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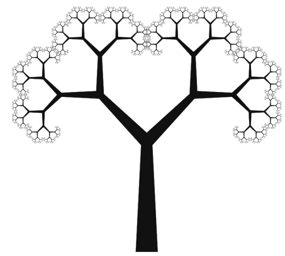
# CSE 143

## Computer Programming II



# Binary Search Trees (BSTs)



## Outline



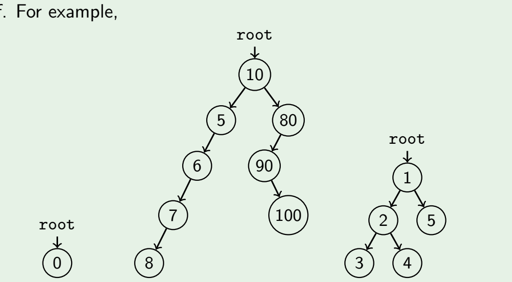
## Another Tree Method 

**`height`**

Write a tree method called `height` (inside the `IntTree` class) with the following method signature:

`public int height()`

that returns the number of nodes on the **longest path** from the root to any leaf. For example,




height is 1 — height is 5 — height is 3

## `height` Solution 

```java
public int height() {
   return height(this.root);
}

private int height(IntTreeNode current) {
   // A null tree has height 0
   if (current == null) {
      return 0;
   }
   else {
      // Find the largest path by taking the max
      // of both branches recursively (and adding 1 for this node)
      return 1 + Math.max(
         height(current.left),
         height(current.right)
      );
   }
}
```

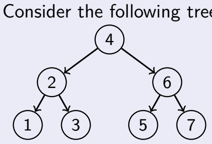
## Back to `contains` 

**Recall `contains()`**

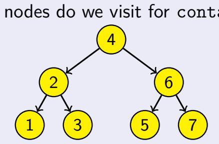
```java
private boolean contains(IntTreeNode current, int value) {
   /* If the tree is null, it definitely doesn't contain value... */
   if (current == null) { return false; }

   /* If current *is* value, we found it! */
   else if (current.data == value) { return true; }

   else {
      return contains(current.left, value) ||
             contains(current.right, value);
   }
}
```

**Runtime of `contains(7)`**

Consider the following tree:




Which nodes do we visit for `contains`(**7**)




**That makes the code $\mathcal{O}(n)$. Can we do better?**

## Doing Better! 

In general, **we can't do better**. BUT, sometimes, we can!

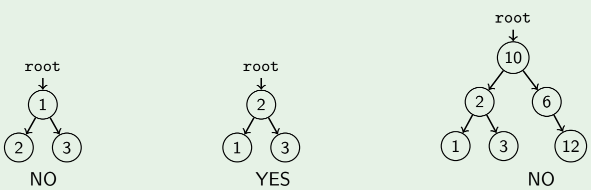
**Definition (Binary SEARCH Tree (BST))**

A binary tree is a **BST** when an **in-order traversal of the tree** yields a sorted list.

To put it another way, a binary tree is a **BST** when:

- All data "to the left of" a node is less than it
- All data "to the right of" a node is greater than it
- All sub-trees of the binary tree are also BSTs

**Example (Which of the following are BSTs?)**




## Checking `isBST` 

**`isBST`**

Write a function `isBST` with the following signature:

```
private boolean isBST(IntTreeNode current, Integer min,
                      Integer max)
```

that returns true if the tree at root `current` is a BST.

```
private boolean isBST(IntTreeNode current, Integer min, Integer max) {
   if (current == null) {
      return true;
   }
   else if ((min != null && current.data < min) ||
            (max != null && current.data > max)) {
      return false;
   }
   else if (!isBST(current.left, min, current.data)) {
      return false;
   }
   else {
      return isBST(current.right, current.data, max);
   }
}
```

## `contains` (AGAIN!) 

**Write `contains()` for a BST**

Fix `contains` so that it takes advantage of the BST properties.

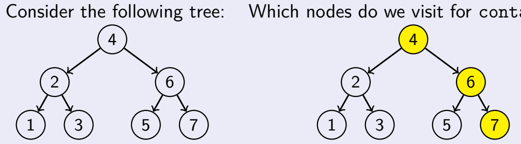
**Recall `contains()`**

```
private boolean contains(IntTreeNode current, int value) {
   /* If the tree is null, it definitely doesn't contain value... */
   if (current == null) { return false; }

   /* If current *is* value, we found it! */
   else if (current.data == value) { return true; }

   else if (current.data < value) {
      return contains(current.right, value);
   }
   else {
      return contains(current.left, value);
   }
}
```

## Tracing the new `contains` 

**Runtime of (better) `contains(7)`**

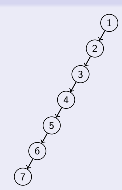
Consider the following tree: Which nodes do we visit for `contains(7)`




**That makes the code $\log n$. Much better!**

**WARNING!**

Consider the following tree:




This is the same tree, but now **we have to visit all the nodes!**

## Adding to a BST! 

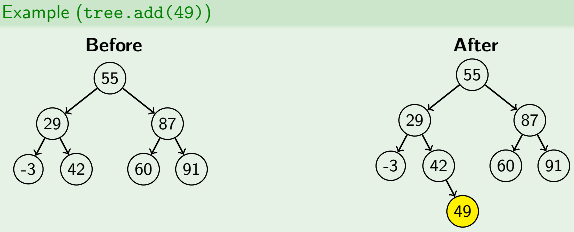
**`add`**

Write a method add in the BST class with the following signature:

```
public void add(int value);
```

that preserves the BST property.

**Example (`tree.add(49)`)**




## Adding to a BST (Attempt #1) 

**Attempt #1**

```
public void add(int value) {
   add(this.root, value);
}
private void add(Node current, int value) {
   if (current == null) {
      current = new Node(value);
   }
   else if (current.data > value) {
      add(current.left, value);
   }
   else if (current.data < value) {
      add(current.right, value);
   }
}
```

**What's wrong with this solution?**

Just like with `LinkedLists` where we must change `front` or `.next`, we're not actually changing anything here. We're discarding the result.

## x = change(x)



Consider the following code:

```
public static void main(String[] args) {
   String s = "hello world";
   s.toUpperCase();
   System.out.println(s);
}
```

OUTPUT

```
>> hello world
```

```
public static void main(String[] args) {
   String s = "hello world";
   s = s.toUpperCase();
   System.out.println(s);
}
```

OUTPUT

```
>> HELLO WORLD
```

**We must USE the result; otherwise, it gets discarded**

## x = change(x)



If you want to write a method that can change the object that a variable refers to, you must do three things:

1. Pass in the original state of the object to the method
2. Return the new (possibly changed) object from the method
3. Re-assign the caller's variable to store the returned result

```
p = change(p); // in main
public static Point change(Point thePoint) {
   thePoint = new Point(99, -1);
   return thePoint;
}
```

## Adding to a BST (Fixed)



Fixed Attempt

```
public void add(int value) {
   this.root = add(this.root, value);
}
private Node add(Node current, int value) {
   if (current == null) {
      current = new Node(value);
   }
   else if (current.data > value) {
      current.left = add(current.left, value);
   }
   else if (current.data < value) {
      current.right = add(current.right, value);
   }
   return current;
}
```

This works because we **always update the result**, **always return the result**, and **always update the root**.

## BST Tips!



- BSTs can make searching/inserting/etc. much faster.
- Make sure that you can figure out if a tree is a BST or not.
- Whenever you are writing a BST method, you **must** use the `x = change(x)` pattern. It won't work otherwise.