## Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {
    int sum = 0;
    for (int i = 1; i <= n; i++) {
        sum += i;
    }
    return sum;
}
```

- Gauss also has a way of solving this public static int sum2(int n) \{

```
    return n * (n + 1) / 2;
```

\}

- Which one is more efficient?


## Runtime Efficiency (13.2)

- efficiency: measure of computing resources used by code.
- can be relative to speed (time), memory (space), etc.
- most commonly refers to run time
- We want to be able to compare different algorithms to see which is more efficient


## Efficiency Try 1

- Let's time the methods!

| $\mathrm{n}=1$ | sum1 took | Oms, | sum2 | took 0ms |
| :---: | :---: | :---: | :---: | :---: |
| $\mathrm{n}=5$ | sum1 took | 0 ms , | sum2 | took 0ms |
| $\mathrm{n}=10$ | sum1 took | 0 ms , | sum2 | took 0ms |
| $\mathrm{n}=100$ | sum1 took | 0 ms , | sum 2 | took 0ms |
| $\mathrm{n}=1,000$ | sum1 took | Qms, | sum2 | took 0ms |
| $\mathrm{n}=10,000,000$ | sum1 took | 18 ms , | sum2 | took 0 ms |
| $\mathrm{n}=100,000,000$ | sum1 took | 12 Bms , | sum2 | took 0ms |
| $\mathrm{n}=2,147,483,647$ | sum1 took1 | 800ms, | sum2 | took 0ms |

- Downsides
- Different computers give different run times
- The same computer gives different results!!! D: <


## Efficiency - Try 2

- Let's count number of "steps" our algorithm takes to run
- Assume the following:
- Any single Java statement takes same amount of time to run.
- int $x=5$;
- boolean $\mathrm{b}=(5+1 * 2)<15+3$;
- System.out.println("Hello");
- A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.
- A method call's runtime is measured by the total runtime of the statements inside the method's body.


## Efficiency examples



## Efficiency examples 2

```
for (int i = 1; i <= N; i++) {
    for (int j = 1; j <= N; j++)
        statement1;
    }
}
for (int i = 1; i <= N; i++) {
    statement2;
    statement3;
    statement4;
    statement5;
}
- How many statements will execute if \(\mathrm{N}=10\) ? If \(\mathrm{N}=1000\) ?
```


## Sum this up for me

- Let's write a method to calculate the sum from 1 to some $n$

```
public static int sum1(int n) {
    int sum = 0;}1
    for (int i = 1; i <= n; i++) {
        sum += i;
    }
    return sum;}1
}
```

- Gauss also has a way of solving this

```
public static int sum2(int n)
    return n * (n + 1) / 2; } 1
```

\}

## Visualizing Difference

Comparing sum1 and sum2


- sum1
- sum2


## Algorithm growth rates (13.2)

- We measure runtime in proportion to the input data size, $N$.
- growth rate: Change in runtime as N changes.
- Say an algorithm runs $\mathbf{0 . 4} \mathbf{N}^{\mathbf{3}}+\mathbf{2 5} \mathbf{N}^{\mathbf{2}}+\mathbf{8 N}+\mathbf{1 7}$ statements.
- Consider the runtime when N is extremely large .
- We ignore constants like 25 because they are tiny next to $N$.
- The highest-order term ( $\mathrm{N}^{3}$ ) dominates the overall runtime.
- We say that this algorithm runs "on the order of" N3.
- or $\mathbf{O}\left(\mathbf{N}^{3}\right)$ for short ("Big-Oh of $N$ cubed")


## Complexity classes

- complexity class: A category of algorithm efficiency based on the algorithm's relationship to the input size N .

| Class | Big-Oh | If you double $\mathbf{N}, \ldots$ | Example |
| :--- | :--- | :--- | :--- |
| constant | $\mathrm{O}(1)$ | unchanged | 10 ms |
| logarithmic | $\mathrm{O}(\log 2 \mathrm{~N})$ | increases slightly | 175 ms |
| linear | $\mathrm{O}(\mathrm{N})$ | doubles | 3.2 sec |
| log-linear | $\mathrm{O}(\mathrm{N}$ log2 N$)$ | slightly more than doubles | 6 sec |
| quadratic | $\mathrm{O}\left(\mathrm{N}^{2}\right)$ | quadruples | 1 min 42 sec |
| cubic | $\mathrm{O}\left(\mathrm{N}^{3}\right)$ | multiplies by 8 | 55 min |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |
| exponential | $\mathrm{O}\left(2^{\mathrm{N}}\right)$ | multiplies drastically | $5 * 10^{61}$ years |

## Complexity classes



## Sequential search

- sequential search: Locates a target value in an array / list by examining each element from start to finish. Used in indexOf.
- How many elements will it need to examine?
- Example: Searching the array below for the value 42:

| index | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | -4 | 2 | 7 | 10 | 15 | 20 | 22 | 25 | 30 | 36 | 42 | 50 | 56 | 68 | 85 | 92 | 103 |

- The array is sorted. Could we take advantage of this?


## Binary search (13.1)

- binary search: Locates a target value in a sorted array or list by successively eliminating half of the array from consideration.
- How many elements will it need to examine?
- Example: Searching the array below for the value 42:

| index | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | -4 | 2 | 7 | 10 | 15 | 20 | 22 | 25 | 30 | 36 | 42 | 50 | 56 | 68 | 85 | 92 | 103 |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

## Sequential search

- What is its complexity class?

```
public int indexOf(int value)
    for (int i = 0; i < size; i++) {
        if (elementData[i] == value) {
        return i;
    } }
}
```

| index | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | -4 | 2 | 7 | 10 | 15 | 20 | 22 | 25 | 30 | 36 | 42 | 50 | 56 | 68 | 85 | 92 | 103 |

- On average, "only" N/2 elements are visited
- $1 / 2$ is a constant that can be ignored


## Binary search

- binary search successively eliminates half of the elements.
- Algorithm: Examine the middle element of the array.
- If it is too big, eliminate the right half of the array and repeat.
- If it is too small, eliminate the left half of the array and repeat.
- Else it is the value we're searching for, so stop.
- Which indexes does the algorithm examine to find value 42 ?
- What is the runtime complexity class of binary search?

| index | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | -4 | 2 | 7 | 10 | 15 | 20 | 22 | 25 | 30 | 36 | 42 | 50 | 56 | 68 | 85 | 92 | 103 |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

## Binary search runtime

- For an array of size $N$, it eliminates $1 / 2$ until 1 element remains.

$$
\text { N, N/2, N/4, N/8, ..., 4, 2, } 1
$$

- How many divisions does it take?
- Think of it from the other direction:
- How many times do I have to multiply by 2 to reach N ?

$$
1,2,4,8, \ldots, N / 4, N / 2, N
$$

- Call this number of multiplications "x".

$$
\begin{aligned}
& 2 x=N \\
& x=\log _{2} N
\end{aligned}
$$

- Binary search is in the logarithmic complexity class.


## Collection efficiency

- Efficiency of our Java's ArrayList and LinkedList methods:

| Method | ArrayList | LinkedList |
| :--- | :--- | :--- |
| add | $\mathrm{O}(1)^{*}$ | $\mathrm{O}(1)^{* *}$ |
| add (index, value) | $\mathrm{O}(\mathrm{N})$ | $\mathrm{O}(\mathrm{N})$ |
| indexOf | $\mathrm{O}(\mathrm{N})$ | $\mathrm{O}(\mathrm{N})$ |
| get | $\mathrm{O}(1)$ | $\mathrm{O}(\mathrm{N})$ |
| remove | $\mathrm{O}(\mathrm{N})$ | $\mathrm{O}(\mathrm{N})$ |
| set | $\mathrm{O}(1)$ | $\mathrm{O}(\mathrm{N})$ |
| size | $\mathrm{O}(1)$ | $\mathrm{O}(1)^{* * *}$ |

* Most of the time!
** Assuming we have a reference to the back of the list
*** Assuming we have a size field

