

CSE 143 Au00 Additional Problems for Sections November 28, 2000

For each of the following sequences of code, use $O()$ notation to describe the running time in terms of N . Give a lower bound for the running time; don't just answer that all are $O(2^n)$.

```
1.   for (int j = 1; j <= N; j++) {  
       int k = 1;  
       while (k < j)  
           k = k * 2;  
       cout << j << " " << k << endl;  
   }
```

Running time is $O(\underline{n \log n})$

Followup: Suppose the statement $k=k*2$; were replaced with $k=k*3$; . How would this change the running time, if any? **No change; still $O(n \log n)$.**

What do you think this code fragment does? **For each integer from 1 to N , prints that integer and the smallest power of 2 greater than or equal to that integer.**

```
2.   double avg(double m[N][N], int r, int c) {  
       double sum = 0.0;  
       for (int j = r-1; j <= r+1; j++)  
           for (int k = c-1; k <= c+1; k++)  
               sum = sum + m[j][k];  
       return sum/9.0;  
   }
```

```
double puzzle[N][N];  
...  
for (int r = 1; r < N-1; r++)  
    for (int c = 1; c < N-1; c++)  
        puzzle[r][c] = avg(puzzle, r, c);
```

Running time is $O(\underline{n^2})$

Followup: Does it make any difference whether `puzzle` is or is not copied each time the function is called? (i.e., Does it matter if the parameter mechanism is call-by-value instead of call-by-reference?) **Yes. If the array is not copied, each function call requires constant time, and there are $O(n^2)$ function calls. If the array were copied, then each function call would require $O(n^2)$ work by itself. Since there are $O(n^2)$ calls, the total work is $O(n^2)$ times $O(n^2)$, or $O(n^4)$.**