| CSE 143 |  |
| :---: | :---: |
| Stacks |  |
| [Chapter 6] |  |
|  |  |
|  |  |
|  |  |
|  |  |

## Collections in 143

We've seens Arrays and Lists (Vectors) already
We'll continue with Stacks, then Queues, then
Trees and perhaps a couple of others
Other collection types can be programmed and/or invented by you!
But knowing what's already been invented can keep you from "reinventing the wheel."

## Collections

- Collections or containers are ADTs that hold many items of data, usually all of the same type - We can think of arrays and lists as collections - Some programming languages (or their libraries) support various types of containers directly C++ has the Standard Template Library or STL includes many generic algorithms and container types - But to use the library you have to understand the container concepts as well as many advanced language features


## Problem: Switching Boxcars

- The back of three trains meet at a Y-junction
- Can only add/remove from the caboose-end of each train
-Goal: get the right cars, in the right order, on the trains.




## Typing and Correcting Chars

- Type characters, use backspace (<) to mean "erase the previous character"
-The most recently typed unerased char is the one erased

| Sample |  |  |
| :---: | :---: | :---: |
| Action | - Result |  |
| - type h | - h |  |
| - type e | - he |  |
| - type I | - hel |  |
| - type o | - helo |  |
| - type < | - hel |  |
| - type I | - hell |  |
| - type w | - hellw |  |
| - type < | - hell |  |
| - type < | - hel |  |
| - type < | - he |  |
| - type < | - h |  |
| - type i | - hi |  |
|  |  | 0.7 |

## What's common?

I have data to process

- boxcars; characters
- The order of adding data is remembered
- I can only remove or affect what I most recently put in
We say the data structure is LIFO or Last In, First Out, and we call it a Stack.
The point where you can add data is called the Top.
boxcar train: Top is the end of the train
character line: Top is the rightmost character



## Abstract Stack Operations

- push(item): Adds an element to top of stack, increasing stack height by one
- item pop(): Removes topmost element from stack and returns it, decreasing stack height by one
- item top(): Returns a copy of topmost element of stack, leaving stack unchanged
- No "direct access"
- cannot index to a particular data item
- No way to traverse the collection

What is the result of...

## Stack s

int v1,v2,v3, v4, v5, v6;
s.push(1);
s.push(2);
v1 = s.pop();
s.push(3);
s.push(4);
v2 = s.pop();
s.push(5);
v3 $=$ s.pop();
v4 = s.pop();
v5 = s.pop();
v6 = s.pop();
 v3 $\square$

## Stacks Around Us

-Stack of books on a desk

- Trays in the Husky Den
- Take one from top only
- Tray on bottom was put there first

Discard pile in a card game

- Discard on to top, draw card from top
- Not allowed to see or draw what's underneath
- Towers of Hanoi


## Stacks in CS

Implementing function calls
Activation records go on a stack
Evaluating expressions
-How does a calculator (or compiler) understand $(3+4) / 5$ ? more later
"Backtracking" to systematically try all combinations of possibilities
e.g., to explore paths through a maze

## Stack Example

Show the changes to the stack in the following example:

Stack s;
int i;
s.push(5);
s.push(3);
s.push(9);
i $=$ s.pop();
i = s.top();
s.push(6);
s.push(4);

$7 / 17 / 00 \quad 0-14$

## A Stack Client

```
// Goal: Read numbers and print in reverse order
void ReverseNumbers() {
    IntStack s;
    int oneNumber;
while ( cin >> oneNumber ) {
    s.push(oneNumber);
}
while (!s.isEmpty() )
    oneNumber = s.pop();
    cout << oneNumber << endl;
}
```


## Ex.: Parenthesis Matching

Are these legal expressions?

$$
\begin{aligned}
& a+\left(b^{*}[c+d] /(c-d)\right) \\
& \left.\left.a+)\left(b^{*}[c+d]\right]\right] /(c-d)\right] \\
& a+\left(b^{*}[c+d) /(c-d)\right]
\end{aligned}
$$

- How to tell?

Method 1: match "inside out
Good "by eye" but hard to program
Method 2: Via Stack - push '(', '['; pop \& match ')', ']' Harder "by eye", but easier to program.

## Parenthesis Matching

```
char c;
charStack s;
while (cin >> c) {
if ( c == '(' || c== '[' )
        s.push( c );
else if ( c == ')' )
        if ( s.pop() != '(' ) ... // error
else if ( c==']' )
        if ( s.pop() != '[' ) ... // error
}
if ( !s.empty() ) ... //error
cout << "Parens all match."
//NEEDS TEST FOR PREMATURE EMPTY, TOO,&1700 0-18
```


## Stack Implementations

Many possible implementations

- Array-based

Linked list

- Or even, using already implemented Vector ADT

As implementer, use other ADTs to make job easier

- Don't reinvent the wheel for every problem

Often simplifies job to reuse pieces when possible
We'll use stack of ints as an example could have stack of any type of data item

## Stack Via Vector (2)

```
    IntStack::IntStack() { }
        // don't need to do anything, why?
    bool IntStack::isEmpty() {
        return items.isEmpty();
}
```

```
Review: Vector Interface
```

```
class Vector {
```

class Vector {
public:
public:
Vector ( );
Vector ( );
bool isEmpty();
bool isEmpty();
int length ();
int length ();
void insert (int newPosition, int newltem);
void insert (int newPosition, int newltem);
int delete (int position);
int delete (int position);
int retrieve (int position);
int retrieve (int position);
}
}
Review: Vector Interface

```

\section*{Stack Via Vector ADT}

We'll use a private Vector variable.
```

\#include "Vector.h"
class IntStack \{
public:
IntStack() ;
bool isEmpty(); // is the stack empty?
// etc etc -- all the Stack operations
private:
Vector items;
\};
Note: no Top variable! Always use the front of the list as the Top.
7/17/00 $\quad 0-20$

```

\section*{Stack Via Vector (3)}
void IntStack::push(int item) \{
        items.insert(0, item);
\}
int IntStack::top() \{
    //FILL THIS IN
    //HINT: can be done in one line of code
    \}
int IntStack::pop() \{
    //FILL THIS IN
\}

Possible Implementations
- Many possible implementations
- Array-based
- Linked list
- Or even, using already implemented List ADT

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\section*{Stack Via Linked List}

Another implementation technique
- Main idea: keep a linked list, with private "top" pointer to the front of the list
Add new data as a new link to the beginning of the linked list
Pop/top: remove/return the beginning of the linked list
- Not the only way -- could have decided to make top be the end of the list
Important thing is to choose a way;document it; and stick with it.

\section*{Stack Via Linked List (3)}
// Push item onto top of this stack
void IntStack::push(int item) \{
Node *newNode = new Node;
assert (newNode != NULL);
newNode->data = item;
newNode->next \(=\) top;
top \(=\) newNode;
\}
// pop an item off the stack
int IntStack::pop() \{

\section*{Possible Implementations}
- Many possible implementations
- Array-based
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As implementer, use other ADTs to make job easier
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Often simplifies job to reuse pieces when possible

\section*{Stack Via Dynamic Arrays}
```

class IntStack {

```
class IntStack {
public: //same as before
public: //same as before
private:
private:
        int size; // # items currently in stack
        int size; // # items currently in stack
        int capacity; // amount of space allocated
        int capacity; // amount of space allocated
        int *data; // Items in stack are stored
        int *data; // Items in stack are stored
            // in data[0.. size-1].
            // in data[0.. size-1].
                //data[0]is the bottom of the stack;
                //data[0]is the bottom of the stack;
        // data[size-1] is the top item on the stack.
        // data[size-1] is the top item on the stack.
    };
    };
The comments are very important to record how we
The comments are very important to record how we
plan to use the variables
```

plan to use the variables

```
```

Stack Via Arrays (2)
// construct empty IntStack
IntStack::IntStack() {
size = 0;
capacity = DEFAULT_CAPACITY;
data = new int[capacity];
assert(data != NULL);
}
// = "this stack is empty"
bool IntStack::isEmpty() {
return (size == 0);
}

```


Picturing the Implementations

All have the same public interface
Picture the private data
- Vector itself allows more than one implementation!

Stack via Vector


\section*{Discussion}
-Why learn three different ways to implement the same ADT?
-What are the pro's and con's of each way?
- Programming effort?
-Speed (efficiency) of execution?
-Suitability to application?
Other factors?

\section*{Stack Application:}

\section*{Evaluating Expressions}

Expressions like " 3 * (4 + 5)" have to be evaluated by calculators and compilers
We'll look first at another form of expression, called "postfix" or "reverse Polish notation"
Turns out a stack algorithm works like magic to do postfix evaluation
And... another stack algorithm can be used to convert from infix to postfix!

\section*{Postfix vs. Infix}

Review: Expressions have operators (+, -, *, /, etc) and operands (numbers, variables)
- In everyday use, we write the binary operators in between the operands
- " \(4+5\) " means "add 4 and 5" - called infix notation
- No reason why we couldn't write the two operands first, then the operator
-"4 5 +" would mean "add 4 and 5"
- called postfix notation

\section*{More on Postfix}
-3 \(45^{*}\) - means same as (3 (4 5 *) -) -infix: 3-(4*5)
Parentheses aren't needed!
- When you see an operator:
both operands must already be available.
Stop and apply the operator, then go on
Precedence is implicit
-Do the operators in the order found, period!
Practice converting and evaluating:
- \(12+7\) * \(2 \%\)
- \((3+(5 / 3) * 6)-4\)

\section*{Postfix Evaluation via a Stack}

Read in the next "token" (operator or data)
- If data, push it on the data stack
- If (binary) operator (call it "op"):

Pop off the most recent data (B) and next most recent (A)

Perform the operation \(R=A\) op \(B\)
Push R on the stack
Continue with the next token
- When finished, the answer is the stack top.
- Simple, but works like magic!

\section*{Example: 34 5 - *}

Draw the stack at each step!
- Read 3. Push it (because it's data)
- Read 4. Push it.
- Read 5. Push it.
-Read -. Pop 5, pop 4, perform 4-5. Push-1
-Read *. Pop -1, pop 3, perform 3*-1. Push -3.
-No more tokens. Final answer: pop the -3 .
- note that stack is now empty

\section*{Why Postfix?}

Does not require parentheses!
- Some calculators make you type in that way
- Easy to process by a program
- The processing algorithm uses a stack for operands (data)
- simple and efficient

\section*{Refinements and Errors}
- If data stack is ever empty when data is needed for an operation:
- Then the original expression was bad
- Too many operators up to that point
- If the data stack is not empty after the last token has been processed and the stack popped:
- Then the original expression was bad
- Too few operators or too many operands

\section*{Converting in- to post-}

A different algorithm converts from infix to postfix - Uses a stack of operators.

Algorithm:
- Read a token
- If operand, output it immediately
- If '(', push the '(' on stack
- If operator:
if stack top is an op of \(=>\) precedence: pop and output
stop when ' (' is on top or stack empty
push the new operator
If ')', pop and output until '(' has been popped
- Repeat until end of input
pop rest of stack

\section*{Another Stack Application}
- Searching for a path through a maze
-Algorithm: try all possible sequences of locations in the maze until you find one that works (or no more to try)
-called "exhaustive search"
- A stack helps keep track of the possibilities - traces a path of locations
- just like the recursive activation records in the mazesolver

\section*{Stack Wrapup}
- Essence: Last In, First Out
- Various ways to implement
- Numerous uses
- In Computer Science
- In modeling the world```

