

CSE / ENGR 142

Programming I

Sorting

(corrected 6/1/99)

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i-1

Sorting:

- The problem:

Given an array $a[0], a[1], \dots, a[n-1]$,
reorder entries so that
 $a[0] \leq a[1] \leq \dots \leq a[n-1]$

- Many different ways to do it (algorithms)

- Lots of applications

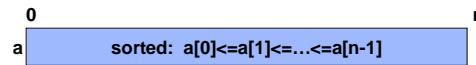
- faster search (allows binary search)
- ordering hits in web search engine
- merging address lists
- etc.

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i-2

Sorting Problem

- What we want: Data sorted in order



- Initial conditions



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Selection Sort

- General situation



- Step:

- Find smallest element x in $a[k..n-1]$
- Swap smallest element with $a[k]$, then increase k



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Subproblem: Find Smallest

```
/* Yield location of smallest element in a[k..n-1] */
int min_loc (int a[], int k, int n) {
    int j, pos; /* a[pos] is smallest element */
    /* found so far */
    pos = k;
    for (j = k + 1; j < n; j = j + 1)
        if (a[j] < a[pos])
            pos = j;
    return pos;
}
```

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Selection Sort

```
/* Sort a[0..n-1] in non-decreasing order (rearrange
```

```
elements in a so that a[0]<=a[1]<=...<=a[n-1] ) */
```

```
int sel_sort (int a[], int n) {
    int k, m;
    for (k = 0; k < n - 1; k = k + 1) {
        m = min_loc(a, k, n);
        swap(&a[k], &a[m]);
    }
}
```

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Example

a [3 | 12 | -5 | 6 | 142 | 21 | -17 | 45]

a [-17 | 12 | -5 | 6 | 142 | 21 | 3 | 45]

a [-17 | -5 | 12 | 6 | 142 | 21 | 3 | 45]

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Example (cont)

a [-17 | -5 | 3 | 6 | 142 | 21 | 12 | 45]

a [-17 | -5 | 3 | 6 | 12 | 21 | 142 | 45]

a [-17 | -5 | 3 | 6 | 12 | 21 | 142 | 45]

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Example (concl)

a [-17 | -5 | 3 | 6 | 12 | 21 | 45 | 142]

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b9

Analysis

- How many steps are needed to sort n things?
 - For each swap, we have to search the remaining array; length is proportional to original array length
 - Need n search/swap operations
 - Total number of steps proportional to n^2
- Can we do better?
 - Selection, insertion, bubble sorts all proportional to n^2
 - Quicksort (CSE143), heap sort, others, proportional to $n \log n$, which is much smaller for big jobs (large n); best we can do in general case
 - Special cases, can do even better: Sort exams by score: drop each exam in one of 101 piles; work is proportional to n

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