n = 10	sum1 took	0ms,	sum2 took	0ms
n = 100	sum1 took	0ms,	sum2 took	0ms
n = 1,000	sum1 took	0ms,	sum2 took	0ms
n = 10,000,000	sum1 took	18ms,	sum2 took	0ms
n = 100,000,000	sum1 took	143ms,	sum2 took	0ms
n = 2,147,483,647	sum1 took	1880ms,	sum2 took	0ms

- Downsides

- Different computers give different run times
- The same computer gives different results!!! D:<

Efficiency – Try 2

- Count number of “simple steps” our algorithm takes to run
- Assume the following:
 - Any single Java statement takes same amount of time to run.
 - `int x = 5;`
 - `boolean b = (5 + 1 * 2) < 15 + 3;`
 - `return x * 3;`
 - A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.
 - A method call's runtime is measured by the total runtime of the statements inside the method's body.



Efficiency examples

```
public static void method1(int N) {
```

```
    statement1;  
    statement2;  
    statement3; }
```

3

```
    for (int i = 1; i <= N; i++) {  
        statement4;  
    }
```

N

```
    for (int i = 1; i <= N; i++) {  
        statement5;  
        statement6;  
        statement7;  
    }  
}
```

3N

$4N + 3$

Efficiency examples 2

```
public static void method2(int N) {  
    for (int i = 1; i <= N; i++) {  
        for (int j = 1; j <= N; j++) {  
            statement1;  
        }  
    }  
  
    for (int i = 1; i <= N; i++) {  
        statement2;  
        statement3;  
        statement4;  
        statement5;  
    }  
}
```

N^2

$4N$

$N^2 + 4N$

- How many statements will execute if $N = 10$? If $N = 1000$?

Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0; } 1  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum; } 1  
}
```

} N

} N + 2

- Gauss also has a way of solving this

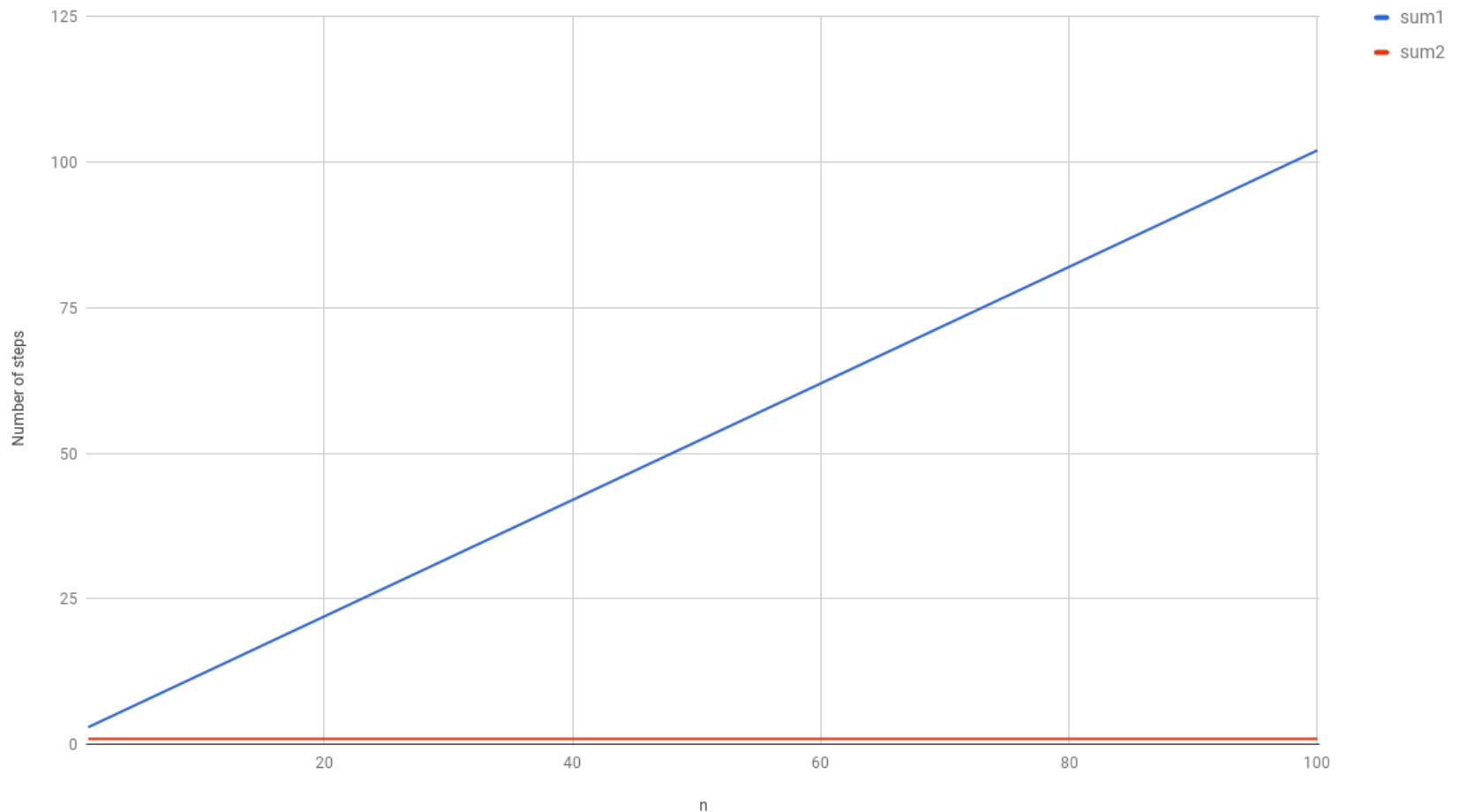
```
public static int sum2(int n) {  
    return n * (n + 1) / 2; } 1  
}
```

} 1

- Which one is more efficient?

Visualizing Difference

Comparing sum1 and sum2



Algorithm growth rates (13.2)

- We measure runtime in proportion to the input data size, N .
 - **growth rate**: Change in runtime as N changes.
- Say an algorithm runs **$0.4N^3 + 25N^2 + 8N + 17$** statements.
 - Consider the runtime when N is *extremely large* .
 - We ignore constants like 25 because they are tiny next to N .
 - The highest-order term (N^3) dominates the overall runtime.
- We say that this algorithm runs "on the order of" N^3 .
- or **$O(N^3)$** for short ("Big-Oh of N cubed")



- Consider this method:

```
public void method(int n) {  
    int value = 0;  
    for (int i = 0; i < 7; i++) {  
        for (int j = 0; j < n; j++) {  
            value += j;  
        }  
    }  
    return value + n / 2;  
}
```

- What is the Big-O efficiency for this method?

- $O(1)$
- $O(n)$
- $O(7n)$
- $O(7n + 4)$
- $O(n^2)$
- $O(n^3)$



Complexity classes

- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N .

Class	Big-Oh	If you double N , ...	Example
constant	$O(1)$	unchanged	10ms
logarithmic	$O(\log_2 N)$	increases slightly	175ms
linear	$O(N)$	doubles	3.2 sec
log-linear	$O(N \log_2 N)$	slightly more than doubles	6 sec
quadratic	$O(N^2)$	quadruples	1 min 42 sec
cubic	$O(N^3)$	multiplies by 8	55 min
...
exponential	$O(2^N)$	multiplies drastically	$5 * 10^{61}$ years

<http://recursive-design.com/blog/2010/12/07/comp-sci-101-big-o-notation/> - post about a Google interview



Range algorithm

What complexity class is this algorithm? Can it be improved?

```
// returns the range of values in the given array;  
// the difference between elements furthest apart  
// example: range({17, 29, 11, 4, 20, 8}) is 25
```

```
public static int range(int[] numbers) {  
    int maxDiff = 0;        // look at each pair of values  
    for (int i = 0; i < numbers.length; i++) {  
        for (int j = 0; j < numbers.length; j++) {  
            int diff = Math.abs(numbers[j] - numbers[i]);  
            if (diff > maxDiff) {  
                maxDiff = diff;  
            }  
        }  
    }  
    return diff;  
}
```


Range algorithm 2

The last algorithm is **$O(N^2)$** . A slightly better version:

```
// returns the range of values in the given array;  
// the difference between elements furthest apart  
// example: range({17, 29, 11, 4, 20, 8}) is 25
```

```
public static int range(int[] numbers) {  
    int maxDiff = 0;        // look at each pair of values  
    for (int i = 0; i < numbers.length; i++) {  
        for (int j = i + 1; j < numbers.length; j++) {  
            int diff = Math.abs(numbers[j] - numbers[i]);  
            if (diff > maxDiff) {  
                maxDiff = diff;  
            }  
        }  
    }  
    return diff;  
}
```

Range algorithm 3

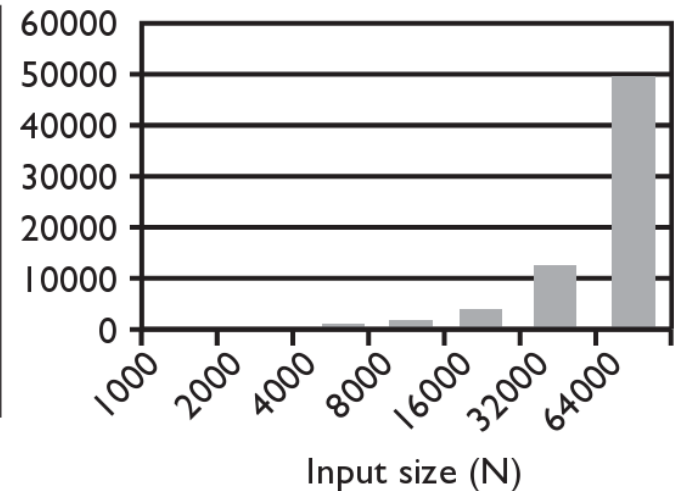
This final version is **O(N)**. It runs MUCH faster:

```
// returns the range of values in the given array;  
// example: range({17, 29, 11, 4, 20, 8}) is 25  
public static int range(int[] numbers) {  
    int max = numbers[0];      // find max/min values  
    int min = max;  
    for (int i = 1; i < numbers.length; i++) {  
        if (numbers[i] < min) {  
            min = numbers[i];  
        }  
        if (numbers[i] > max) {  
            max = numbers[i];  
        }  
    }  
    return max - min;  
}
```

Runtime of first 2 versions

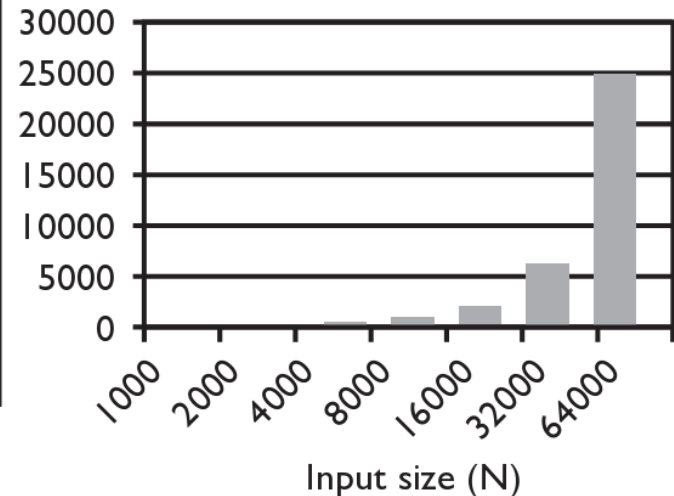
- Version 1:

N	Runtime (ms)
1000	15
2000	47
4000	203
8000	781
16000	3110
32000	12563
64000	49937



- Version 2:

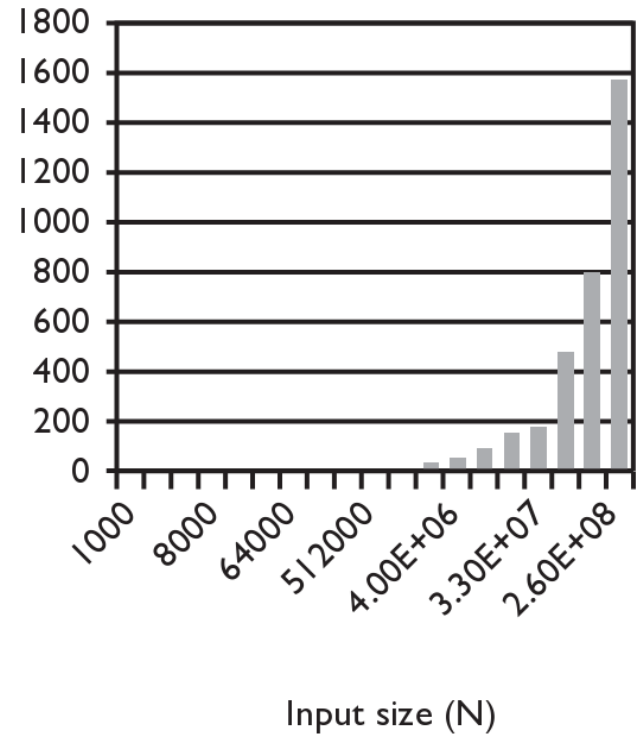
N	Runtime (ms)
1000	16
2000	16
4000	110
8000	406
16000	1578
32000	6265
64000	25031



Runtime of 3rd version

- Version 3:

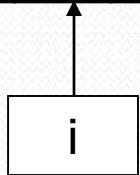
N	Runtime (ms)
1000	0
2000	0
4000	0
8000	0
16000	0
32000	0
64000	0
128000	0
256000	0
512000	0
1e6	0
2e6	16
4e6	31
8e6	47
1.67e7	94
3.3e7	188
6.5e7	453
1.3e8	797
2.6e8	1578



Sequential search

- **sequential search:** Locates a target value in an array / list by examining each element from start to finish. Used in `indexOf`.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



Sequential search

- What is its complexity class?

```
public int indexOf(int value) {  
    for (int i = 0; i < size; i++) {  
        if (elementData[i] == value) {  
            return i;  
        }  
    }  
    return -1;    // not found  
}
```

} N

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

- On average, "only" $N/2$ elements are visited
 - $1/2$ is a constant that can be ignored

Binary search (13.1)

- **binary search:** Locates a target value in a *sorted* array or list by successively eliminating half of the array from consideration.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

min

mid

max

Binary search

- **binary search** successively eliminates half of the elements.
 - *Algorithm:* Examine the middle element of the array.
 - If it is too big, eliminate the right half of the array and repeat.
 - If it is too small, eliminate the left half of the array and repeat.
 - Else it is the value we're searching for, so stop.
 - Which indexes does the algorithm examine to find value **42**?
 - What is the runtime complexity class of binary search?

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

min

mid

max

Binary search runtime

- For an array of size N , it eliminates $\frac{1}{2}$ until 1 element remains.

$N, N/2, N/4, N/8, \dots, 4, 2, 1$

- How many divisions does it take?
- Think of it from the other direction:
 - How many times do I have to multiply by 2 to reach N ?
 $1, 2, 4, 8, \dots, N/4, N/2, N$
 - Call this number of multiplications " x ".

$$2^x = N$$

$$x = \log_2 N$$

- Binary search is in the **logarithmic** complexity class.

<http://recursive-design.com/blog/2010/12/07/comp-sci-101-big-o-notation/> - post about a Google interview



Collection efficiency

- Efficiency of our Java's `ArrayList` and `LinkedList` methods:

Method	ArrayList	LinkedList
add	$O(1)^*$	$O(1)$
add(index , value)	$O(N)$	$O(N)$
indexOf	$O(N)$	$O(N)$
get	$O(1)$	$O(N)$
remove	$O(N)$	$O(N)$
set	$O(1)$	$O(N)$
size	$O(1)$	$O(1)$

* Most of the time!

Throw Back: Unique words

- Recall two weeks ago when we counted the number of unique words in a file. Our first attempt

```
public static int uniqueWords(Scanner input) {  
    List<String> words = new LinkedList<String>();  
    while (input.hasNext()) {  
        String word = input.next();  
        if (!words.contains(word)) {  
            words.add(word);  
        }  
    }  
    return words.size();  
}
```


Throw Back: Unique words

- Recall two weeks ago when we counted the number of unique words in a file. Our second attempt
- We saw briefly that operations on `HashSet` are $O(1)$

```
public static int uniqueWords(Scanner input) {  
    Set<String> words = new HashSet<String>();  
    while (input.hasNext()) {  
        String word = input.next();  
        words.add(word);  
    }  
    return words.size();  
}
```