Comparable + Huffman

Hitesh Boinpally Summer 2023



Agenda

- Comparable
- P3 Huffman



Agenda

- Comparable 🖛
- P3 Huffman



Interfaces Review

- Define set of behavior a class can implement
- "Contract", "Certification"
- List and Set are some interfaces we've used
- To utilize:
 - public class Tesla implements Car
- Says that Tesla "implements" the Car interface





The Comparable Interface

- Say you had a FootballTeam class and wanted to sort them
 - How could you tell Java how to sort them?





The Comparable Interface

- Say you had a FootballTeam class and wanted to sort them
 - How could you tell Java how to sort them?
- Utilize the Comparable interface!
 - Fixed definition of how to compare objects
 - This is needed to allow objects to be added to TreeSets and TreeMaps





Agenda

• Comparable

• P3 - Huffman 🖛





Priority Queues and Huffman Encoding Introduction to the Final Project

Hunter Schafer

CSE 143, Autumn 2021

Priority Queue

A collection of ordered elements that provides fast access to the minimum (or maximum) element.

public class PriorityQueue<E> implements Queue<E>

PriorityQueue <e>()</e>	constructs an empty queue
add(E value)	adds value in sorted order to the queue
peek()	returns minimum element in queue
remove()	removes/returns minimum element in queue
size()	returns the number of elements in queue

```
Queue<String> tas = new PriorityQueue<String>();
tas.add("Watson");
tas.add("Sherlock");
tas.remove();
```

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size()	returns the number of elements in queue

```
Queue<String> tas = new PriorityQueue<String>();
tas.add("Watson");
tas.add("Sherlock");
tas.remove(); // "Sherlock"
```

Final Project: Huffman Coding

File Compression

Compression

Process of encoding information so that it takes up less space.

Compression applies to many things!

- Store photos without taking up the whole hard-drive
- Reduce size of email attachment
- Make web pages smaller so they load faster
- Make voice calls over a low-bandwidth connection (cell, Skype)

Common compression programs:

• WinZip, WinRar for Windows



ASCII (American Standard Code for Information Interchange)

Standardized code for mapping characters to integers

- Many text files on your computer are in ASCII.
- But, computers need numbers represented in binary!

Character	ASCII value
6 9	32
'a'	97
ʻb'	98
'C'	99
'e'	101
'Z'	122

ASCII (American Standard Code for Information Interchange)

Standardized code for mapping characters to integers

- Many text files on your computer are in ASCII.
- But, computers need numbers represented in binary!

Every character is represented by a byte (8 bits).

Character	ASCII value	Binary Representation
6 9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
'C'	99	01100011
'e'	101	01100101
'Z'	122	01111010

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻc'	99	01100011
'e'	101	01100101
'Z'	122	01111010

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
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'e'	101	01100101
"Z'	122	01111010

<u>c</u>ab z

Answer 01100011

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
'Z'	122	01111010

c<u>a</u>b z

Answer

01100011 01100001

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
'Z'	122	01111010

ca<u>b</u> z

Answer 01100011 01100001 01100010

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
"Z'	122	01111010

 cab_z

Answer

01100011 01100001 01100010 00100000

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
"Z'	122	01111010

cab <u>z</u>

Answer

01100011 01100001 01100010 00100000 01111010

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
'c'	99	01100011
'e'	101	01100101
'Z'	122	01111010

Answer

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
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'C'	99	01100011
'e'	101	01100101
'Z'	122	01111010

How do we read the following binary as ASCII? 011000010110001101100101

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
'Z'	122	01111010

How do we read the following binary as ASCII? 01100001 01100011 01100101

Answer

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
'C'	99	01100011
'e'	101	01100101
'Z'	122	01111010

How do we read the following binary as ASCII? 01100001 01100011 01100101

Answer

а

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
'C'	99	01100011
'e'	101	01100101
'Z'	122	01111010

How do we read the following binary as ASCII? 01100001 <u>01100011</u> 01100101

Answer

ac

Character	ASCII value	Binary Representation
6.9	32	00100000
'a'	97	01100001
ʻb'	98	01100010
ʻC'	99	01100011
'e'	101	01100101
'Z'	122	01111010

How do we read the following binary as ASCII? 01100001 01100011 <u>01100101</u>

Answer

ace

Huffman's Insight

Use variable length encodings for different characters to take advantage of frequencies in which characters appear.

- Make more frequent characters take up less space.
- Don't have codes for unused characters.
- Some characters may end up with longer encodings, but this should happen infrequently.

Huffman Encoding

- Create a "Huffman Tree" that gives a good binary representation for each character.
- The path from the root to the character leaf is the encoding for that character; left means 0, right means 1.

Character	Binary Representation
6.9	00100000
'a'	01100001
ʻb'	01100010
'C'	01100011
'e'	01100101
'Z'	01111010

ASCII Table

Huffman Tree



The final project asks you to write a class that manages creating and using this Huffman code.

- (A) Create a Huffman Code from a file and compress it.
- (B) Decompress the file to get original contents.

Input File Contents

bad cab

Input File Contents

bad cab

Step 1: Count the occurrences of each character in file

{ ' '=1, 'a'=2, 'b'=2, 'c'=1, 'd'=1}

Input File Contents

bad cab

Step 1: Count the occurrences of each character in file

{ ' '=1, 'a'=2, 'b'=2, 'c'=1, 'd'=1}

Step 2: Make leaf nodes for all the characters. Place in a PriorityQueue

$$\mathsf{pq} \leftarrow \boxed{ \left(\begin{array}{c} \overbrace{freq: 1} \\ \hline{freq: 1} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 1} \\ \hline{freq: 1} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 1} \\ \hline{freq: 1} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \overbrace{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \\ \hline{freq: 2} \end{array} \right) \left(\begin{array}{c} \hline{freq: 2} \end{array} \right)$$

Input File Contents

bad cab

Step 1: Count the occurrences of each character in file

{ ' '=1, 'a'=2, 'b'=2, 'c'=1, 'd'=1}

Step 2: Make leaf nodes for all the characters. Place in a PriorityQueue

$$pq \leftarrow \boxed{\begin{pmatrix} \cdot & \cdot \\ \hline freq: 1 \end{pmatrix}} \left(\begin{matrix} \hline \frac{\cdot c_{c}'}{freq: 1} \end{matrix} \right) \left(\begin{matrix} \hline \frac{\cdot c_{d}'}{freq: 1} \end{matrix} \right) \left(\begin{matrix} \hline \frac{\cdot c_{d}'}{freq: 2} \end{matrix} \right) \left(\begin{matrix} \hline \frac{\cdot c_{b}'}{freq: 2} \end{matrix} \right) \leftarrow$$

Step 3: Use Huffman Tree building algorithm (described soon)

Input File Contents

bad cab

Step 1: Count the occurrences of each character in file

{' '=1, 'a'=2, 'b'=2, 'c'=1, 'd'=1}

Step 2: Make leaf nodes for all the characters. Place in a PriorityQueue

$$\mathsf{pq} \leftarrow \boxed{\begin{pmatrix} & \cdot & \cdot \\ \hline freq: & 1 \end{pmatrix}} \left(\begin{matrix} \frac{\mathsf{r}_{\mathsf{C}'}}{freq: & 1} \end{matrix} \right) \left(\begin{matrix} \frac{\mathsf{r}_{\mathsf{C}'}}{freq: & 1} \end{matrix} \right) \left(\begin{matrix} \frac{\mathsf{r}_{\mathsf{C}'}}{freq: & 2} \end{matrix} \right) \left(\begin{matrix} \frac{\mathsf{r}_{\mathsf{D}'}}{freq: & 2} \end{matrix} \right) \leftarrow$$

Step 3: Use Huffman Tree building algorithm (described soon)

Step 4: Save encoding to .code file to encode/decode later.

{'d'=00, 'a'=01, 'b'=10, ' '=110, 'c'=111}

Input File Contents

bad cab

Step 1: Count the occurrences of each character in file

{' '=1, 'a'=2, 'b'=2, 'c'=1, 'd'=1}

Step 2: Make leaf nodes for all the characters. Place in a PriorityQueue

$$\mathsf{pq} \leftarrow \boxed{\left(\overbrace{freq: 1}^{i \ \cdot \ \cdot} \right) \left(\overbrace{freq: 1}^{i \ c'} \right) \left(\overbrace{freq: 1}^{i \ c'} \right) \left(\overbrace{freq: 1}^{i \ c'} \right) \left(\overbrace{freq: 2}^{i \ c'} \right) \left(\overbrace{freq: 2}^{i \ c'} \right)} \leftarrow$$

Step 3: Use Huffman Tree building algorithm (described soon)

Step 4: Save encoding to .code file to encode/decode later.

{ 'd'=00, 'a'=01, 'b'=10, ' '=110, 'c'=111 }

Step 5: Compress the input file using the encodings Compressed Output: 1001001101110110 We do this step for you

Input File		
bad cab		

Generate Counts Array:



- Store each character and its frequency in a HuffmanNode object.
- Place all the HuffmanNodes in a PriorityQueue so that they are in ascending order with respect to frequency

$$\mathsf{pq} \leftarrow \boxed{\left(\begin{array}{c} \overset{\cdot}{}, \\ \hline freq: 1 \end{array}\right) \left(\begin{array}{c} \overset{\cdot}{} \overset{\cdot}{} \\ \hline freq: 1 \end{array}\right) \left(\begin{array}{c} \overset{\cdot}{} \overset{\cdot}{} \\ \hline freq: 1 \end{array}\right) \left(\begin{array}{c} \overset{\cdot}{} \overset{\cdot}{} \\ \hline freq: 2 \end{array}\right) \left(\begin{array}{c} \overset{\cdot}{} \overset{\cdot}{} \\ \hline freq: 2 \end{array}\right)} \leftarrow$$



















• What is the relationship between frequency in file and binary representation length?

Algorithm Pseudocode

```
while P.Q. size > 1:
    remove two nodes with lowest frequency
    combine into a single node
    put that node back in the P.Q.
```



Save the tree to a file to save the encodings for the characters we made.



Output of save



Output of save
100
00



Output of save
100
00
97
01



Output of save
100
00
97
01
98
10



Output of save
100
00
97
01
98
10
32
110



Output of save
100
00
97
01
98
10
32
110
99
111

We do this step for you

Input File	Huffman Encoding
bad cab	100
Comproseed Output	00
compressed output	97
	01
	98
	10
	32
	110
	99
	111

We do this step for you

Input File	Huffman Encoding
bad cab	100 'd'
Comprossed Output	00
compressed output	97 'a'
	01
	98 <mark>'b</mark> '
	10
	32 ''
	110
	99 'c'
	111

We do this step for you

Input File	Huffman Encoding
bad cab	100 ' <mark>d</mark> '
Compressed Output	00
10 01 00 110 111 01 10	97 'a'
	98 'b'
	10
	32 ''
	110
	99 'c'
	111

We do this step for you

Input File	Huffman Encoding
bad cab	100 <mark>'d</mark> '
Commenced Output	00
Compressed Output	97 'a'
10 01 00 110 111 01 10	01
Uncompressed Output	50 D
01100010 01100001 01100100	10
01100010 01100001 01100100	32 ''
00100000 01100011 01100001	110
01100010	99 'c'
	111

Step 1: Reconstruct the Huffman tree from the code file Step 2: Translate the compressed bits back to their character values.

Input code File	
97	
0	
101	
100	
32	
101	
112	
11	

Initially the tree is empty



Input code File	
97	
0	
101	
100	
32	
101	
112	
11	

Tree after processing first pair



Input code File	
97	
0	
101	
100	
32	
101	
112	
11	

Tree after processing second pair



Input code File	
97	
0	
101	
100	
32	
101	
112	
11	

Tree after processing third pair



Input code File	
97	
0	
101	
100	
32	
101	
112	
11	

Tree after processing last pair



After building up tree, we will read the compressed file bit by bit.

Input 0101110110101011100

Output



After building up tree, we will read the compressed file bit by bit.

Input 0101110110101011100

Output

a papa ape



Reading bits in Java is kind of tricky, we are providing a class to help!

BitInputStream(String file)	Creates a stream of bits from file
hasNextBit()	Returns true if bits remain in the stream
nextBit()	Reads and returns the next bit in the stream

public class BitInputStream

Review - Final Project

Part A: Compression

public HuffmanCode(int[] counts)

• Slides 11-13

public void save(PrintStream out)

• Slide 14

Part B: Decompression

public HuffmanCode(Scanner input)

• Slide 17

```
public void translate(BitInputStream in,
```

PrintStream out)

• Slide 18