#### **BEFORE WE START**

Talk to your neighbors:

How do you remember where you put things?

#### Instructor: Nathan Brunelle

Arohan	Ashar	Neha	Rohini	Rushil
TAS: Ido	Zachary	Sebastian	Joshua	Sean
Hayden	Caleb	Justin	Heon	Rashad
Srihari	Benoit	Derek	Chris	Bhaumik
Kuhu	Kavya	Cynthia	Shreya	Ashley
Ziao	Kieran	Marcus	Crystal	Eeshani
Prakshi	Packard	Cora	Dixon	Nichole
Niyati	Trien	Lawrence	Evan	Cady

CSE 123 Hashing

**Questions during Class?** 

Raise hand or send here

sli.do #cse123A

#### Announcements

- Programming Project 3 due today, Friday 5/30 at 11:59pm
- Creative Project 3 out, due Friday, June 6 at 11:59pm
- Resubmission Cycle 6 due tonight at 11:59pm
  - <u>**P1**</u>, C2, P2 eligible
- R7 / R-Bucks will open on Monday
- Final Exam: Wednesday, June 11 at 12:30pm 2:20pm
  - <u>Left-handed desk request form</u>, closes Tuesday, June 3
  - Details and resources posted later today

# Data structures so far

#### • Lists

- Maintain an ordered sequence of elements
- Provides get(), add(), remove(), ...
- Studied two implementations: ArrayIntList and LinkedIntList

#### • Sets

- Maintain a collection of elements
- Provides contains(), add(), remove(), ...
- Implementations?

	ArraySet(?)	LinkedSet(?)	TreeSet	HashSet
contains()	O(n)	O(n)		
add()	O(n)	O(n)		
remove()	O(n)	O(n)		

	ArraySet(?)	LinkedSet(?)	TreeSet	HashSet
contains()	O(n)	O(n)	O(log(n))*	
add()	O(n)	O(n)	O(log(n))*	
remove()	O(n)	O(n)	O(log(n))*	

\* assuming tree is balanced

## Hash Table

- Define a hash function h(x) that turns any value into an integer
  - Call this the values hash code
- Create a big array
- Store each value in the array at index h(x)



	ArraySet(?)	LinkedSet(?)	TreeSet	HashSet
contains()	O(n)	O(n)	O(log(n))*	
add()	O(n)	O(n)	O(log(n))*	
remove()	O(n)	O(n)	O(log(n))*	

\* assuming tree is balanced

	ArraySet(?)	LinkedSet(?)	TreeSet	HashSet
contains()	O(n)	O(n)	O(log(n))*	O(1)**
add()	O(n)	O(n)	O(log(n))*	O(1)**
remove()	O(n)	O(n)	O(log(n))*	O(1)**
			<ul> <li>* assuming tree</li> <li>is balanced</li> </ul>	<pre>** with some assumptions</pre>

# What is Linear Probing?

A way to resolve collisions by adding the element in the next available spot

**Regular Hash Function** 

Hash Function (if collision)

h(x) = x % size

h'(x) = [h(x) + f(i)] % size f(i) = i

# What is Linear Probing?

h(x) = x % size h'(x) = [h(x) + f(i)] % size f(i) = i



# What is Quadratic Probing?

A way to resolve collisions by adding the element in the next available spot (quadratically)

**Regular Hash Function** 

Hash Function (if collision)

h(x) = x % size

h'(x) = [h(x) + f(i)] % size f(i) = i<sup>2</sup>

# What is Quadratic Probing?

h(x) = x % size h'(x) = [h(x) + f(i)] % size  $f(i) = i^2$ 

# What is <a href="https://www.chaining">Chaining?</a>

A way to resolve collisions by creating a LinkedList at that Index (also called a "bucket")

• Combines both features of ArrayList Indexing and the ease of adding values using LinkedLists

# What is Chaining? Chaing?

h(x) = x % size



## Recap (Comparison)

Linear Probing	Quadratic Probing	📌 Chaining
A way to resolve collisions by adding the element in the next available spot	A way to resolve collisions by adding the element in the next available spot (quadratically)	A way to resolve collisions by creating a LinkedList at that Index (also called a "bucket")

# Why **+** Chaining?

Clustering - A tendency for data to clump together when using solutions to Collisions like Linear and Quadratic probing

- Linear and Quadratic Probing often result in "Clustering"
- Inefficient use of space in the table
- This means the Runtimes will also be slower

# Why **+** Chaining?

- Hashing can reduce it down to O(1)
- "Load Factor" lambda (λ)
  - the number of values in each LinkedList
- Finding the index in the Table is O(1)
- Finding value in LinkedList is  $O(\lambda)$  or essentially O(1)

#### Standard Java hashCode