LEC 11

CSE 123

Binary Tree Modification; Binary Search Trees

Questions during Class?

Raise hand or send here

sli.do #cse123

BEFORE WE START

Talk to your neighbors:

What's your favorite English word?
What page is it on in the dictionary?

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TAs: Trien Nichole Chris Packard Eeshani

Lecture Outline

- **Announcements**
- Binary Tree Modification
- Binary Search Review
- Binary (Search) Trees Review
- More runtime!

Announcements

- Quiz 2 Completed! 😮💨
- Creative Project 2 due today (8/6) at 11:59pm
- Creative Project 3 out tomorrow, due Friday, 8/13 at 11:59pm
- Resubmission Cycle 4 is open, due on Friday, 8/08 at 11:59pm
  - ***C1***, P1, P2 eligible
  - Reminder: In R7, all assignments will be eligible!

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- **Binary Tree Modification**
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Modifying Binary Trees

- Like linked lists, cannot modify nodes
  - Because data field is `final` (there are good reasons for this)
- Will need to create and insert new nodes
- Use `x = change(x)`, usually ***3 times***
  - overall root (in public method)
  - left subtree
  - right subtree
- Order might matter!
  - Does operation on root depend on children?

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Looking through a dictionary

- Assuming a *sorted order* of elements to search through `list`
- Suppose you're looking for a specific element `target`
- Return the index of the given `target`, or -1 if it's not in the `list`

```
begin with the dictionary, from the first to last word,
        looking for target

search(dictionary, left, right, target):
    if there are no more words to look through
        give up
    else
        pick a midpoint between left and right
        pick the word at that midpoint
        if target is that word
            found it!
        else if target comes before that word
            search(dictionary, left, midpoint-1, target)
        else (target comes after that word)
            search(dictionary, midpoint 1, right, target)
```

Binary Search

- Assuming a *sorted order* of elements to search through `list`
- Suppose you're looking for a specific element `target`
- Return the index of the given `target`, or -1 if it's not in the `list`

```
begin with search(list, 0, list.size() – 1, target)

search(list, left, right, target):
    if (left > right):
        return -1
    else:
        mid = (left + right) / 2
        if (target == list[mid]):
            return mid;
        else if (target < list[mid]):
            return search(list, left, mid – 1, target)
        else
            return search(list, mid + 1, right, target)
```

-5	4	18	23	30	49	55	108	184

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- **Binary (Search) Trees Review**
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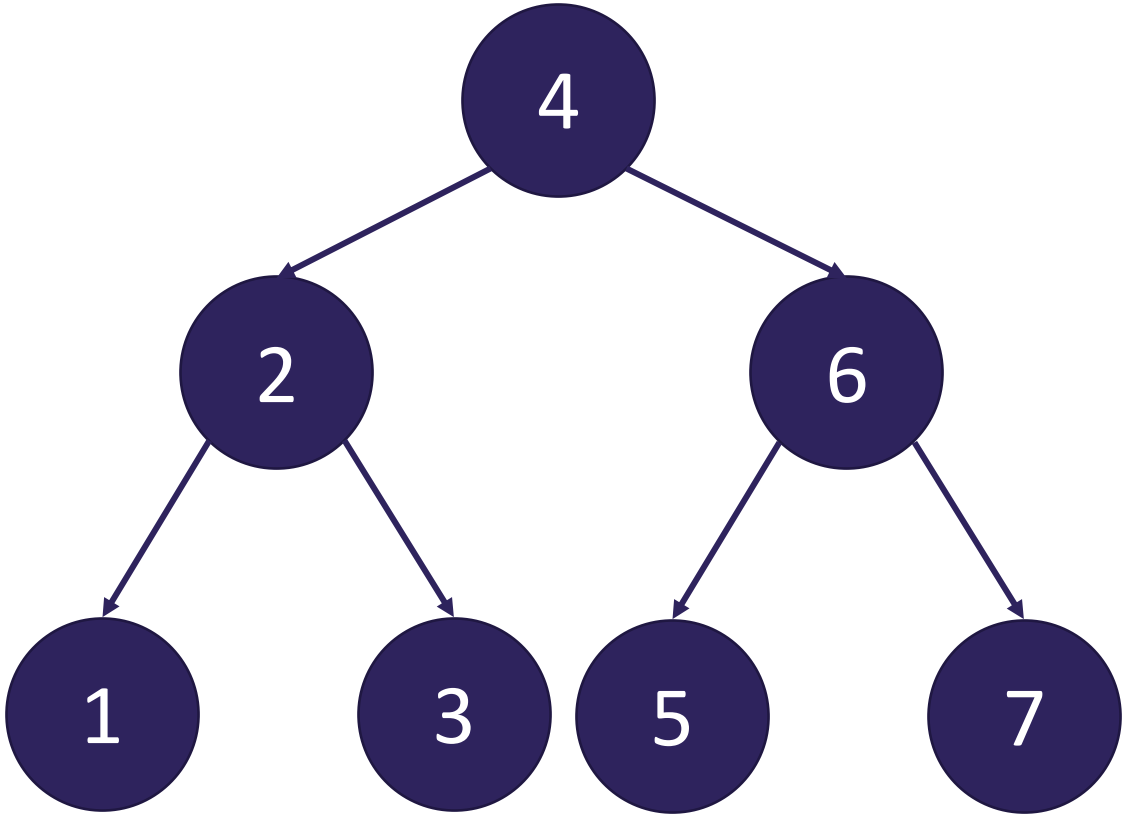
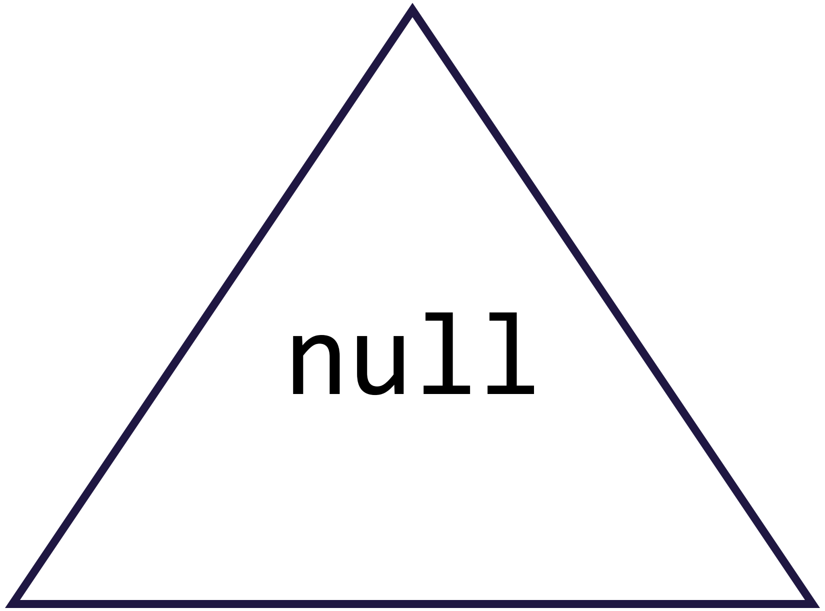
Example Tree: contains



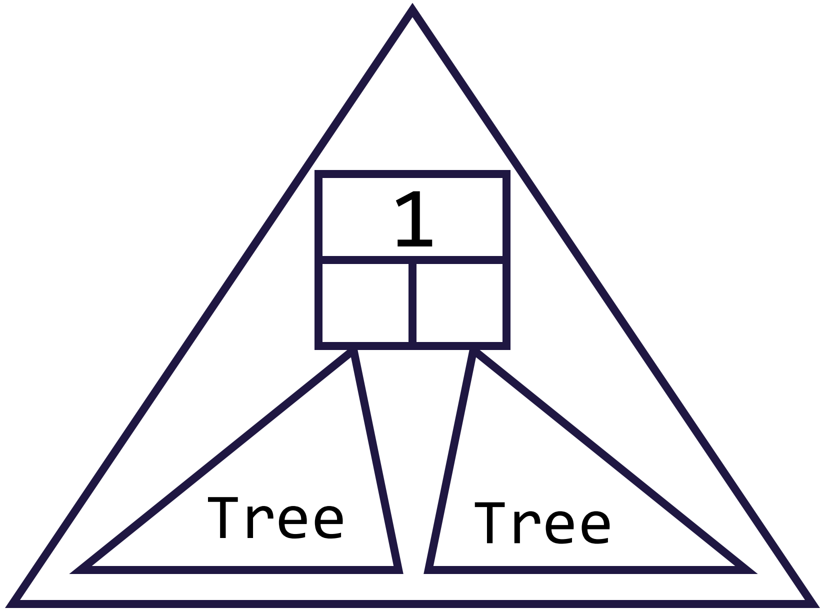
Binary Trees [Review]

- We'll say that any Binary Tree falls into one of the following categories:



Empty tree

`root == null`



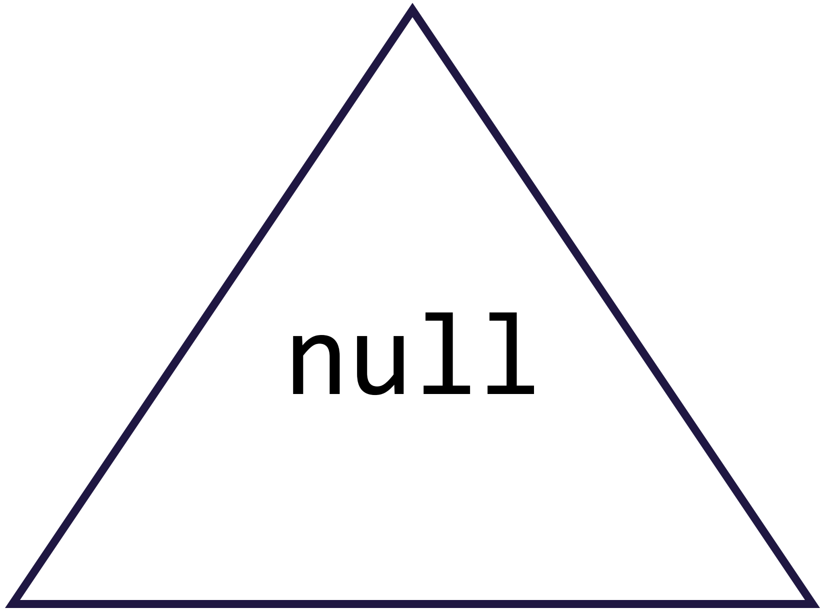
Node w/ two subtrees

`root != null`
`root.left / root.right = Tree`

This is a recursive definition! A tree is either empty or a node with two more trees!

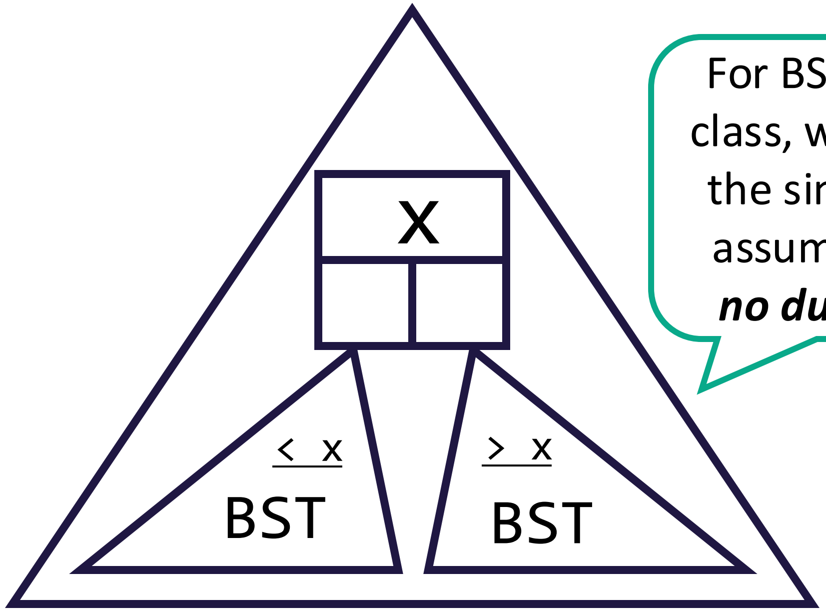
Binary Search Trees (BSTs)

- We'll say that any Binary Search Tree falls into the following categories:



Empty BST

`root == null`



For BSTs in this class, we'll make the simplifying assumption of ***no duplicates***

Node w/ two subBSTs

```
root != null
root.left / root.right = Tree
max(root.left) < x && min(root.right) > x
```

Note that not all Binary Trees are Binary Search Trees

Why BSTs?

- Our `IntTree` implementation to `contains(int value)`

```
private boolean contains(int value, IntTreeNode root) {
    if (root == null) {
        return false;
    } else {
        return root.data == value ||
               contains(value, root.left) ||
               contains(value, root.right);
    }
}
```

- Which direction(s) do we travel if `root.data != value`?
  - Both left and right
- In a Binary Search Tree, should we check both sides?
  - Remember, additional constraint: `max(root.left) < root.data &&`
    `min(root.right) > root.data`

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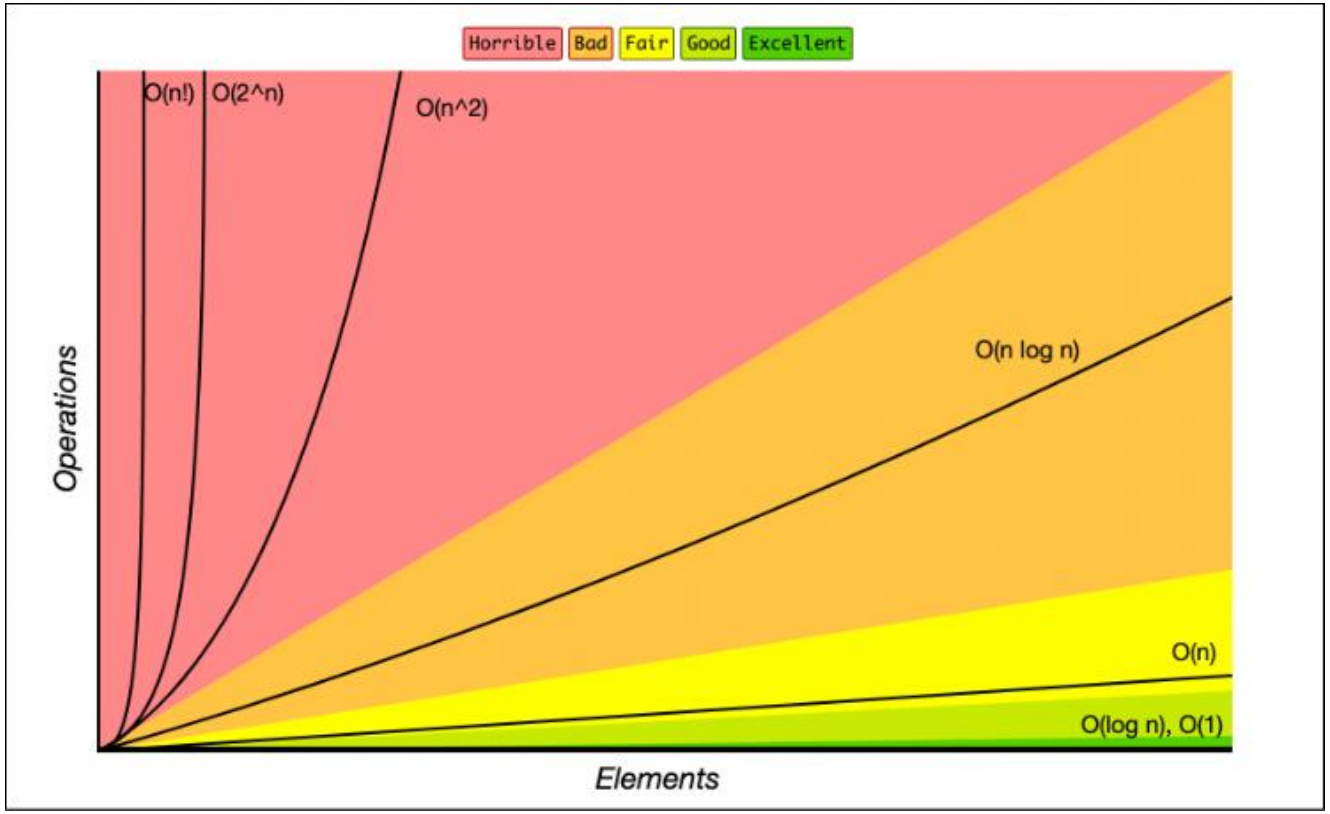
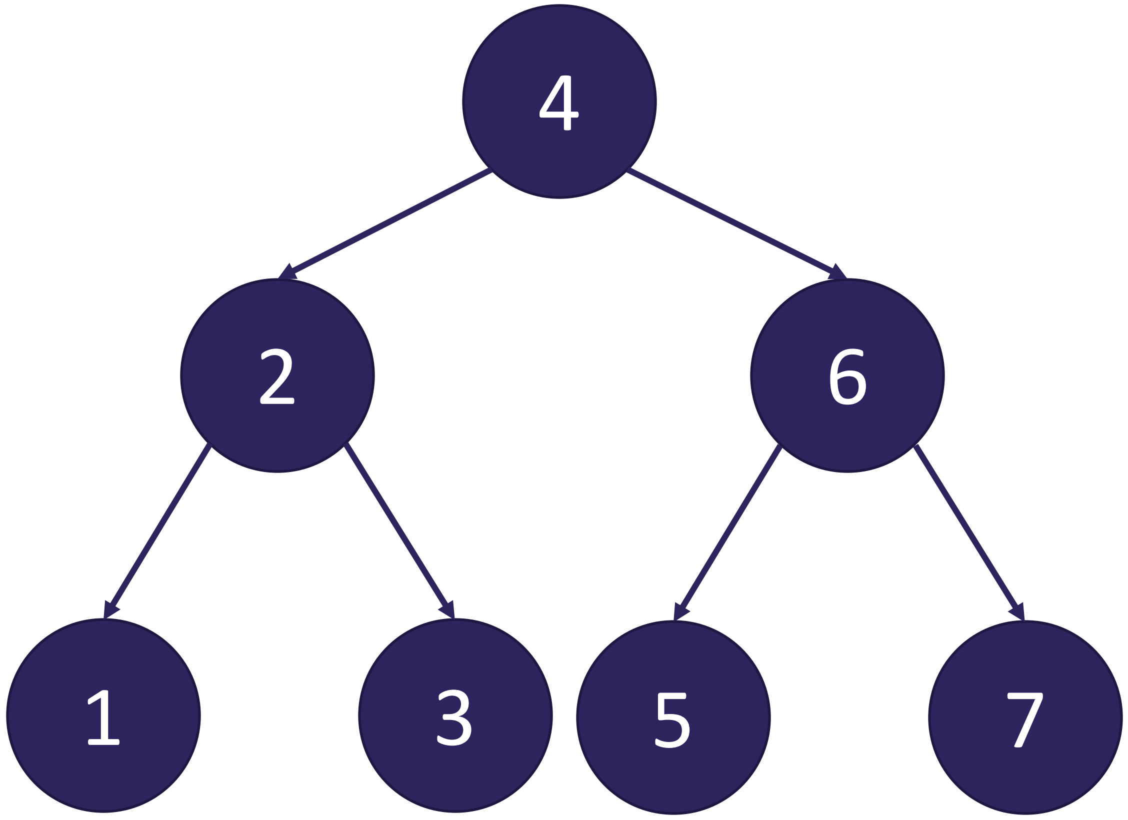
BSTs & Runtime (1)

- Contains operation on a <u>balanced</u> BST runs in `O(log(n))`
  - Leverages removing half of the values at each step
  - *New runtime class unlocked!*



Example Tree: contains for balanced BST



BSTs & Runtime (2)

- Contains operation on a <u>balanced</u> BST runs in `O(log(N))`
  - Leverages removing half of the values at each step
  - *New runtime class unlocked!*

- Comparison between data structures:

Operation	ArrayIntList	LinkedIntList	IntSearchTree
contains(x)	O(N)	O(N)	O(log(N)) ?

BSTs & Runtime (3)

- Contains operation on a balanced BST runs in `O(log(N))`
  - Leverages removing half of the values at each step
  - *New runtime class unlocked!*

- Comparison between data structures:

Operation	ArrayIntList	LinkedIntList	IntSearchTree
contains(x)	O(N)	O(N)	***O(N)***

`O(Log(N))` runtime is only guaranteed for **BALANCED** *BSTs. If your tree isn't balanced, we see `O(N)` runtime!*

BSTs In Java

- Self-balancing BST implementations (AVL / Red-black) exist
  - AVL better at contains, Red-black better at adding / removing
- Both the `TreeMap` / `TreeSet` implementations use self-balancing BSTs
  - Determines said ordering via the `Comparable` interface / `compareTo` method
  - Printing out shows natural ordering – preorder traversal

- Complete table comparing data structures:

Operation	ArrayList	LinkedList	TreeSet
contains(x)	O(N)	O(N)	O(log(N))
add(x)	O(1*)	O(1)	O(log(N)*)
remove(x)	O(N)	O(N)	O(log(N)*)

**It's slightly more complicated but we'll leave that for a higher level course*