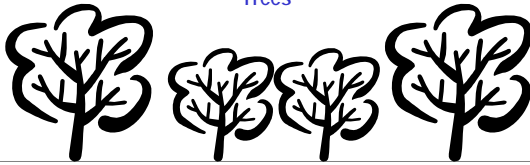


CSE 143 Notes 5/15/06

Trees



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Overview

- Topics
 - Trees: Definitions and terminology
 - Binary trees
 - Tree traversals
 - Recursive tree algorithms



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Trees

- Most of the structures we've looked at so far are linear
 - Arrays
 - Linked lists
- There are many examples of structures that are not linear
 - Organization charts
 - Book contents (chapters, sections, paragraphs)
 - Class inheritance diagrams
- **Trees** can be used to represent hierarchical structures

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Looking Ahead To An Old Goal

- Finding algorithms and data structures for fast searching
 - Sorted arrays are faster than unsorted arrays, for searching
 - Can use binary search algorithm
 - Not so easy to keep the array in order
 - **LinkedLists** were faster than arrays (or ArrayLists), for insertion and removal operations
 - The extra flexibility of the "next" pointers avoided the cost of sliding
 - But... LinkedLists are hard to search, even if sorted
- Is there a way to get the best of both worlds?
- The answer will be...Yes: a particular type of tree

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Drawing Trees

- For whatever reason, computer scientists usually draw trees upside down with the root at the top



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Tree Definitions (1)

- A **tree** is a collection of **nodes** connected by **edges**
- A **node** contains
 - Data (e.g. an int, an Object, or whatever we want)
 - References (edges) to two or more **subtrees** or **children**
- Equivalently: a **tree** is either
 - An empty tree, or
 - A root node with left and right subtrees
- Both definitions are recursive: the first focuses on the implementation (nodes, edges), while the second is a bit more abstract (trees & subtrees)
 - We'll look at trees both ways depending on the situation
 - Often we will use this structure to help formulate algorithms and analysis (recursive data ↔ recursive algorithms)

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Tree Definitions (2)

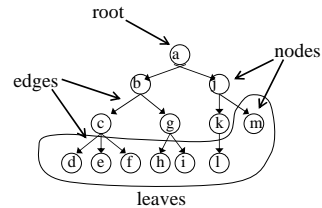
- Trees are hierarchical
 - A node is said to be the *parent* of its *children* (subtrees)
 - We can also speak of the collection of *ancestors* (parent, grandparent, ...) and *descendants* (children, grandchildren) of a node
 - There is a single unique *root* node that has no parent
 - Nodes with no children are called *leaf nodes*
 - Nodes with one or more children are often called *interior* nodes
 - A tree with no nodes is said to be *empty*

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Tree Terminology



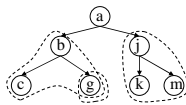
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Subtrees

- A *subtree* in a tree is any node in the tree together with all of its descendants (its children, and their children, recursively)



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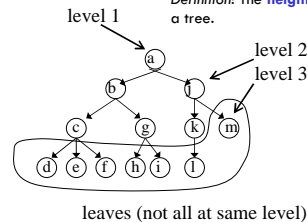
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Level and Height

Definition: The root has *level 1*

Children have level 1 greater than their parent

Definition: The *height* is the highest level of any node in a tree.



leaves (not all at same level)

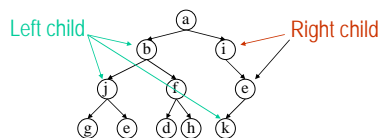
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Binary Trees

- A *binary tree* is a tree each of whose nodes has no more than two children
 - The two children are called the *left child* and *right child*
 - The subtrees rooted at those children are called the *left subtree* and the *right subtree*



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Binary Tree Nodes

- A node for a binary tree holds some sort of data and references to its subtrees

- For example, tree nodes to hold a integer values

```
class TreeNode {
    public int data;           // data item in this node
    public TreeNode left;      // left subtree, or null if none
    public TreeNode right;     // right subtree, or null if none
    public TreeNode(int data, TreeNode left, TreeNode right) { ... }
}
```

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Binary Tree Implementation

- A collection that uses a tree as its underlying data structure normally just needs a single instance variable pointing to the root node, or null if the tree is empty

(The fact that a tree is the underlying data structure is usually a private detail, just as the use of an array or linked list is private in a list structure)

```
// collection of integers
public class IntCollection {
    private TreeNode root; // root of tree, or null if empty
    public IntCollection() { root = null; }
    ...
}
```

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Tree Algorithms

- The definition of a tree is naturally recursive:
 - A tree is either null (empty),
or data + left (sub-)tree + right (sub-)tree
 - Base case(s)?
 - Recursive case(s)?
- Given a recursively defined data structure, recursion is often a very natural technique for algorithms on that data structure
 - Don't fight it!

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A Typical Tree Algorithm: nPositive()

```
public class IntCollection {
    ...
    /** Return the number of positive (>0) ints stored in this tree */
    public int nPositive() {
        return nPositive(root);
    }
    // Return the number of nodes with positive ints in the (sub-)tree with root r
    private int nPositive(TreeNode r) {
        if (r == null) {
            return 0;
        } else {
            // ...
        }
    }
}
```

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Tree Traversal

- Functions like nPositive systematically “visit” each node in a tree
 - This is called a *traversal*
 - We also used this word in connection with lists
- Traversal is a common pattern in many algorithms
 - The processing done during the “visit” varies with the algorithm
- What order should nodes be visited in?
 - Many are possible
 - Three have been singled out as particularly useful for binary trees: *preorder*, *postorder*, and *inorder*

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Traversals

- **Preorder** traversal:
 - “Visit” the (current) node first
i.e., do whatever processing is to be done
 - Then, (recursively) do preorder traversal on its children, left to right
- **Postorder** traversal:
 - First, (recursively) do postorder traversals of children, left to right
 - Visit the node itself last
- **Inorder** traversal:
 - (Recursively) do inorder traversal of left child
 - Then visit the (current) node
 - Then (recursively) do inorder traversal of right child

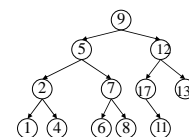
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Example of Tree Traversal

In what order are the nodes visited, if we start the process at the root?



Preorder:

Inorder:

Postorder:

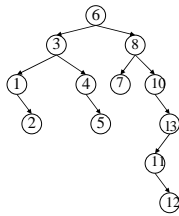
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More Practice

What about this tree?



Preorder:

Inorder:

Postorder:

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New Algorithm: *contains*

- Return whether or not a value is in the tree

```

public class IntCollection {
    ...
    /** Return whether n is in the tree */
    public boolean contains(int n) {
        return contains(root, n);
    }
    // Return whether n is in (sub-)tree with root r
    private boolean contains(TreeNode r, int n) {
        if (r == null) {
            return _____;
        } else if (r.data == n) {
            return _____;
        } else {
            return _____;
        }
    }
}
  
```

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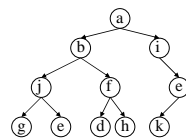
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Test

contains(d)

contains(c)



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Cost of *contains*

- Work done at each node:
- Number of nodes visited:
- Total cost:
- Can we do better?

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