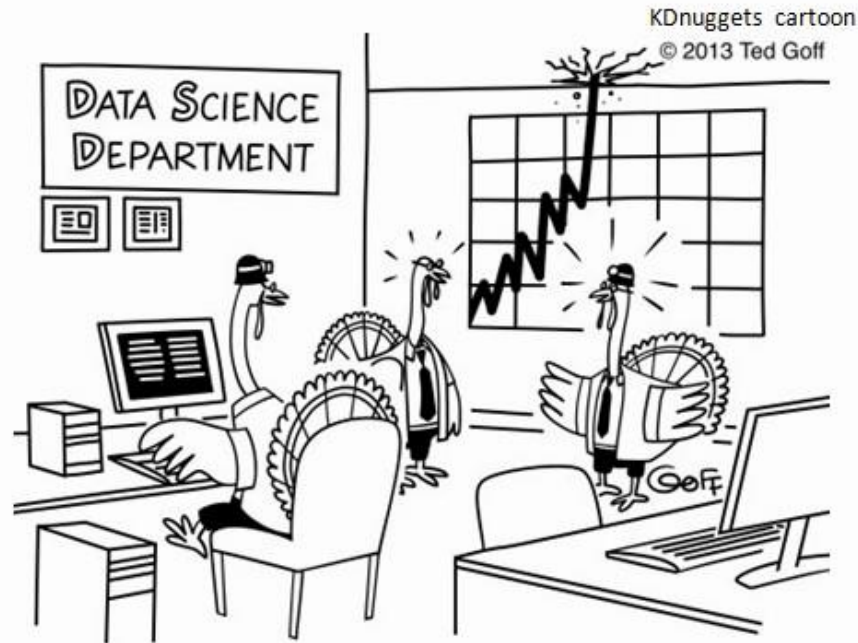


CSE 143

Lecture 23: quick sort



**“I don’t like the look of this.
Searches for gravy and turkey stuffing
are going through the roof!”**

Quick sort

- **quick sort:** Orders a list of values by partitioning the list around one element called a *pivot*, then sorting each partition.
 - invented by British computer scientist C.A.R. Hoare in 1960
- Quick sort is another divide and conquer algorithm:
 - Choose one element in the list to be the pivot.
 - *Divide* the elements so that all elements less than the pivot are to its left and all greater (or equal) are to its right.
 - *Conquer* by applying quick sort (recursively) to both partitions.
- Runtime: $O(N \log N)$ average, $O(N^2)$ worst case.
 - Generally somewhat faster than merge sort.

Quick sort example

index	0	1	2	3	4	5	6	7	8	9
value	65	23	81	43	92	39	57	16	75	32

choose pivot=65

32	23	81	43	92	39	57	16	75	65
32	23	16	43	92	39	57	81	75	65
32	23	16	43	57	39	92	81	75	65
32	23	16	43	57	39	92	81	75	65
32	23	16	43	57	39	65	81	75	92

swap pivot (65) to end

swap 81, 16

swap 57, 92

swap pivot back in

recursively quicksort each half

32	23	16	43	57	39
39	23	16	43	57	32
16	23	39	43	57	32
16	23	32	43	57	39

pivot=32

swap to end

swap 39, 16

swap 32 back in

81	75	92
92	75	81
75	92	81
75	81	92

pivot=81

swap to end

swap 92, 75

swap 81 back in

...

...

Choosing a "pivot"

- The algorithm will work correctly no matter which element you choose as the pivot.
 - A simple implementation can just use the first element.
- But for efficiency, it is better if the pivot divides up the array into roughly equal partitions.
 - What kind of value would be a good pivot? A bad one?

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	8	18	12	-4	27	30	36	50	7	68	91	56	2	85	42	98	25

Choosing a better pivot

- Choosing the first element as the pivot leads to very poor performance on certain inputs (ascending, descending)
 - does not partition the array into roughly-equal size chunks
- Alternative methods of picking a pivot:
 - *random*: Pick a random index from $[min .. max]$
 - *median-of-3*: look at left/middle/right elements and pick the one with the medium value of the three:
 - $a[min]$, $a[(max+min)/2]$, and $a[max]$
 - better performance than picking random numbers every time
 - provides near-optimal runtime for almost all input orderings

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	8	18	91	-4	27	30	86	50	65	78	5	56	2	25	42	98	31

Stable sorting

- **stable sort**: One that maintains relative order of "equal" elements.
 - important for secondary sorting, e.g.
 - sort by name, then sort again by age, then by salary, ...
- All of the N^2 sorts shown are stable.
 - bubble, selection, insertion, shell
- Merge sort is stable.
- Quick sort is *not* stable.
 - The partitioning algorithm can reverse the order of "equal" elements.
 - For this reason, Java's `Arrays/Collections.sort()` use merge sort.

Unstable sort example

- Suppose you want to sort these points by Y first, then by X:
 - $[(4, 2), (5, 7), (3, 7), (3, 1)]$
- A stable sort like merge sort would do it this way:
 - $[(3, 1), (4, 2), (5, 7), (3, 7)]$ sort by y
 - $[(3, 1), (3, 7), (4, 2), (5, 7)]$ sort by x
 - Note that the relative order of (3, 1) and (3, 7) is maintained.
- Quick sort might leave them in the following state:
 - $[(3, 1), (4, 2), (5, 7), (3, 7)]$ sort by y
 - $[(3, 7), (3, 1), (4, 2), (5, 7)]$ sort by x
 - Note that the relative order of (3, 1) and (3, 7) has reversed.