

Building Java Programs

Chapter 13
binary search and complexity

reading: 13.1-13.2

Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```

- Gauss also has a way of solving this

```
public static int sum2(int n) {  
    return n * (n + 1) / 2;  
}
```

- Which one is more efficient?

Runtime Efficiency (13.2)

- **efficiency**: measure of computing resources used by code.
 - can be relative to speed (time), memory (space), etc.
 - most commonly refers to run time
- We want to be able to compare different algorithms to see which is more efficient

Efficiency Try 1

- Let's time the methods!

n = 1	sum1 took 0ms,	sum2 took 0ms
n = 5	sum1 took 0ms,	sum2 took 0ms
n = 10	sum1 took 0ms,	sum2 took 0ms
n = 100	sum1 took 0ms,	sum2 took 0ms
n = 1,000	sum1 took 0ms,	sum2 took 0ms
n = 10,000,000	sum1 took 18ms,	sum2 took 0ms
n = 100,000,000	sum1 took 123ms,	sum2 took 0ms
n = 2,147,483,647	sum1 took 1880ms,	sum2 took 0ms

- Downsides

- Different computers give different run times
- The same computer gives different results!!! D:<

Efficiency – Try 2

- Let's count number of "steps" our algorithm takes to run
- Assume the following:
 - Any single Java statement takes same amount of time to run.
 - `int x = 5;`
 - `boolean b = (5 + 1 * 2) < 15 + 3;`
 - `System.out.println("Hello");`
 - A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.
 - A method call's runtime is measured by the total runtime of the statements inside the method's body.


Efficiency examples

```
statement1;  
statement2;  
statement3;
```




3

```
for (int i = 1; i <= N; i++) {  
    statement4;  
}
```



N

```
for (int i = 1; i <= N; i++) {  
    statement5;  
    statement6;  
    statement7;  
}
```



3N



$4N + 3$

Efficiency examples 2

```
for (int i = 1; i <= N; i++) {  
    for (int j = 1; j <= N; j++) {  
        statement1;  
    }  
}
```

} N^2

```
for (int i = 1; i <= N; i++) {  
    statement2;  
    statement3;  
    statement4;  
    statement5;  
}
```

} $4N$

} $N^2 + 4N$

- How many statements will execute if $N = 10$? If $N = 1000$?

Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0; } 1  
    for (int i = 1; i <= n; i++) { } N  
        sum += i;  
    }  
    return sum; } 1  
}
```

$N + 2$

- Gauss also has a way of solving this

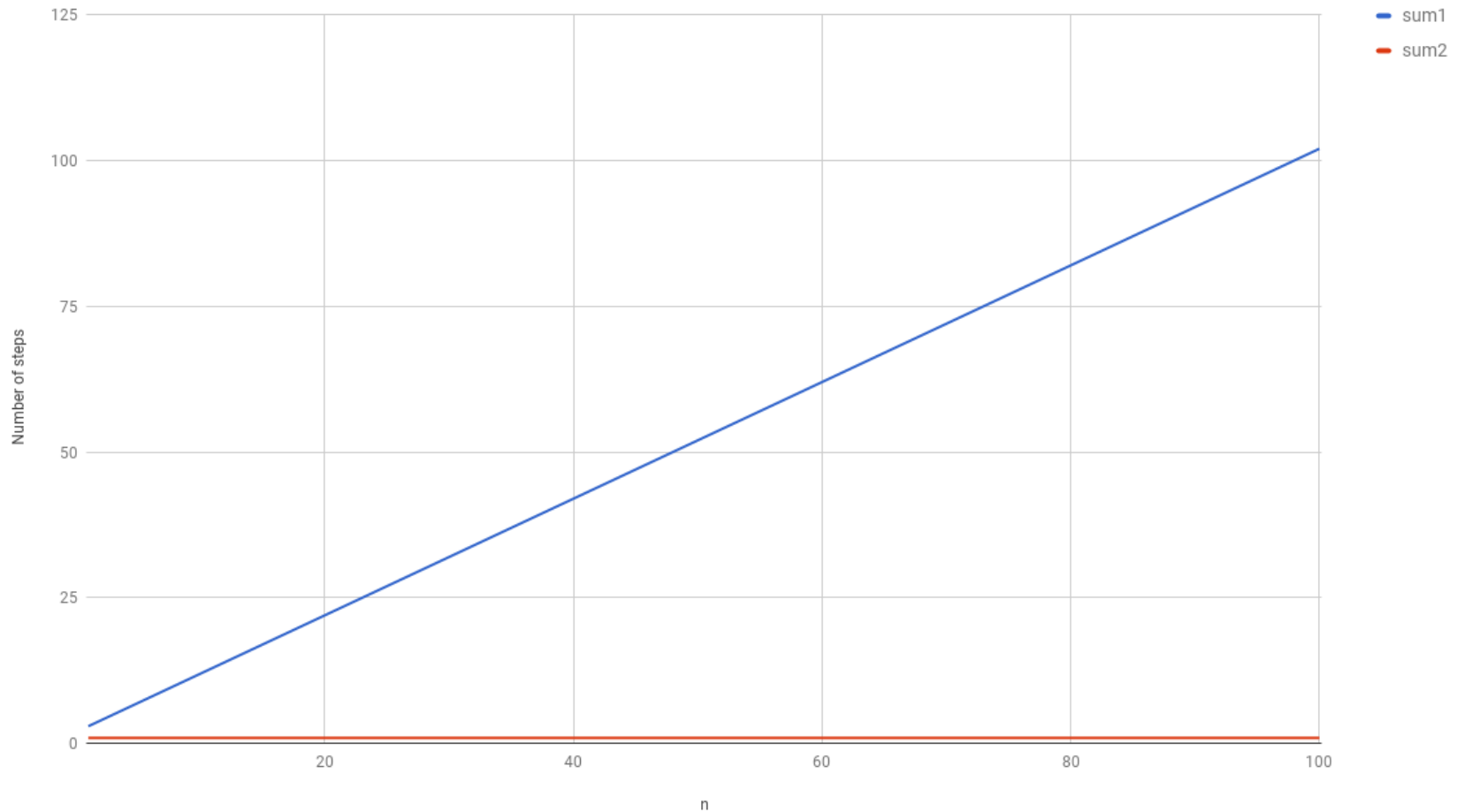
```
public static int sum2(int n) {  
    return n * (n + 1) / 2; } 1  
}
```

1

- Which one is more efficient?

Visualizing Difference

Comparing sum1 and sum2



Algorithm growth rates (13.2)

- We measure runtime in proportion to the input data size, N .
 - **growth rate**: Change in runtime as N changes.
- Say an algorithm runs **$0.4N^3 + 25N^2 + 8N + 17$** statements.
 - Consider the runtime when N is *extremely large* .
 - We ignore constants like 25 because they are tiny next to N .
 - The highest-order term (N^3) dominates the overall runtime.
 - We say that this algorithm runs "on the order of" N^3 .
 - or **$O(N^3)$** for short ("Big-Oh of N cubed")

Complexity classes

- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N .

Class	Big-Oh	If you double N , ...	Example
constant	$O(1)$	unchanged	10ms
logarithmic	$O(\log_2 N)$	increases slightly	175ms
linear	$O(N)$	doubles	3.2 sec
log-linear	$O(N \log_2 N)$	slightly more than doubles	6 sec
quadratic	$O(N^2)$	quadruples	1 min 42 sec
cubic	$O(N^3)$	multiplies by 8	55 min
...
exponential	$O(2^N)$	multiplies drastically	$5 * 10^{61}$ years

Complexity classes

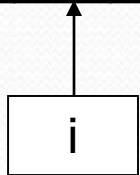
- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N .

Input Size	$O(1)$ steps	$O(N)$ steps	$O(N^2)$ steps	$O(N^3)$ steps
X	1	X	X^2	X^3
$2X$				
$3X$				

Sequential search

- **sequential search:** Locates a target value in an array / list by examining each element from start to finish. Used in `indexOf`.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



- What is a value we could search for that would be “fast”
- The array is sorted. Could we take advantage of this?

Binary search (13.1)

- **binary search:** Locates a target value in a *sorted* array or list by successively eliminating half of the array from consideration.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

Diagram illustrating a binary search on a sorted array. The array is shown with indices 0 to 16 and corresponding values. The value 42 is highlighted in yellow at index 10. Three boxes labeled 'min', 'mid', and 'max' are positioned below the array, with arrows pointing to indices 0, 8, and 16 respectively, indicating the current search range.

Sequential search

- What is its complexity class?

```
public int indexOf(int value) {  
    for (int i = 0; i < size; i++) {  
        if (elementData[i] == value) {  
            return i;  
        }  
    }  
    return -1;    // not found  
}
```

} N

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

- On average, "only" $N/2$ elements are visited
 - $1/2$ is a constant that can be ignored

Binary search

- **binary search** successively eliminates half of the elements.
 - *Algorithm:* Examine the middle element of the array.
 - If it is too big, eliminate the right half of the array and repeat.
 - If it is too small, eliminate the left half of the array and repeat.
 - Else it is the value we're searching for, so stop.
 - Which indexes does the algorithm examine to find value **42**?
 - What is the runtime complexity class of binary search?

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

Diagram illustrating the binary search process on an array. The array contains values: -4, 2, 7, 10, 15, 20, 22, 25, 30, 36, 42, 50, 56, 68, 85, 92, 103. The search target is 42. The current search range is from index 0 (min) to index 16 (max). The middle element (index 10) is highlighted in yellow, indicating it is the current element being examined.

Binary search runtime

- For an array of size N , it eliminates $\frac{1}{2}$ until 1 element remains.

$N, N/2, N/4, N/8, \dots, 4, 2, 1$

- How many divisions does it take?
- Think of it from the other direction:
 - How many times do I have to multiply by 2 to reach N ?
 $1, 2, 4, 8, \dots, N/4, N/2, N$
 - Call this number of multiplications " x ".

$$2^x = N$$

$$\mathbf{x = \log_2 N}$$

- Binary search is in the **logarithmic** complexity class.