



# Building Java Programs

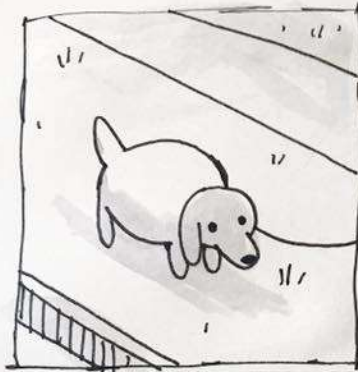
Chapter 13  
binary search and complexity

**reading: 13.1-13.2**



10/07/2018

# DOGS SPOTTED THIS WEEKEND



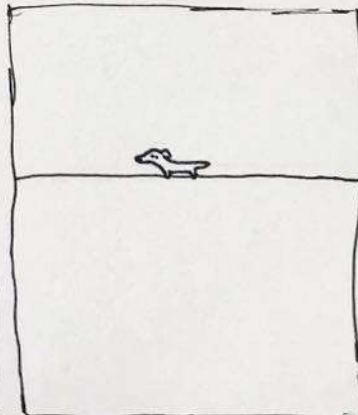
CHUNK ON THE STREET



DAZZLING SMILE



WAITING PATIENTLY



LITTLE WEENIE, TOO FAR AWAY



# Road Map

## CS Concepts

- Client/Implementer
- Efficiency
- Recursion
- Regular Expressions
- Grammars
- Sorting
- Backtracking
- Hashing
- Huffman Compression

## Data Structures

- Lists
- Stacks
- Queues
- Sets
- Maps
- Priority Queues

## Java Language

- Exceptions
- Interfaces
- References
- Comparable
- Generics
- Inheritance/Polymorphism
- Abstract Classes

## Java Collections

- Arrays
- ArrayList ✕
- LinkedList ✕
- Stack
- TreeSet / TreeMap
- HashSet / HashMap
- PriorityQueue



# Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```

- Gauss also has a way of solving this

```
public static int sum2(int n) {  
    return n * (n + 1) / 2;  
}
```

- Which one is more efficient?



# Runtime Efficiency (13.2)

- **efficiency**: measure of computing resources used by code.
  - can be relative to speed (time), memory (space), etc.
  - most commonly refers to run time
- We want to be able to compare different algorithms to see which is more efficient



# Efficiency Try 1

- Let's time the methods!

n = 1	sum1 took 0ms,	sum2 took 0ms
n = 5	sum1 took 0ms,	sum2 took 0ms
n = 10	sum1 took 0ms,	sum2 took 0ms
n = 100	sum1 took 0ms,	sum2 took 0ms
n = 1,000	sum1 took 0ms,	sum2 took 0ms
n = 10,000,000	sum1 took 18ms,	sum2 took 0ms
n = 100,000,000	sum1 took 123ms,	sum2 took 0ms
n = 2,147,483,647	sum1 took 1880ms,	sum2 took 0ms

- Downsides

- Different computers give different run times
- The same computer gives different results!!! D:<




# Efficiency – Try 2

- Count number of “simple steps” our algorithm takes to run
- Assume the following:
  - Any single Java statement takes same amount of time to run.
    - `int x = 5;`
    - `boolean b = (5 + 1 * 2) < 15 + 3;`
    - `System.out.println("Hello");`
  - A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.
  - A method call's runtime is measured by the total runtime of the statements inside the method's body.




# Efficiency examples


```
public static void method1() {  
    statement1;  
    statement2;  
    statement3;  
}
```



```
    for (int i = 1; i <= N; i++) {  
        statement4;  
    }
```



```
    for (int i = 1; i <= N; i++) {  
        statement5;  
        statement6;  
        statement7;  
    }
```

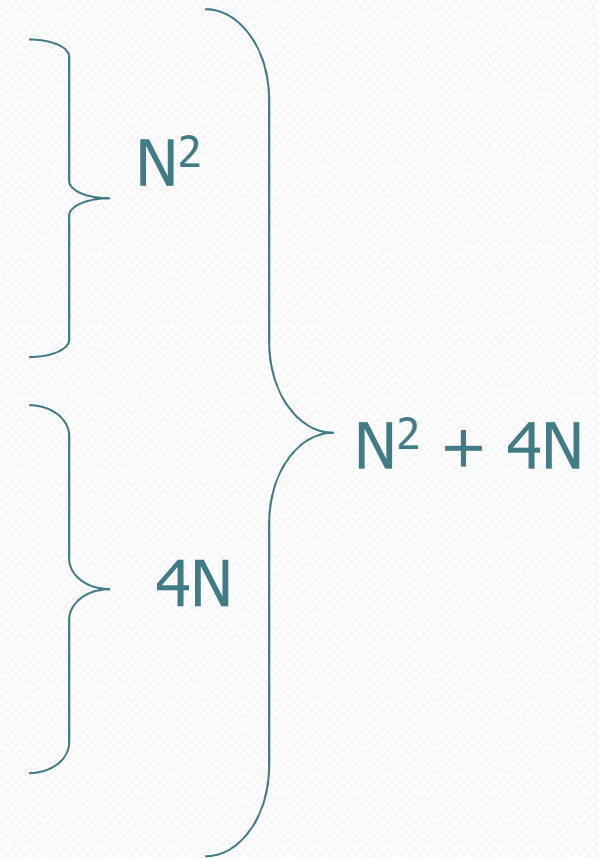


$4N + 3$



# Efficiency examples 2

```
public static void method2() {  
    for (int i = 1; i <= N; i++) {  
        for (int j = 1; j <= N; j++) {  
            statement1;  
        }  
    }  
  
    for (int i = 1; i <= N; i++) {  
        statement2;  
        statement3;  
        statement4;  
        statement5;  
    }  
}
```



- How many statements will execute if  $N = 10$ ? If  $N = 1000$ ?



# Sum this up for me

- Let's write a method to calculate the sum from 1 to some n

```
public static int sum1(int n) {  
    int sum = 0; } 1  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum; } 1  
}
```

} N

} N + 2

- Gauss also has a way of solving this

```
public static int sum2(int n) {  
    return n * (n + 1) / 2; } 1  
}
```

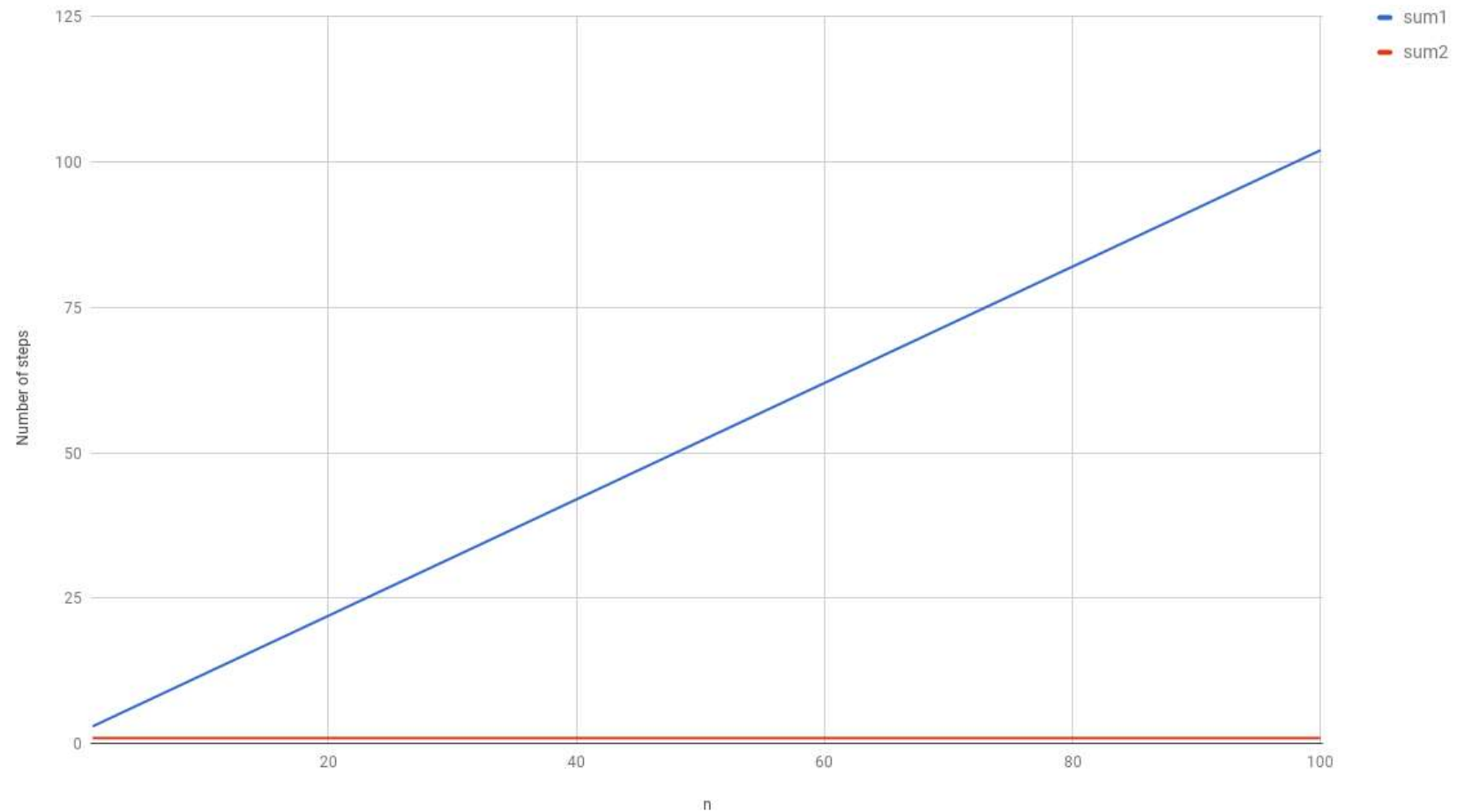
} 1

- Which one is more efficient?



# Visualizing Difference

Comparing sum1 and sum2





# Algorithm growth rates (13.2)

- We measure runtime in proportion to the input data size,  $N$ .
  - **growth rate**: Change in runtime as  $N$  changes.
- Say an algorithm runs  **$0.4N^3 + 25N^2 + 8N + 17$**  statements.
  - Consider the runtime when  $N$  is *extremely large* .
  - We ignore constants like 25 because they are tiny next to  $N$ .
  - The highest-order term ( $N^3$ ) dominates the overall runtime.
- We say that this algorithm runs "on the order of"  $N^3$ .
- or  **$O(N^3)$**  for short ("Big-Oh of  $N$  cubed")

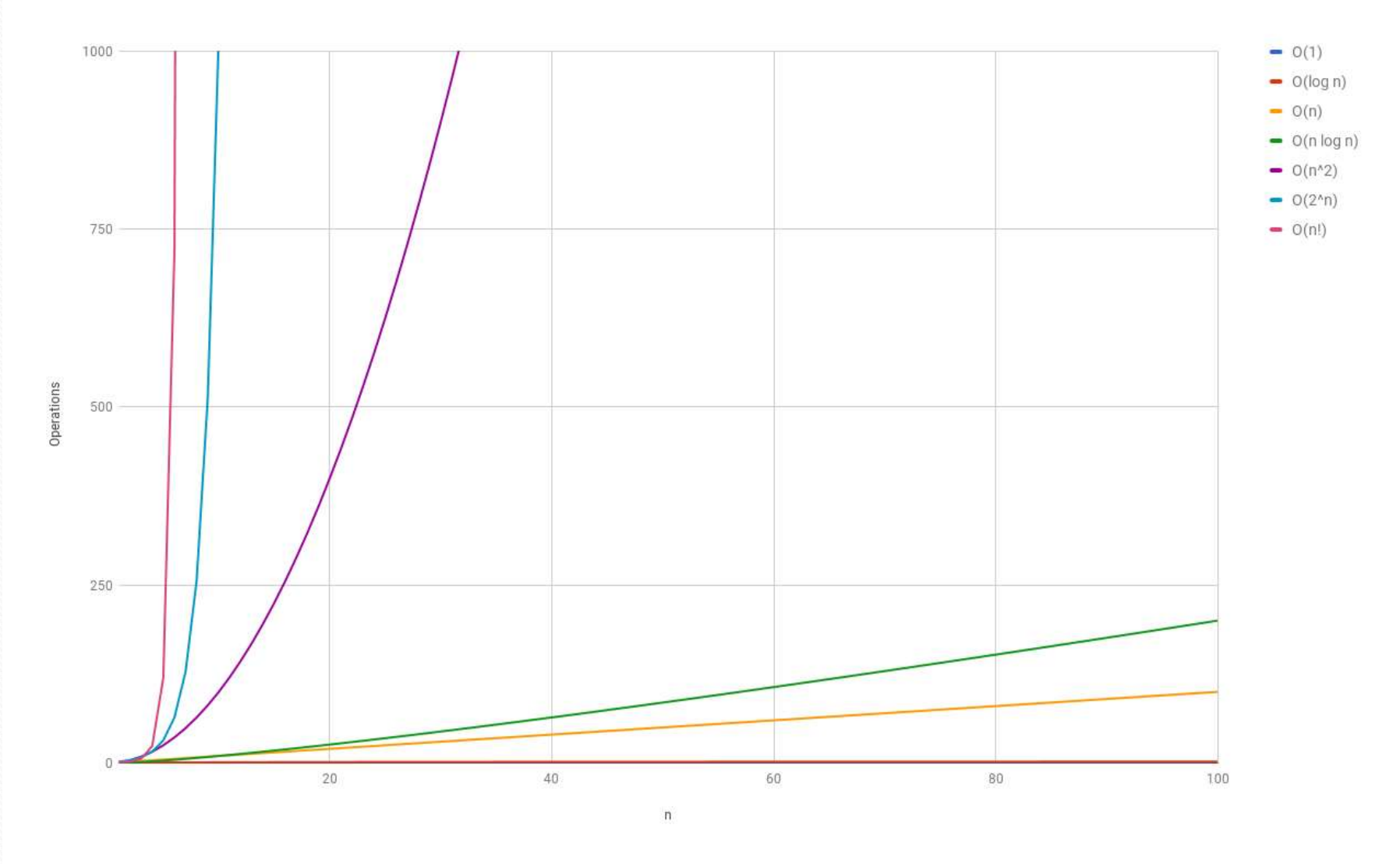


# Complexity classes

- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size  $N$ .

Class	Big-Oh	If you double $N$ , ...	Example
constant	$O(1)$	unchanged	10ms
logarithmic	$O(\log_2 N)$	increases slightly	175ms
linear	$O(N)$	doubles	3.2 sec
log-linear	$O(N \log_2 N)$	slightly more than doubles	6 sec
quadratic	$O(N^2)$	quadruples	1 min 42 sec
cubic	$O(N^3)$	multiplies by 8	55 min
...	...	...	...
exponential	$O(2^N)$	multiplies drastically	$5 * 10^{61}$ years



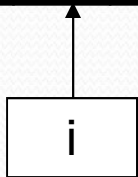




# Sequential search

- **sequential search:** Locates a target value in an array / list by examining each element from start to finish. Used in `indexOf`.
  - How many elements will it need to examine?
  - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



- The array is sorted. Could we take advantage of this?



# Sequential search

- What is its complexity class?

```
public int indexOf(int value) {  
    for (int i = 0; i < size; i++) {  
        if (elementData[i] == value) {  
            return i;  
        }  
    }  
    return -1;    // not found  
}
```

} N

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

- On average, "only"  $N/2$  elements are visited
  - $1/2$  is a constant that can be ignored



# Binary search (13.1)

- **binary search:** Locates a target value in a *sorted* array or list by successively eliminating half of the array from consideration.
  - How many elements will it need to examine?
  - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

min

mid

max



# Binary search

- **binary search** successively eliminates half of the elements.
  - *Algorithm:* Examine the middle element of the array.
    - If it is too big, eliminate the right half of the array and repeat.
    - If it is too small, eliminate the left half of the array and repeat.
    - Else it is the value we're searching for, so stop.
  - Which indexes does the algorithm examine to find value **42**?
  - What is the runtime complexity class of binary search?

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

min

mid

max



# Binary search runtime

- For an array of size  $N$ , it eliminates  $\frac{1}{2}$  until 1 element remains.

$N, N/2, N/4, N/8, \dots, 4, 2, 1$

- How many divisions does it take?
- Think of it from the other direction:
  - How many times do I have to multiply by 2 to reach  $N$ ?  
 $1, 2, 4, 8, \dots, N/4, N/2, N$
  - Call this number of multiplications " $x$ ".

$$2^x = N$$

$$x = \log_2 N$$

- Binary search is in the **logarithmic** complexity class.



# Collection efficiency

- Efficiency of our Java's `ArrayList` and `LinkedList` methods:

Method	ArrayList	LinkedList
add	$O(1)^*$	$O(1)$
add( <b>index</b> , <b>value</b> )	$O(N)$	$O(N)$
indexOf	$O(N)$	$O(N)$
get	$O(1)$	$O(N)$
remove	$O(N)$	$O(N)$
set	$O(1)$	$O(N)$
size	$O(1)$	$O(1)$

\* Most of the time!