

LEC 17

CSE 123**Hashing**

Questions during Class?
Raise hand or send here

sli.do #cse123

**BEFORE WE START*****Talk to your neighbors:****Favorite breakfast cereal?***Respond on sli.do!**

Instructor: Miya Natsuhara**TAs:**

| | | | |
|---------|--------|---------|--------|
| Arohan | Shiven | Yuntong | Anya |
| Sreshta | Vrinda | Amiya | Anisha |
| Rushil | Gavin | Sahana | Trien |
| Sean B | Shreya | Anirudh | Neal |
| Chloe | Jonah | Rohan | Evan |
| Jenny | Renee | Crystal | Rena |
| Nate | Chris | Eeshani | |
| Saachi | Ishita | Prakshi | |
| Hawa | Kuhu | Aidan | |
| Maggie | Kavya | Cora | |
| Sean E | Misha | Nhan | |

Music:  [CSE 123 26sp Lecture Tunes](#) 

Announcements

- Programming Project 3 due today, Friday 5/29 at 11:59pm
- Creative Project 3 out, due Friday, 6/5 at 11:59pm
- Resubmission Cycle 6 due **tonight at 11:59pm**
 - P1, C2, P2 eligible
- R7 / R-Gumball will open on Monday
- Final Exam: Thursday, June 11 12:30pm – 2:20pm
 - Information has been posted on the course website
 - Most of next week's class sessions are dedicated to final exam review!
 - TA-led review session on Tuesday, June 9 4:30pm – 7:00pm in JHN 102
- Gumball and friends visit on Monday, June 8 1:00pm – 2:30pm around Drumheller Fountain 🐾

Data structures so far

- **Lists**

- Maintain an ordered sequence of elements
- Provides `get()`, `add()`, `remove()`, ...
- Studied two implementations: `ArrayList` and `LinkedList`

- **Sets**

- Maintain a collection of elements
- Provides `contains()`, `add()`, `remove()`, ...
- Implementations?

Set implementations: `ArraySet`, `LinkedSet`

| | <code>ArraySet(?)</code> | <code>LinkedSet(?)</code> | <code>SortedArraySet</code> | <code>TreeSet</code> | <code>HashSet</code> |
|-------------------------|--------------------------|---------------------------|-----------------------------|----------------------|----------------------|
| <code>contains()</code> | $O(n)$ | $O(n)$ | | | |
| <code>add()</code> | $O(n)$ | $O(n)$ | | | |
| <code>remove()</code> | $O(n)$ | $O(n)$ | | | |

Set implementations: SortedArraySet

| | ArraySet(?) | LinkedSet(?) | SortedArraySet | TreeSet | HashSet |
|------------|-------------|--------------|----------------|---------|---------|
| contains() | $O(n)$ | $O(n)$ | $O(\log(n))$ | | |
| add() | $O(n)$ | $O(n)$ | $O(n)$ | | |
| remove() | $O(n)$ | $O(n)$ | $O(n)$ | | |

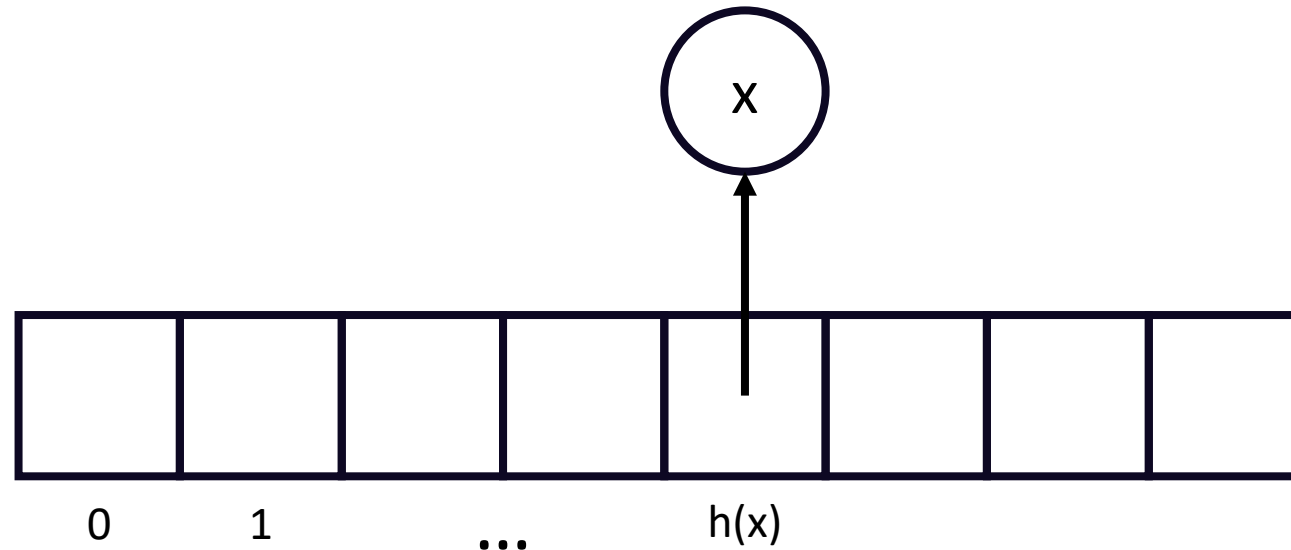
Set implementations: TreeSet

| | ArraySet(?) | LinkedSet(?) | SortedArraySet | TreeSet | HashSet |
|------------|-------------|--------------|----------------|----------------|---------|
| contains() | $O(n)$ | $O(n)$ | $O(\log(n))$ | $O(\log(n))^*$ | |
| add() | $O(n)$ | $O(n)$ | $O(n)$ | $O(\log(n))^*$ | |
| remove() | $O(n)$ | $O(n)$ | $O(n)$ | $O(\log(n))^*$ | |

* assuming tree
is balanced

Hash Table

- Define a **hash function** $h(x)$ that turns any value into an integer
 - Call this the values **hash code**
- Create a big array
- Store each value in the array at index $h(x)$



Set implementations: HashSet

| | ArraySet(?) | LinkedSet(?) | SortedArraySet | TreeSet | HashSet |
|------------|-------------|--------------|----------------|----------------|-------------|
| contains() | $O(n)$ | $O(n)$ | $O(\log(n))$ | $O(\log(n))^*$ | $O(1)^{**}$ |
| add() | $O(n)$ | $O(n)$ | $O(n)$ | $O(\log(n))^*$ | $O(1)^{**}$ |
| remove() | $O(n)$ | $O(n)$ | $O(n)$ | $O(\log(n))^*$ | $O(1)^{**}$ |

* assuming tree
is balanced

** with some
assumptions

What if two items hash to the same slot?

- Nothing stops the hash function from returning the same number for two different items!
- Our hash table needs some way to handle this...
- Any ideas?
- Common solutions:
 - Probing: look at some other related slots
 - Chaining: make each slot a list(!)

What is Linear Probing?

A way to resolve collisions by adding the element in the next available spot

Regular Hash Function

$$h(x) = x \% \text{size}$$

Hash Function (if collision)

$$h'(x) = [h(x) + f(i)] \% \text{size}$$

$$f(i) = i$$

What is Quadratic Probing?

A way to resolve collisions by adding the element in the next available spot (quadratically)

Regular Hash Function

$$h(x) = x \% \text{size}$$

Hash Function (if collision)

$$h'(x) = [h(x) + f(i)] \% \text{size}$$

$$f(i) = i^2$$

What is Chaining?

A way to resolve collisions by creating a LinkedList at that Index (also called a “bucket”)

- Combines both features of ArrayList Indexing and the ease of adding values using LinkedLists

Recap (Comparison)

| Linear Probing | Quadratic Probing | Chaining |
|------------------------------------------------------------------------------|----------------------------------------------------------------------------------------------|---------------------------------------------------------------------------------------------|
| A way to resolve collisions by adding the element in the next available spot | A way to resolve collisions by adding the element in the next available spot (quadratically) | A way to resolve collisions by creating a LinkedList at that Index (also called a “bucket”) |

Why Chaining? Clustering

Clustering - A tendency for data to clump together when using solutions to Collisions like Linear and Quadratic probing

- Linear and Quadratic Probing often result in “Clustering”
- Inefficient use of space in the table
- This means the Runtimes will also be slower

| | | | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|--|--|--|---|--|--|--|--|--|--|---|---|--|--|
| 1 | 2 | 3 | 4 | 5 | 6 | | | | 7 | | | | | | | 8 | 9 | | |
|---|---|---|---|---|---|--|--|--|---|--|--|--|--|--|--|---|---|--|--|

Why Chaining? Efficiency

- Hashing can reduce it down to $O(1)$
- “Load Factor” - lambda (λ)
 - the number of values in each LinkedList
- Finding the index in the Table is $O(1)$
- Finding value in LinkedList is $O(\lambda)$ or essentially $O(1)$

Standard Java hashCode

```
hash = 0  
for (each field) {  
    hash = 31 * hash + hash(field)  
}
```